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THE MINIMUM MAXIMAL-CUBE COVERING APPROACH TO SWITCHING
THEORY

New Jersey Institute of Technology

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THE MINIMUM MAXIMAL-CUBE COVERING

APPROACH TO SWITCHING THEORY

BY

YU-TSANG HWANG

A DISSERTATION

PRESENTED IN PARTIAL FULFILLMENT OF

THE REQUIREMENTS FOR THE DEGREE

OF

DOCTOR OF SCIENCE IN ELECTRICAL ENGINEERING

AT

NEW JERSEY INSTITUTE OF TECHNOLOGY

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Newark, New Jersey

1980

APPROVAL OF DISSERTATION

THE MINIMUM MAXIMAL-CUBE COVERING
APPROACH TO SWITCHING THEORY

BY

YU-TSANG HWANG

FOR

DEPARTMENT OF ELECTRICAL ENGINEERING
NEW JERSEY INSTITUTE OF TECHNOLOGY

BY

FACULTY COMMITTEE

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ABSTRACT

This dissertation is devoted to the construction of a new algebraic structure for the K-map so that algorithms and theorems can be derived for the minimization of Boolean functions and at the same time provide for the manipulation of Boolean functions as freely as one could on the K-map.

Computer programs have been included to demonstrate the theory and numerous examples of various complexity have been worked out in the appendices.

PREFACE

With the advent of microprocessors and programmable logical arrays, low speed logic circuits can now be realized without performing the tedious and complicated process associated with switching-function minimization. However, as speeds and complexities of switching functions become ever faster and more complex, the classical two-level random logic approach, far from being obsolete, deserves serious consideration. Having the least propagation delay is sufficient advantage for the two-level AND/OR gating logic to be considered in future switching circuits. In an ultra-fast system, if propagation delays are of the same order of magnitude as pulse widths, as the clock period approaches twice the smallest pulse width, the system cycle time can actually become faster than the execution time of a microprocessor or the access time of a logic array. When this happens, two-level random logic is the only alternative. In the report of a study conducted by Rudolf E. Thun and Edward T. Lewis¹ at Raytheon Company, the impact of gigahertz logic on system design is discussed. Berthod G. Bosch² also did a study on the prospect of ultra-fast logic systems in the 1980's. These studies indicate that the coming of ultra-fast logic systems in the 1980's is clear and certain.

Research ranging from the hardware packaging to the new scope of the applications is in full swing; the Military VLSI ultra-fast speed logic program is also progressing swiftly.

The most recent attempts in the classical two-level logic minimizations are to unify the two steps of Quine's original concepts, i.e. to search the minimum set directly without finding the prime implicant tables. McKinney's³ attempt with the direct search method, Curtis'⁴ with the adjacent table method and Arevalo's⁵ four-tuple method are all in this trend. The most natural and logical route is actually the algebraic equivalent to the Karnaugh-map method. However, the lack of an algebraic structure prevents one from manipulating the K-map analytically.

This dissertation is devoted to the construction of a new algebraic structure for the K-map so that algorithms and theorems can be derived for the minimization of Boolean functions and at the same time provide for the manipulation of Boolean functions as freely as one could on the K-map. We call it the minimum maximal-cube covering approach.⁶

In Chapter One, a brief historical survey and review of switching circuits, Boolean algebra and Boolean function minimization is presented. The comparison of

various approaches to the minimization of switching circuits leads to a set of most desirable characteristics of the mechanics of minimization of switching circuits.

Chapter Two is devoted to the construction of the algebraic structure and the theoretical foundation of the minimum maximal-cube covering approach to switching theory. Starting from definitions and theorems of minimum maximal-cube covering approach, a deeper and richer insight of switching theory is shown. Applications and examples will be demonstrated in Chapters Three, Four and Five.

Chapter Three covers the algorithms of a cube graph approach with demonstrated example problems.

Chapter Four covers a different and more computer-oriented algorithm, the algebraic-geometric approach. Examples, problems and demonstrative computer programs are also shown in Chapter Four.

In Chapter Five, applications in other areas are explored - the decomposition of Boolean functions, the symmetry of Boolean functions, the sequential circuits, the coding theory, et cetera.

In Chapter Six, the concluding chapter, we summarize the results and also look beyond this dissertation to the prospective of minimum maximal cube covering approach.

- ¹ Thun, E. Rudolf and Edward T. Lewis, "Advanced Packaging Techniques, Final Report on Task I" AFAL-TR-77-165, part I, Raytheon Company, Missile Systems Division, Hartwell Road, Bedford, Massachusetts, 01730 pp. 10-16, August 1977.
- ² Bosch, B. G., "Gigabit Electronics", Proc. 1977 Europ. Microwave Conference, Copenhagen. (invited paper) pp.1-6.
- ³ McKinney, M. H., "A directed search algorithm for the canonical minimization of switching functions", Ph.D. dissertation, Texas A&M University, College Station, TX, 1974.
- ⁴ Curtis, H. A., "Adjacency table Method of Deriving Minimal Sums," IEEE Transaction on Computers Vol. C-26, No. 11, Nov. 1977, pp. 1136-1141.
- ⁵ Arevalo, Z and Bredeson, J. G. "A Method to Simplify a Boolean Function into a near minimal sum-of-products for programmable logic arrays", IEEE Transaction of Computers Vol C-27, No. 11, Nov. 1978, pp. 1026-1039.
- ⁶ For detailed explanation of the meaning of Minimum Maximal-Cube Covering, please refer to Chapter 1 and Chapter 2.

DEDICATION

To my wife, Rene, for her patience, her understanding, and her encouragement.

ACKNOWLEDGEMENTS

The author gratefully acknowledges the guidance and assistance of Professor Robert E. Anderson. He also expresses his gratitude to Drs. Kenneth Sohn, William J. Troop, and Simon Cohen.

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CHAPTER I

HISTORICAL REVIEW OF SWITCHING THEORY AND MINIMIZATION

1.1 Historical Review

At the beginning, primitive mathematical ideas tend to arise from the need of a better way to describe and analyze some special physical system; then through mathematician's hands it is further developed by deductive reasoning into a new branch of mathematics. Engineering tools for analysis and design of apparatus, starting from craftsman's rudimentary cut-and-try method to the borrowing of rigorous mathematics for refined and systematic synthesis, coincides with the progress of technological demands. In the 1680's, G.W. Leibniz' attempts to create a systematic mathematical logic led to the discovery of some of the basic Boolean algebra.¹ His contemplation of a calculating machine, although not rendered into hardware, was the starting point of modern digital systems and computers. During the following 200 years, little progress was made in mathematical logic. In 1854, British mathematician George Boole published his work, "An Investigation of the Laws of Thoughts"², which marked the beginning of modern symbolic logic and Boolean Algebra. Modern telegraphy and control systems have been using relay switching circuits since the late 19th

century and engineers and designers depended mostly upon experience and the art of trial and error to the design of switching circuits. In 1938, C. E. Shannon's introduction of the idea of using Boolean algebra in the analysis and design of relay contact circuits started the rapid development of switching theory and digital systems.⁴

Thus engineers in digital or switching circuits acquired the proper tools to handle the analysis and synthesis of hardware; however, as the circuits and systems became more and more complex, the basic Boolean algebra manipulation became awkward and inefficient in handling the Boolean functions encountered in real systems. The first systematic attempt at more efficient synthesis of logic circuits was made by Harvard University's Computation Laboratory.⁵ The next important advancement was made by W. V. Quine in the early 1950's.⁶ In 1955, E. J. McCluskey generalized Quine's idea and also systematically developed the engineering notation and procedures for the synthesis of switching functions.

Another technique, the graphical method, also evolved. Veitch and Karnaugh, in slightly different manners independently derived a map from the Venn diagram which enables one to do the synthesis by visual inspection of the geometric patterns. Since then, there have been

numerous improvements in these two basic approaches. However, the covering problem in Quine-McCluskey's method and the limitation in the number of variables in the map method were always problems. The most recent papers offer some improvements and refinements in these areas. More detailed comparisons of various methods will be made in the next section.

In the late 1950's, R. P. Roth⁸ introduced the topological and simplex concept of treating Boolean functions, furthering the understanding of the inter-relationship between the Boolean expressions and Boolean functions. Roth also introduced the extraction algorithms which were the basis of the early computer algorithms developed at the University of Wisconsin⁹ and the computer programs at IBM.

As the technology in hardware improved, new devices and new circuits altered the cost value functions in minimizing hardware. Also, multi-valued logics, threshold logic, and sequential logic demands more generalized mathematical tools. With the coming of integrated circuits, the search for a minimum library of basic gates for all logic functions began. Thus a different approach to the synthesis of switching circuits developed. New technology like MOS also required a different class of algorithms. Different approaches to switching

theory also arose from the increase in size of the integrated circuits and systems. The decomposition of Boolean functions, the symmetry of Boolean functions, and the multiple-output problem also became of more practical than theoretical interest.¹⁰

With the increasingly realizable hardware devices for threshold logic and multi-valued logic, newer and more generalized mathematical tools are in order for the practical engineer and designer.

By the late 1960's, Tison's Generalization of Consensus Theory and the computer algorithms derived from it dominated the industry; however, in the 1970's the need to refine the computer algorithms directed the synthesis of switching functions into a different horizon, that of faster and more efficient computation. The algorithms developed in Auburn University by C.C. Carrol's group started to look for relationships between minterms of a Boolean function, so that the covering problem could be avoided in the algorithmized methods. This approach conceptually coincides with the notion of realizing that synthesis can be done with a Karnaugh Map by recognizing geometric patterns. McKinney's directed search algorithm and Arevalo's Four-tuples Methods are recent contributions in this direction. More detailed comparisons will be presented in the next section. A more

rational approach to the synthesis of switching functions is discussed in Section 3.

Footnotes: 1.1 Historical Review

¹Bell, E. T., The Development of Mathematics (New York: McGraw Hill Company, 1962), pp. 553-554

²Boole, George, An Investigation of the Law of Thought (London: 1854).

³Bell, pp. 553-577, has a very good account of the historical development of mathematical logics.

⁴Shannon, C. E., A Symbolic Analysis of Relay and Switching Circuits, Transcript AIEE, Vol. 57 (1938) pp. 713-723. Other people in Russia and Japan had similar concepts, notably, P. S. Ehrenfist of Russia in 1910, and Nakashima and Hanzawa of Japan in 1938. For more historical events, refer to G. C. Moisil, The Algebraic Theory of Switching Circuits (New York: Pergamon Press, 1969), and International Series of Monographs in Pure and Applied Mathematics, Vol. 41, pp. 11-20.

⁵See Book ref. (12) which stemmed from simplifying vacuum tube flip-flop circuits, which procedure can be applied to all switching functions.

⁶See ref. Q1, Q2, Q3

⁷See ref. K1, B7

⁸See ref. B6, S4, S5, S16, S17

⁹See ref. C10, S11, S15

¹⁰See ref. A6, B1, C10

¹¹See ref. T1

1.2 Comparison of Various Approaches

The basic concern of engineering design is to save cost and to improve performance. When we say cost, we have to consider both the software cost, hardware cost, as well as the difference in labor cost and the yield of constructing or manufacturing an item. When we say performance, we mean not only the performance of the specified parameters but also the reliability.

To increase reliability, we have to increase redundancy which is the opposite of minimizing the circuits; to improve the tolerance of faulty hardware we also have to sacrifice minimization of circuits. Therefore, in search of an optimal mathematical tool for the analysis and synthesis of switching circuits, mathematical completeness may not yield the most economical hardware, minimized hardware may not be worth the software cost of generating the design of hardware, and the most economical circuits may not have the desirable reliability.

(1) The Harvard Computation Laboratory Method¹

Listing all the minterms and all the possibilities of the basic functions which contain them is thorough and systematic; however, it is too bulky and cumber-

some to handle in any large size practical problem.

(2) The Standard Quine-McCluskey Method²

Using every minterm to build up a prime implicant table is also systematic and thorough, but suffers from the long process of enumeration and the very difficult problem of finding a minimum covering for the non-irredundant prime implicants.³ Many other tabular methods⁴ including R. P. Roth's cubical method and Tison's generalized consensus methods are also influenced by this method. In fact, this method presently represents the main stream of tabular or algorithmized methods.

(3) The Karnaugh and Veitch Map Method.

Using grid squares to represent minterms and the relationship between minterms has both a graphical and an algebraic advantage; however, the graphical aspect of the method so dominates the significance of the method, that its deeper algebraic implications are not fully developed. In this paper, some of the algebraic aspects of the method will be presented in a different manner.⁵

(4) Boolean Algebraic Methods.

R. T. Nelson's method⁶ of alternatively changing back and forth from products-of-sums to a-sum-of-

products during the algebraic cancellation or simplification of the Boolean expression has the advantage of not needing to convert Boolean expression into minterms. But this method has little usage because of the number of variables it can handle. Nevertheless, the characteristic of not requiring expansion of the Boolean expression into minterms is certainly a desirable feature. There are other algorithms which have this property?

(5) Topological and Simplex Cubical Method.

Urbano, Mueller, and R. P. Roth, in a series of papers, applied the topological similarity of a n-cube and the Boolean minterms, and advocated the cubical complex methods. Subsequently, Tison extended it to a generalized consensus method.⁹ At the University, of Wisconsin a series of computer programs were generated based on this method. This method contributes a better representation and notation for Boolean functions and a better visualization of the interrelationship between minterms. The n-cube notation and concept is adopted by most authors in the field¹⁰ and we owe our notation and terminology to this approach.

(6) Graphical Methods

There have been several instructive graphical methods proposed in the last twenty years, notably T. M. Booth's Vertex-Frame Method, M.E. Arthur's Geometric Mapping Method, and G.E. Marihugh and R.E. Anderson's H-Diagram Methods.¹¹ The H-Diagram, especially, has also proved to be fruitful in applications to other areas.¹² In this dissertation we offer a new graphical representation. This approach illustrates the concept of graphical depiction of Boolean functions.

(7) Mechanical Devices Method

A. Swoboda's contact grids^(S22), F. B. Hall's Binary Sieve^(H1), and A. K. Choudbury and M. S. Basu's Mechanical Charts represent attempts to use mechanical devices to achieve the synthesis of switching functions^(C5). This approach demonstrates the feasibility of constructing a slide rule like device to synthesize switching functions. Combined with the characteristic of algorithm orientation, we believe a small hand-held portable calculator is also highly feasible.

(8) Computer Algorithms

Originally, the computer programs all stemmed from computerizing the Quine-McCluskey Method. There was

limited success in programming Karnaugh's Map Method. The following group of computer programs have been reported in the literature:

- (a) IBM's R. P. Roth Group computer program based on R.P. Roth's Algebraic Topological approach.^(R5-R11)
- (b) C. C. Carrol, S. E. Jordan, W. A. Hornfeck, S. G. Shiva, and H. T. Nagle's Fast Algorithms at Auburn University.^(C13), ^(C14)
- (c) R. Slage, C. Chang, and R. C. Lee's sprouting algorithm from NIH.^(C13)
- (d) D. L. Dietmeyer and Y-H Su's Consensus Methods from Wisconsin University.^(D3), ^(S16), ^(S17)
- (e) M. H. McKinney's direct-search method algorithms at Texas A&M University.^(Mc5)
- (f) Z. Arevalo, J. G. Bredeson's Four-Tuple Algorithm.^(A8)
- (g) Harmonic Analysis to generate prime implicant.¹⁴
- (h) Bantee, and Mott have some older computer programming which is of historical interest.
- (i) Other Notable Methods
 H. R. Hwa^(H12), A. H. Sheinman^(S9), S. K. Das and N.S. Khabra^(D2), J. Mott^(M15), A. K. Cloudbury

and S. K. Das^{C12}, D. M. Y. Chang and T. H. Mott^{C2},
 Angelo R. Meo^{M6}, J. F. Gimpel^{G3}, N.N. Biswas^{B11}
 S. C. DeSarkar, A. K. Basu, A. K. Choudbury^{D7},
 A. Natapoff^{N1}, O. L. Ostapko and S. J. Hong^{O2}.

In this dissertation, computer programs are given to illustrate the new general algorithm. One general characteristic is to economize computer memory and computing time.

- (1) For detail procedure refer to ref. Book 12
- (2) Refer to reference Mc1, Mc2, Mc3, Mc4, or any standard logic design and switching theory text book.
- (3) Many good methods of finding the minimum set are theoretically interesting but not practical. References C6, C12, C15, C2, C3, P7, P8, G3.
- (4) See reference H12, S9, D2.
- (5) For original paper, refer to reference K1, V2, For examples and method, refer to a textbook, e.g. reference 4, 15, 36, 55, 39.
- (6) Refer to reference N4.
- (7) Refer to reference S13.
- (8) Refer to reference V1, R6, R7, R8, R9, R10.
- (9) Refer to reference T1
- (10) A good exposition of this notation and terminology can be found in reference Ref. Book 19.
- (11) Refer to reference B7, A5, M3
- (12) To Walch functions

(13) Refer to reference C5, S22, H1

(14) Refer to reference Book 41

1.3 Desireable Characteristics and Discussion

A comparison of the approaches in synthesizing Boolean functions and switching circuits indicates the following characteristics to be desirable:

1. Algorithm Oriented Characteristic

Because of the increasing use of computing machines and size, any approach to the synthesis of switching circuits has to be computerizable. F. Mileto and G. Putzolu^(M8),^(M9) did some initial work in estimating the merit and efficiency of computer algorithms for the minimization of Boolean functions based on the computing time and storage. However, because of the different criteria and concepts used in the different approaches, it was difficult to make an accurate and fair estimation. The size of the memory and the speed of calculation varies with different problems and settings. The skill of implementation into actual programs is also important. Although it is a very tedious job to evaluate the merits of the algorithms, it is mandatory to measure their effectiveness and efficiency. In this aspect, Roth's cube-array method is an improvement over Quine-McCluskey's, and the consensus algorithm is another step forward. Each of the newer algorithms claim some order of improvement

on time and storage saving.

2. Limitation of the Size Characteristic

Most of the approaches do not differ much in efficiency when the size of the problem is both small in number of variables and small in the number of elements of the on-set and don't cares. Therefore, the real measure of the effectiveness is on the size of problem that a specific approach can handle. The comparative merits have to be measured against each other with respect to a reference frame for size. The basic limitation of the map method is the difficulty of programming on the computer. Therefore it basically can be worked efficiently only with graphical methods, and the difficulty of depicting 8 variables limits its usefulness.

3. Visual Depictable Characteristic

The synthesis of switching circuits is not only to find a minimum circuit for a Boolean function, but also to observe the influence of changes of one part of the circuit on the remainder. It is not desirable to have to minimize the whole circuit every time we perturb the system. In this aspect, the map method offers the best opportunity. In fact, most of the

graphical methods try to gain this freedom and at the same time increase the number of usable variables. The H-Diagram is the most useful improvement since the original Karnaugh Map and can handle more variables. The fact that H-Diagrams have been applied to the Welch function indicates that this is important, and the wide use of the Karnaugh Map in industry for smaller problems shows the desirability of this feature.

4. Algebraical Structure Characteristic

Two-valued Boolean logic being the base for other logics, and logic design being the core of digital system and switching circuits, the algebraic structure of an approach would not only enrich the switching theory of Boolean functions, but could also be useful in other areas of switching theory.

Table 1.1 shows a comparison of the major approaches mentioned in Section 1.2 with these four characteristics. In Chapter 2, we shall develop a structure so that we can maximize all the advantages of different approaches.

CHARACTERISTICS		CHAR. 1	CHAR. 2	CHAR. 3	CHAR. 4
APPROACHES		(ALGO- RITHM)	(SIZE)	(VISUAL)	(STRUC- TURE)
AKERS	(A1)			X	X
AREVALO	(A3)	X	X		X
ARTHUR	(A5)			X	
BATNI	(B4)				X
BEATSON	(B5)				X
BOOTH	(B7)	X			X
BOWMAN	(B8)				X
BREUER	(B9,10)				X
BUBENIK	(B11)				X
CHANG	(C2 & 3)	X	X		
CHOUDHURY	(C4)			X	
CHOUDHURY	(C5)			X	
CHU	(C6)				X
COBHAM	(C7)				X
CURTIS	(C11)			X	X
DAS	(D1)			X	X
DAS	(D2)			X	X
DUNHAM	(D6)				
GILL	(G1)				X
GIMPEL	(G2,G3,G4)	X	X		X
GNAZALA	(G5)	X			X
GOODRICH	(G6)				X
GORDON	(G7)				X
HALL	(H1)				X
HARRIS	(H1)				X

Table 1 Comparision of Different Approaches

CHARACTERISTICS APPROACHES	CHAR.1 (ALGO- RITHM)	CHAR. 2 (SIZE)	CHAR. 3 (VISUAL)	CHAR.4 (STRUC- TURE)
POLANSKY (P4)				X
PYNE (P7,P8)				X
QUINE (Q1,Q2,Q3)				X
REUSH (R1)				X
RHYNE (R2)	X			X
ROBINSON (R3)				X
RODIN (R4)				X
ROTH (R5,R11)	X	X		X
ROY (R12)				X
SCHKOLNIK (S2)				X
SCHNEIDER (S4,S5)				X
SHANNON (S6,S7,S8)	X			X
SHEINMAN (S9)				X
SHEN (S10)				X
SHIRA (S12)				X
SLAGLE (S13)	X			X
STOFFERS (S15)				X
SURESCHANDER (S18)				X
SVOBODA (S20,21,22)			X	X
TISON (T1)	X			X
URBANO (U1)			X	X
VANDLING (V1)	X			X
VEITCH (V2)			X	X
VINK (V3)				X
ZISSOS (Z1)				X

Table 1 (continued)

CHARACTERISTICS APPROACHES		CHAR. 1 (ALGO- RITHM)	CHAR. 2 (SIZE)	CHAR. 3 (VISUAL)	CHAR. 4 (STRUC- TURE)
HARVARD	(H13)	X	X		X
HIRSCHORN	(H6)				X
HONG	(H7)	X			X
HOUSE	(H8)	X			X
HULME	(H10)			X	X
HWA	(H12)				X
IKRAM	(I1)				X
KARNAUGH	(k1)			X	X
KARP	(K3)	X			X
LAWLER	(L1,L2)	X			X
LIN	(L4)			X	X
LUCCIO	(L6)			X	X
MC CLUSKEY	(M3)		X		X
MARCOVITS	(M3)				X
MARIHUGH	(M3)				X
MEO	(M6)				X
MORREALE	(M10,M11)				X
MOTT	(M14,M15)			X	X
NATAPOFF	(N1)	X			X
NECULA	(N2)				X
NELSON	(N4,N5,N6)	X			X
OKADA	(O1)				X
OSTAPKO	(O2)				X
PARKHOMENKO	(P1)				X
PETRICK	(P2)				X

Table 1 (continued)

CHAPTER II

ALGEBRAIC STRUCTURE AND THEORY2.1 THE MINIMUM MAXIMAL-CUBE COVERING

The basic notation and terminology of the switching theory in this dissertation can be found in the textbook "Mathematical Theory of Switching and Automata" by S. T. Hu, a well executed attempt to organize all the known results in switching theory into a simple branch of pure mathematics with unified notation and terminology.⁽¹⁾ To develop our concepts properly, we found further structure had to be constructed to facilitate the progress in some areas. However, our main interest is in the hardware used in computer and digital systems rather than in the pure mathematics. The format chosen is not of a pure mathematical symbolic-logic approach, nor of the engineering journal's approach, which goes directly from results to examples placing definitions and proofs in the appendix. We chose to integrate the two. In this section, we will first discuss the properties of Karnaugh maps. Some of them have been utilized by other authors in recent algorithms independent of our development.⁽²⁾ In following sections, we shall construct an algebraic structure to facilitate the use of those properties.

1. Adjacency:

Quine-McCloskey's method demonstrates the importance of this property. Roth's and Tison's studies of the general consensus approach, expanded the use of this adjacency property,⁽³⁾ and most of the later authors also recognized the importance of this characteristic.⁽⁴⁾ Recently, H. C. Curtis's adjacency table method^{C11}, M. H. McKinney's directed search method,^{Mc8, Mc9} and Z. Arevalo's 4-tuple- Method^{A3} all made much use of the property. Our approach is to use the number of adjacent minterms to any particular minterm as one of its characteristic numbers.⁽⁵⁾ In the graphical methods, one of the advantages is that the adjacency can be determined by visual observation. In the analytical method, however, another kind of measure must be used to determine the adjacencies. The Hamming distance is the most logical metric and thus we shall adopt it.⁽⁶⁾ There are two parts of information pertaining to the adjacency. One is the number of cubes adjacent to a given cube; the second is the coordinates or directions of the adjacencies. In graphical methods, both are easily preserved except where there is a miscount due to human error. In our approach, we choose only to list the first part, the number of adjacent cubes, rather than identify each one of them individually, because we claim that the second part of

the information need only to be generated as needed. This saves labor of enumeration considerably.

2. Tops and Bottoms

Another important characteristic is the tops and bottoms of a cube. All the conventional approaches recognize this characteristic for any k-cube that may be chosen. However, their identification is still more or less a trial-and-error method as in the graphical approaches. We recognize the importance of this characteristic and will construct our algebraic structure in such a way that we can deduce the tops and bottoms analytically. Our investigation shows that the tops and bottoms not only need the information of adjacency, but also the two new concepts that we introduced in the investigation, the degree of coverability, and the order of coverability for a specific minterm point in the on-set. These will be discussed in detail in the section. These concepts have been implied in some of the recent results in a different form,^{S12} that is, to identify the largest cube covering for a group of on-set elements. According to Hu's terminology, this is the maximal cube covering for a specific vertex in the on-set. Consequently, the minimum number of these maximal cube-coverings consists of the minimum set of prime implicants for any Boolean function. Our approach

especially emphasizes this point of finding the minimum maximal cube covering directly. At the start, it seems that we must identify the maximal cube expended from any two vertices. The difficulty arises from the lack of analytic ways to prove that it is indeed the maximal cube covering. Our effort of constructing an algebraic structure is due to this demand. After we achieved this goal, we found the algebraic structure we built has more implication to switching theory than was thought. This shall be treated in later chapters.

3. Interrelationship Between Vertices

One of the most important advantages of a graphical approach is that the detailed relationship of one particular vertex with respect to all the rest of vertices of any on-set can be seen. In our investigation, it was found that from the concept of distance matrix and distance frequency matrix for the on-set, we were able to deduce most of these relationships easily. This new concept is a process similar to working with the map without drawing it. Additionally, we can employ most of the operation of the map method analytically.

4. Relationship Between Cubes

The prime implicant table covering problem stemmed from the intrinsic properties of the relationship between cubes.

The analytical approaches center their efforts in finding efficient transformation of Boolean expressions⁽⁷⁾ while graphical approaches focus their attention on a closer inspection of the graphical representation of the cubes.⁽⁸⁾ There seems lacking a fundamental treatment of this aspect of the cubes. The cyclic cubes, the less-than cubes has haunted most of the researchers in switching theory since the day it was born. Our first approach to it is the theory of cube-graphs and branching graphs which will be fully developed in Section 2.3. Our second attack on it is in Section 2.4, where the geometric meanings of the cube-covering in the metric space (Q^n, P) are introduced. In graphical approaches, the covering is drawn so that a minimum number of covering cubes can be found by trial-and-error, but no known theory has been developed to prove the validity of the existence theorem of a minimum set. We offer an attempt to prove the existence of such a minimum set with respect to a specific value function.

5. Beyond the Cube-Covering

We constructed the metric space (Q^n, d) out of the convenience for synthesising Boolean switching functions. It may well be worthwhile to look for more meaningful properties of Boolean switching functions direct from this discrete space. Further work in generalizing it into

other kinds of switching functions is feasible and will be discussed in Chapter V.

Applications with illustrative examples for the Branching graphs and the geometric approach will be given in Chapters III and IV.

Footnotes for 2.1

- (1) Hu's book represented the kind of effort by an eminent pure mathematician we referenced in the opening statements of this dissertation (refer to Chapter I, paragraph 1.1, page 1 of this dissertation). Reference Book 19.
- (2) Our initial results were submitted to IEEE Transaction on computer in December 1974. See repository on Computer magazine. (Reference H13, H14)
- (3) Adjacency is actually the basis for consensus.
- (4) Refer to Reference (H12)
- (5) We use column two of our characteristic matrix to represent this characteristic.
- (6) We use (Q^n, d) metric space. The definitions for (Q^n, d) and adjacency refer to Chapter II, para 2.2.
- (7) For example, the work of R. T. Nelson (Reference N4), M. J. Gnzala (Reference G5), F. Luccio (Reference L6), Petrick (Reference P2,P3), I.B. Byne and E. J. McCluskey (References P7, P8) and J. F. Gimpel (Reference G3).
- (8) Our deep conviction that the process of finding an optimal covering by inspection on Karnaugh Map must be proveable led us to devise a way of gaging the distance between two cubes and then to determine how close they really are.

2.2. Algebraic Structure of Minimum Maximal-Cube Covering Approach

Definition 1: Q^n is the cartesian product of Q sets
where $Q = \{0,1\}$. Q^n is called a n -cube.

Definition 2: (Q^n, d) is defined as a metric space with
metric function d being defined as follows:
If $x = (x_1, x_2, \dots, x_n)$, and $y = (y_1, y_2, \dots, y_n)$
 $x, y \in Q^n$, then: $d(x, y) = \sum_{i=1}^n |x_i - y_i|$ is
the distance between x and y . (The Hamming
distance = the number of positions in which
 x and y differ)

Definition 2 and the definitions and theorems derived from
it will suffice to render all the necessary information of
a Karnaugh map. Although the main application of this
approach is for the minimization of AND-OR (NAND) logic
circuits. It can be shown that the theory developed in
chapter 2 can be applied in many other areas.

Definition 3: Let $F = \{a_1, a_2, \dots, a_m\}$ be the on-set of
a switching function F . The $m \times m$ matrix $\{A = [a_{ij}]\}$ is defined
as $a_{ij} = d(a_i, a_j)$. The i -th row of A will be denoted by
 $d(a_i, F)$. The $m \times m$ matrix $B = [b_{ij}]$ the distance frequency
distribution matrix of F , that is, b_{ij} is equal to the
number of points in F that have distance j to a_i . We de-
note by $R(a_i)$ the i -th row of B and by $C(j)$ the j^{th} column of B .

Definition 4: $Q^* = \{1,0,*\}$, where * is the sign of don't care. If $N = \{1,2,\dots,n\}$ is the first N positive integer, then $\phi: N \rightarrow Q^*$. The coordinate function defines all the coordinates of a k -cube in Q^n . Furthermore $d \text{ in } (k) =$ the dimension of a k -cube = the cardinality of the subset $\phi^{-1}(*)$ of N . and $\text{Cod}(k) = n - \text{dim}(k) =$ the cardinality of the subset $\phi^{-1}(0)$ of N .²

Definition 5: A cube is covered by B cube if $A \subseteq B$.

Definition 6: A maximal cube is defined as the largest cube covering a subset.

When A cube in definition 5 reduces to a 0-cube, and B cube is taken from the union of on-set and don't care set, the definitions 5 and 6 are equivalent to the prime implicants in the Quine-McCluskey's minimization technique.

Example 2.2.1: $f(x) = \Sigma (0,2,3,4,5,6,7,13)$ with on-set $F = (0,2,3,4,5,6,7,13)$ then: $d(0,2) = 1$,
 $d(0,3) = 2$. etc.

Thus Matrix A and Matrix B for F are obtained in Figures 2.1 and 2.2 respectively. Figure 2.3 is the corresponding Karnaugh map.

We notice from the Karnaugh Map in Figure 2.3 that Vertex 5 has 3 adjacent Vertices, 4, 7 and 13, and they are on the left hand, right hand and above Vertex 5 respectively. Similarly, from Row Five and Column 2 of Matrix B we know Vertex 5 has 3 adjacent vertices, and from Row 5 of

	0	2	3	4	5	6	7	13
0	0	1	2	1	2	2	3	3
2	1	0	1	2	3	1	2	4
3	2	1	0	3	2	2	1	3
4	1	2	3	0	1	1	2	2
5	2	3	2	1	0	2	1	1
6	2	1	2	1	2	0	1	3
7	3	2	1	2	1	1	0	2
13	3	4	3	2	1	3	2	0

Figure 2.1 Matrix A

Distance Matrix for Example 2.2.1.

	P ₀	P ₁	P ₂	P ₃	P ₄	
1	2	3	2	0	0	
1	3	2	1	1	2	
1	2	3	2	0	3	
1	3	3	1	0	4	
1	3	3	1	0	5	
1	3	3	1	0	6	
1	3	3	1	0	7	
1	1	2	3	1	13	

Figure 2.2 Matrix B

Distance Matrix for Example 2.2.1.

8	9	11	10
12	13 	15	14
4 	5 	7 	6
0 	1	3 	2

FIGURE 2.3. Karnaugh Map for Example 2.2.1.

Matrix A we know exactly they are 4, 7 and 13, because Column 4, 7 and 13 of Row 5 have values of 1. It is not known whether they are on its left hand side or right hand side. To find this, we have to compute them from $5 = (0101)$, $4 = (0100)$, $7 = (0111)$ and $13 = (1101)$, 4 differs at the fourth coordinate, 7 at the 3rd, and 13 at the 1st respectively. the relationship of 5 to any other vertices can also be derived similarly.

Definition 7: $\langle a, b \rangle_k$ is defined as the k-cube having a and b as its diametrical extreme points. It is also called a k-cube spanned by a and b.

Theorem 1: Let $a \in F$ and $K \subseteq F$ be a k-cube covering a. then there exists a unique point $b \in K$ such that $d(a, b) = k$ and $K - \langle a, b \rangle_k = \{x \in F: d(a, x) + d(b, x) = k\}$

Proof: 1) Let b be the point in k obtained from a by complementing the varying coordinates of a. then $d(a, b) = k$ and no other point in k can be at a distance k from a!

2) a: clearly $\langle a, b \rangle_k = \{x \in F: d(a, x) + d(b, x) = k\}$

b: In computing the integer $d(a, x) + d(b, x)$ each of the k positions where a and b differ makes a contribution of 1 to this number. Therefore if $d(a, x) + d(b, x) = k$, x must agree with a and b in the constant coordinate positions. Thus $x \in \langle a, b \rangle_k$.

Corollary 1 $\# \{x \in F: d(z,x)+d(b,x)\} = 2^k$ (trivial)

Corollary 2: If a is a point in a k -cube in F , then the sum of the distance of x to any pair of diametrical extreme points is k .

Corollary 3: If a point $a \in F$ is covered by a k -cube in F , then there exists a k -cube in F with a as a diametrical extreme point.

Now we can examine example 2.2.1 again. Suppose the vertex 5 is chosen, $d(5,F) = (2,3,2,1,0,2,1,1)$. Since $\#(F) = 8 = 2^3$ and since at least one of the points (13), is not 3-cube coverable, we consider a 2-cube covering. Let's look for a 2-cube in F of the form $\langle 5.b \rangle_2$ that will cover 5. By theorem 1, we consider vertices that are 2-distance away. In this case, vertices 0,3. and 6 are all candidates for b . By corollary 1 of theorem 1, we need $2^2=4$

where $b=0,3, or 6. We find $d(5,F)+d(0,F) = (2,4,4,2,2,4,4,4)$. There are only three 2's, less than $2^k=2^3=4$.$

Therefore, we know $\langle 5.0 \rangle_2$ is not the 2-cube we are seeking nor is $\langle 5,3 \rangle_2$. However, $\langle 5,6 \rangle_2$ does satisfy the requirements. The only 2-cube existing is $\langle 5,6 \rangle_2$. Thus, the cubes can be found without the use of the Karnaugh maps. We trade off the immediate identifiable missing vertices of a cube to gain the advantage of algorithmizing the process.

Theorem 2: (Q^n, d) and Distance Matrix A are equivalent to the Karnaugh Map with on-set F.

Proof: 1) By definition 2. matrix A gives us all the information about adjacency between elements of an on-set F.

2) By theorem 1, we are able to identify all the cubes on the map.

3) The physical relation of the map can be deduced from matrix A, and the n-cube coordinates representation.

from 1) to 3), we conclude that we can generate the same kind of information from (Q^n, d) and matrix A as the K-map. Therefore, they are equivalent. Q.E.D.

The minimization of Boolean functions will be demonstrated. Instead of using visual inspection, we can now use deduction. The disadvantage of this approach is that m^2 units of storage is necessary to replace n squares of the map. Although matrix A is symmetrical, it still is necessary to use $m^2/2$ storage, where m is the number of minterms, and n is number of variables. However, in practice the memory is not the only factor, the calculation time is also important and matrix A does not have to be random accessable, it can be disc file .

In the next section, the theory for the synthesis of

switching functions will be developed. In Section 2.4, the geometric interpretation and additional theory for more efficient mechanization of the algorithms will be developed. Chapters III and IV will implement the algorithms for the results developed in Sections 2.3 and 2.4 respectively. Chapter V will extend theory into other areas.

Definition 8: If there exists a cube that covers a point in the on-set F of dimension j , this point is said to be j -coverable by a j -cube. If the point can be covered by an i number of j -cubes, then it is j -coverable of i -th order.

Definition 9: The subset F_j^i of F is the set of all the points in F that are j -coverable of order i .

Lemma 1: If a point is j -coverable, it is automatically $(j-1)$ - coverable. ($j \geq 1$)

Proof: If a is j -coverable, there exists a j -cube $\langle a, b \rangle_j$ which covers a . Let c be a point in $\langle a, b \rangle_j$ such that $d(a, c) = j-1$. Since $(c \in \langle a, b \rangle_j)$ C 's coordinates differ from a at $j-1$ places while the remaining place it differs from b , there $\langle a, c \rangle_{j-1} \subset \langle a, b \rangle_j$: $\langle a, c \rangle_{j-1}$ is a cube covering for a . Therefore, a is $j-1$ coverable. Q.E.D.

Theorem 3: Let $l_{ij} = \min_{k=0}^{k=j} \binom{b_{ik}}{k}$. Then l_{ij} is the maximal number of j -cubes that may exist to cover the point a_i of the on-set F which is represented by the i -th row in matrix A and matrix B . The matrix $L = [l_{ij}]$ represents the upper limit of the number of all maximal cubes that cover the points in the on-set F .

Proof: Since a cube is always symmetrical, the distance function is preserved under any rotation. Without loss of generality, the point I can be considered as the origin of the cube. If the cube is of dimension j , then by calculating the distance from every point in the cube to the origin and determining the number of points with equal distance, a distance frequency distribution pattern similar to the coefficient of binomial expansion (or Pascal triangle) is found. Thus, if there is a j -cube that covers the point I , then the number b_{ik} must be greater than or equal to the corresponding binomial coefficient $\binom{j}{k}$, and it is obvious that the most it can have is $\binom{b_{jk}}{k}$ of them.

Therefore, for the range, $k = 0$ to $k = j$, the smallest number represents the maximum number of j -cubes that may exist to cover a_i . Q.E.D.

Lemma 2: If a j -cube covering exists for point a_i then the number of points in F with $l_{ij} \geq 1$ is at least 2^j .

Proof: Since it takes 2^j points to form a cube, there must be 2^j points in the on-set F satisfying this condition.

Q.E.D.

Theorem 4: $\langle a_i, a_j \rangle_i$ is a cube of dimension j that covers a_i and a_j if and only if, the i -th row added to the j -th row in the distance matrix A produces exactly 2^j entries of j 's in the resultant row. Furthermore, these 2^j entries correspond to the points contained in $\langle a_i, a_j \rangle_j$ and a_i, a_j are the diametrical extremes.

Proof: 1. "If part":

Since there are exactly 2^j points which have the property of $d(a_i, a_k) + d(a_k, a_j) = j$, it suffices to prove that any point k having this property is contained in the j -cube, $\langle a_i, a_j \rangle_j$. For any point a_k having this property, a_k differs from a_i or a_j by $d(a_i, a_k)$ or $d(a_k, a_j)$ in the corresponding coordinates respectively. These places are all disjoint since any overlapping would imply that the distance between a_i and a_j is less than j , contradicting the fact that a_i and a_j are exactly the j -places in which a_i and a_j differ. Therefore, a_k differs from a_j at exactly those $d(a_i, a_k)$ places where a_i and a_j differ, and a_k differs from a_i at exactly those $d(a_k, a_j)$ places where the remaining a_i and a_j differ. Thus, a_k is one of the points belonging to the j -cube with a_i and a_j as the diametrical extremes and $d(a_i, a_k) + d(a_k, a_j) = j$. Since there are 2^j

of them, it must contain every point of the j -cube, $\langle a_i, a_j \rangle_j$.

2. "only if part":

If $\langle a_i, a_j \rangle_j$ is a j -cube covering a_i, a_j , then (a) $d(a_i, a_j) = j$, and (b) there are 2^j points in $\langle a_i, a_j \rangle_j$ including a_i, a_j . Now, if any point $a_k \in \langle a_i, a_j \rangle_j$, differs from a_i in x places, then $d(a_i, a_k) = x$; if a_k differs from a_j in y places, then $d(a_k, a_j) = y$. $a_k \in \langle a_i, a_j \rangle_j$, a_k can not differ from a_i or a_j other than where a_i and a_j differ. Otherwise a_k is not in $\langle a_i, a_j \rangle_j$. $\therefore x+y = j = d(a_i, a_j)$, $\therefore d(a_i, a_k) = j$ and there are 2^j points in $\langle I, J \rangle$. Thus, there must be exactly 2^j entries of j 's in the resultant row formed by adding the rows a_i and a_j .

Corollary: Given a_i, a_j and a_k of the on-set F : If $d(a_i, a_k) + d(a_k, a_j) = d(a_i, a_j) = j$, and if there exists any cube-covering that includes a_i, a_j and a_k , then the smallest cube is of dimension j .

Proof: If there is a covering, then $d(a_i, a_k) + d(a_k, a_j) = d(a_i, a_j) = j$ implies that the cube I, J covering K is of dimension j . Any cube smaller than dimension j cannot cover I and J . Therefore $\langle a_i, a_j \rangle_j$ of j dimension is the smallest cube that covers a_i, a_j and a_k .

Definition 10: A maximal cube which has to be included in any covering is essential and is called a core cube.

Lemma 3: Any maximal cube that can be included in the union of essential maximal cubes is redundant (trivial).

Definition 11: A maximal cube of a cover for a cube that can be included in the union of lower dimensioned cubes is a less than cube, thus redundant.

Definition 12: A point is said to be a j -coverable core point, if, and only if, it has one and only one maximal cube covering of dimension j , with all smaller cube covering included.

Theorem 5: If x is a j -coverable core point, then there are exactly j points in the on-set having distance of 1 from x .

Proof: Trivial (By definition 12)

The inverse of this theorem is not necessarily true.

Footnotes for 2.2

1. Reference book 18, 19.
2. Reference book 19, page 63-65

2.3 CUBE GRAPH AND BRANCHING GRAPH

To facilitate the analysis and synthesis, the concept of j -coverable is introduced so that the on-set F of a switching function can be divided into disjoint subsets of j -coverables. To further improve the ability to identify and select the cyclic cubes and the less-than cubes, the concept of a candidate for a core point, and the core point as well as the cube graph and a branching graph are used.

Definition 13: Two points are connected, if and only if they are covered by a common point.

Definition 14: Two cubes are connected, if and only if they both cover at least one common point.

Definition 15: A j -coverable branching graph is a linear graph representing all the connected j -coverable points.

The nodes in a j -coverable branching graph represent the diametrical extremes of the connecting cubes and the branching representing the connecting cubes.

The initial node can be any point in any set of connected j -coverable points. The next nodes are those points forming a connecting j -cube with the present node such that the diametrical extremes are the present node and the next node.

The terminal nodes are: (1) the nodes that have only one branch connected to them; (2) the nodes whose remaining

connecting branches going to the new nodes consist of cubes already represented by existing branches; (3) the nodes belonging to don't care or higher coverable sets.

Definition 16: A j -coverable cube-graph is a linear graph representing the j -cubes covering the connected j -coverable points.

Definition 17: Two j -coverable branching graphs are said to be equivalent if and only if they have the same j -coverable cube-graph.

Lemma 4: For every j -coverable branching-graph, there is one unique j -coverable cube-graph.

Proof: Let G be a j -coverable branching graph with C and C' being cube-graphs of G . If C is not coincident with C' , then there is at least one of the cubes which is different. Since every branch in a branching graph represents a distinct cube, the C must coincide with C' .

Lemma 5: A set of connected j -coverable points has a unique cube-graph.

Proof: Let F_j be a set of connected j -coverable points with C and C' being the cube graphs of F_j . Let cube $c \in C$ and $c' \notin C'$. There is then at least one point x in c that is not in any cube of C' . Since x is not in any of the cubes in C' , x is not connected to F_j . This contradicts the fact that x is a member of F_j .

Theorem 6: A set of connected j -coverable points must have 2^{j-1} different but equivalent branching graphs.

Proof : Since there are 2^{j-1} different pairs of diametrical extremes to a j -cube, there are 2^{j-1} different pairs of end nodes for each branch of a branching graph. When joined, there are 2^{j-1} distinct branching graphs. By Lemma 4 and Lemma 5, all are equivalent.

Theorem 7 : The minimum set of non-redundant maximal cubes is the union of minimum covering for the branching graph.

Proof: Since each branch of a branching graph represents a distinct maximal cube, the union of the minimum covering for each branching graph of a j -coverable subset ($j = 0$ to k) is the minimum set of non-redundant maximal cubes.

This theorem transforms the process of finding a minimum covering for a table of prime implicants into the finding of a minimum covering for each branching graph of a j -coverable subset of the on-set, taking the union over $j = 0$ to k , the highest dimension of the maximal cubes.

The advantages are: 1) No redundant maximal cubes are generated; 2) The on-set is partitioned into subsets of j -coverable connected points so that the covering problem can be treated separately for different subsets; and 3) The branching graphs are easier to work with.

The following theorems relate to the covering of branching graphs.

Theorem 8 : A minimum covering of a branching graph is the sum of the least number of branches needed to cover every circuit in the branching graph.

Proof: Since a branching graph represents the maximal cubes covering connected j -coverable points, and each loop in the branching graph represents a set of cyclic maximal cubes, therefore the sum of the least number of branches needed to cover each loop is the minimum covering of a branching graph. An open loop is a set of incomplete cyclic maximal cubes.

Theorem 9 : The least number of branches needed to cover a loop of n branches is $\frac{n}{2}$ if n is even or is $(\frac{n}{2} + 1)$ if n is odd.

Proof: Let the loop have n branches, then there are n nodes. Therefore, at least one half of these branches is needed to cover all those nodes because a loop is a closed plane curve and, within a loop, each node has exactly two branches connected to it. If n is odd, since a half-branch can not be used, an additional branch must be used.

Corollary 1: If the number of branches in a loop is even, the minimal covering of the loop is found by choosing every other branch.

Proof: Since in any loop, a cube represented by a branch can be divided into two parts: 1) the common part of

this cube and the cube represented by the preceding branch;
 2) The other is the common part of this cube and the cube represented by the following branch. Therefore if the preceding and the following branches are chosen, then the cube represented by this branch is covered and thus is redundant. Q.E.D.

Corollary 2: If the number of branches in a loop is odd, then except for the starting or finishing of the covering which has to be two consecutive branches, every other branch is the minimum covering.

Proof: Since the number of branches is odd, every other branch can not be chosen to cover a loop. There must be one part of a cube covered by two consecutive branches. Thereafter, it is the same as Corollary 1.

Theorem 10: If there are terminal nodes in a branching graph of the second kind, then these nodes have to be treated first for the elimination of the higher-dimensioned less-than cube.

Proof: If a terminal node is of the third kind, then it must be a higher-dimensioned cube coverable point. Since it is a terminal node of a branching graph, it has only one branch connected to it and that branch is necessary unless it is a less-than cube of some lower dimensioned cubes. Therefore, the higher dimensioned cube which also

covers this node may not be necessary, if the rest of the cube is covered. Thus, the set of all terminal nodes of the second kind must be tested for redundancy of the higher dimensioned maximal cubes in a higher coverable branching graph.

Theorem11: If a terminal node of a branching graph is of the first kind, then it is a core point, and the only branch connected to it represents an essential maximal cube.

Proof: Since a terminal node of the first kind by definition has one and only one maximal cube covering it and since no higher cube covers it, therefore it meets the definition of a core point. The only branch connected to the node represents the maximal cube mentioned above. Therefore, it is essential.

Theorem12: The more loops sharing a branch, the higher is the priority for choosing this branch.

Proof: The more loops sharing a branch, the more this branch can be used to cover different loops. Since this branch represents a partial covering for each loop, a smaller number of branches is needed. Q.E.D.

If there is more than one branch connected to a node, and no branch is shared by two or more loops, then the priority for choosing a branch for this node is inversely

proportional to the number of branches connected to the node.

Definition 18: If two branching graphs of different dimensions are linked such that the branches connected to the linking node in the higher dimensioned branching graph can be omitted in a minimum covering, the linkage is inseparable.

2.4 Geometric Structure of Maximal-Cube Covering

Since the use of Venn Diagrams, graphic depicting and geometric representation has been widely used in minimization of switching functions. In the author's earlier work, the geometric structure of the Karnaugh-Map and generalized finite discrete point representation of Boolean functions had been discussed. ^{H14,H15} The basic aim of using the geometric structure to represent functions is to seek visual recognition of the complex interrelationship. The analysis by the geometric structure of Maximal-cube covering approach is different than the graphical means.

(a) Diametric Extremes

The importance of using diametric extreme terminology is that cubes can be treated as if they were circles or spheres. Thus the drawing of circles tangent to each other would be a way of finding the minimum covering. Another advantage of this terminology is that the sense of interior points, boundary points, and exterior points can be introduced. In the case of Maximal-cubes of Boolean function, there are boundary points and exterior points, but no interior points. If x, y are a pair of diametric extremes of cube $\langle x, y \rangle$, then any point z is either on the boundary of the cube or outside the cube.

Another feature of the diametric extreme is that the

absolute origin of coordinates is not needed, because any pair of diametric extremes can be used as the reference.

(b) Awayness

The major advantage of graphical methods is that the distribution of cubes (a minterm point can be treated as 0-cube) can be visually observed. However, pattern recognition is always needed. The purpose of introducing the concept of awayness and aloofness in the maximal-cube covering approach is to simulate the kind of pattern recognition used in the map method.

To arrive at a numerical measure of the awayness or closeness of a cube to the existing cube covering, the concept of distance function is expanded.

Definition 9: A k-cube is re-defined as a pseudo-circle with diameter of k distance units.

Definition 20: The diameter of a pseudo-circle is the longest span of a pseudo-circle [equal to the dimension of a cube].

Definition 21: The two end-points of a diameter of a pseudo-circle is called the pair of diametric extreme points, or the diametric opposite points.

Definition 22: $d(x, A)$ is defined as the distance from vertex x to any pseudo-circle A (cube A). $d(x, A) = \text{Min} (d(x, a) : a \in A)$

Definition 23: $D(A)$ is defined as the distance units of diameters of pseudo-circle A (cube A) $D(A) = \text{Max} \{d(a_1, a_2) : a_1, a_2 \in A\}$

Definition 24: $d(A, B)$ is defined as the distance between two cubes A and B.

$$d(A, B) = \text{Min} \{d(a_1, a_2) : a_1 \in A, a_2 \in B\}$$

Theorem 13: The distance of a point x, the exterior point of a pseudo-circle (cube) to the pseudo-circle is the number of changes of coordinates between x and any point inside the pseudo-circle excluding those places where the points of the pseudo-circle vary their coordinates among themselves.

Proof: By Definition 18, $d(x, A) = \text{Min} \{d(s, a) : a \in A\}$

Let $L = m + n = d(s, a)$ for all $a \in A$.

m = the number of changes of coordinates between x and a at those places where the pseudo-circle (or cube) A varying coordinates among its points.

n = the number of changes of coordinates between x and a, where the coordinates for all points in the pseudo-circle (or cube) A are invariant.

Then n is constant for all points in A. Only m varies.

Therefore: $d(x, A) = \text{Min} \{d(x, a) : a \in A\}$

$$\begin{aligned} & m = k \\ & = \text{Min} (m + n) = n. \\ & m = 0 \end{aligned}$$

Q.E.D.

Corollary 1: The distance from interior points of pseudo-circle to the pseudo-circle is zero.

Proof: As in the proof of Theorem 13, when $n = 0$.

Definition 25: The smallest cube or pseudo-circle having a point of a cube A (or pseudo-circle A) and any point x is defined as the connecting cube (or pseudo-circle) between A and x.

Definition 26: The smallest pseudo-circle (or cube) having a pair of diametric extreme points where one is from pseudo-circle A (or cube A) and the other is from pseudo-circle B (or cube B) is called the connecting pseudo-circle (or cube) between pseudo-circle A (or cube A) and pseudo-circle B (or cube B).

Theorem 14: $d(x, A)$ = the dimension of the connecting cube between x and A,

$d(A, B)$ = The dimension of the connecting cube between cube A and cube B. (Since pseudo-circle = cube, we will use cube from now on.)

Definition 27: The foot of a connecting cube between a cube A and a point χ is a point in A such that $C = \langle \chi, a \rangle$

Theorem 15: The distance between a point to a cube is equal to the distance between the point to the foot of the connecting cube.

Proof: Let a be the foot of the connecting cube for point

x to cube A. By Definition 27, $d(a, x)$ = dimension of connecting cube. By Definition 25, $d(a, x) = \text{Min } \{d(x, y) : y \in A\} = d(x, A)$ Q.E.D.

Theorem 16: If a and b are a pair of diametric extremes, then all the places of the coordinates where the points of the cube vary their coordinates among themselves are complementary.

Proof: Let $\langle a, b \rangle$ be a cube of k dimension, then there are k places among the coordinates where the points of the cube vary among themselves. Since $d(a, b) = k$ and only k place can vary, therefore a and b have to be complementary at these places. Q.E.D.

Lemma 6: If a_1, a_2 are diametric extremes of a k-cube, then $d(a_1, a_2) = D(A) = k = \text{dimension of the cube} = \text{the diameter of the sphere or circle. (Trivial)}$

Theorem 17: If a_1, a_2 are diametric extremes of a k-cube and $x \in Q^n$, then

$$(1) \quad d(a_1, x) + d(a_2, x) = d(a_1, a_2) = k \rightarrow x \in \langle a_1, a_2 \rangle_k.$$

$$(2) \quad d(a_1, x) + d(a_2, x) < d(a_1, a_2) \rightarrow x \text{ is outside of } \langle a_1, a_2 \rangle_k$$

$$(3) \quad d(a_1, x) + d(a_2, x) < d(a_1, a_2) \text{ does not exist.}$$

Proof: (d, Q^n) is a metric space, implies

$$d(a_1, x) + d(a_2, x) \geq d(a_1, a_2).$$

by Theorem 1, Corollary 1 (1) holds.

(2) implies that point x has some coordinates differing from a_1 , or a_2 which are not among the places a_1 and a_2 are different. Therefore x is not included in $\langle a_1, a_2 \rangle$. Thus x is outside of $\langle a_1, a_2 \rangle$. Q.E.D.

Theorem 18: If $A = \langle a_1, a_2 \rangle_k$, then for any $\chi \in Q^n$

$$d(a_1, x) + d(a_2, x) - k = 2d$$

where $d = d(x, A) =$ the distance from x to A .

Proof: If $\chi = \phi$ (Trivial)

$$\text{If } \chi \neq \phi, d(a_1, x) + d(a_2, X) \geq d(a_1, a_2) = k$$

(by definition of metric space)

1) If $\chi \in \langle a_1, a_2 \rangle_k$ then $d(\chi, A) = 0$

By Theorem 1, Corollary 1;

$$d(a_1, x) + d(a_2, x) = k$$

$$d(a_1, x) + d(a_2, x) - k = 0 = d(X, A)$$

2) If $\chi \in \langle a_1, a_2 \rangle_k$, then

By Theorem 13;

$$d(a_1, x) = M_1 + d(X, A) \dots (1)$$

$$d(a_2, x) = M_2 + d(X, A) \dots (2)$$

Where $d(x, A) =$ The number of coordinate positions

where x differs from A , among the subset of coordinate positions where A does not vary its coordinates among its elements.

M_1 = The number of coordinate position x differ from a_1 , among the subset of coordinate positions where A varying its coordinates.

M_2 = The number of coordinate position x differ from a_2 , among the subset of coordinate positions where A varying its coordinates.

Since a_1 and a_2 are complementary at the varying coordinate positions, M_1 and M_2 must also be complementary.

Thus $M_1 + M_2 = k$.

(1)+(2) resulted in

$$d(a_1, x) + d(a_2, x) = M_1 + M_2 + 2d(x, A)$$

$$d(a_1, x) + d(a_2, x) - k = 2d(x, A)$$

Q.E.D.

Since $\langle a_1, a_2 \rangle$ is an arbitrary pair of diametrical extremes, and x is any point in Q^n , therefore, this is true for any pair of diametrical extremes.

Thus, $d(a_1, x) + d(a_2, x) - k$ is an invariant.

Corollary 1: For all $x \in Q^n$, $a \in A$.

$$d(x, a) = M + d(x, A)$$

Where M is the number of positions x differ from a , among those coordinate positions A is varying. and $d(x, A)$ is the distance of x to A .

Proof: See main theorem.

Corollary 2: For all $x \in Q^n$, $a_1 \in A$, $a_2 \in A$. if $\{d(a_1, x) - d(x, A)\} + \{d(a_2, x) - d(x, A)\} = k$. a_1, a_2 form a pair

of diametrical extreme points of A.

Proof: See theorem 18 (gives an invariant which is going to be useful in minimization. It will be called awayness.)

Definition 28: Awayness is defined as the distance of a point to cube, $W_{ab}(s) = d(a,x) + d(b,x) - k = 2d$
 Definition 28 can also be extended to the distance between two cubes A, B.

Definition 29: Awayness of two cubes A, B is defined as

$$W_A(B) = d(a_1, B) + d(a_2, B) - k_A + \#(R)$$

where k_A is the dimension of cube A, and R is the subset of coordinate positions when both A, and B are varying its coordinates. We will define $\#(R)$ as the degree of skewness.

Definition 30: The skewness is defined as the number of coordinate positions both cube A and cube B are varying. We will denote it by $SK_A(B) = \#(R) = SK_B(A) = SK(AB)$

Theorem 19: Awayness is independent of diametric extremes

Proof: (1) If one of the cube is a 0-cube then theorem 19 reduces to theorem 18.

(2) If cube A and cube B are not 0-cubes, then:

$$\begin{aligned} W_A(B) &= d(a_1, B) + d(a_2, B) - K_A = 2d - \#(R) \\ &= d(b_1, A) + d(b_2, A) - K_B = W_B(A) \end{aligned}$$

where $\langle a_1, a_2 \rangle_{KA=A}$, $\langle b_1, b_2 \rangle_{KB=B}$.

d is the distance between A and B. $\#(R)$ is the number of common parts in the varying coordinates of cubes A, B.

There are four different categories of coordinate positions for cubes A, B denoted as follows.

p = those positions only cube A varying its coordinates;

Q = those positions only cube B varying its coordinates;

R = both A and B varying its coordinates;

S = neither A or B varying its coordinates

It is obvious $\#(P+R) = K_A$

$$\#(Q+R) = K_B$$

$\#(P+Q+R+S) = n$, the dimension of Q^n .

Let $d(x, y) = \{x \in A, y \in B\}$.

Then $d(x, y) = M_p + M_q + m_r + M_s$

where M_p = The number of the positions x and y differ in values among P .

M_q = The number of the positions x and y differ in values among Q .

M_r = The number of the positions x and y differ in values among R .

M_s = The number of the positions x and y differ in values among S .

We observe the following facts.

$M_S = \text{Constant}$, because both cubes A, B are not varying.

M_P varies from 0 to $K_A - \#(R)$

M_Q varies from 0 to $K_B - \#(R)$

M_R varies from 0 to $\#(R)$

Thus $d(a,b) = d(A,B) = \min d(x,y)$

$$= M_S$$

That is, the closest points a, b in A and B only differ in those places neither A and B are varying. If we complement a , and b in A and B , we obtain c in A , and d in B . and by definition $\langle c,a \rangle_{K_A} = A$, $\langle b,d \rangle_{K_B} = B$.

Furthermore $d(c,d) = K_A + K_B - 2\#(R) + M_S$, because in the process of complementing a , and b , we change only those P and R values for a , and Q and R values for b and none in S . Since we are simultaneously complementing R for a and b , there is no difference in the resultant c , and d . Therefore none contribute to distance number; while in S subset; c and d differ the same as a and b , thus contribute M_S . However in the P and Q subsets, c and d differ $\#P$ and $\#Q$ as against no difference for a and b . Totalizing all the differences,

$$\begin{aligned} d(c,d) &= \#(P) + \#(Q) + M_S \quad \dots \quad d(c,d) = K_A - \#(R) + K_B - \#(R) + M_S \\ &= K_A + K_B - 2\#(R) + M_S \end{aligned}$$

(These are the most distance points among A and B .)

i.e. $d(c,d) = \text{Max} \{d(x,y) \mid x \in A, y \in B\}$

$$d(a,b) = M_S = \text{Min} d(x,y) \dots \dots (1)$$

$$d(c,d) = K_A + K_B - 2\#(R) + M_S \dots \dots (2)$$

for any pair of diametric extremes a_1, a_2 of A.

$$\text{Thus, } d(a_1, b) + d(a_2, b) = K_A + 2d(b, B) \dots \dots \dots (3)$$

$$d(a_1, d) + d(a_2, d) = K_A + 2d(d, B) \dots \dots \dots (4)$$

$$\text{Since } 2d(d_1, B) = d(c_1, b) + d(a_1, b) - K_A \dots \dots \dots (5)$$

$$2d(d_1, B) = d(c_1, d) + d(a_1, d) - K_A \dots \dots \dots (6)$$

and since c , and d were obtained from a , and b ,

$$d(c, b) = \#(P) + \#(R) + M_S = K_A + M_S \dots \dots \dots (7)$$

$$d(d, a) = \#(Q) + \#(R) + M_S = K_B + M_S \dots \dots \dots (8)$$

Substituting (1), (2), (7), (8) into (5), (6)

$$2d(b, B) = K_A + M_S - K_A = 2M_S \dots \dots \dots (9)$$

$$\begin{aligned} 2d(d, B) &= K_A + K_B - 2\#(R) + M_S + K_B + M_S - K_A \\ &= 2[K_B + M_S - \#(R)] \dots \dots \dots (10) \end{aligned}$$

Substituting (9) and (10) into (3), and (4).

$$d(a_1, b) + d(a_2, b) = K_A + 2M_S \dots \dots \dots (11)$$

$$d(a_1, d) + d(a_2, d) = K_A + 2[K_B + M_S - \#(R)] \dots \dots \dots (12)$$

add (11) (12) and rearrange the left hand side to obtain.

$$\begin{aligned} &d(a_1, b) + d(a_1, d) + d(a_2, b) + d(a_2, d) \\ &= K_B + 2d(a_1, B) + K_B + 2d(a_2, B) \\ &= 2[K_B + d(a_1, B) + d(a_2, B)] \end{aligned}$$

While the right hand side is

$$\begin{aligned} &2[K_A + K_B + 2M_S - \#(R)] \\ &\therefore K_B + d(a_1, B) + d(a_2, B) \\ &= K_A + K_B + 2M_S - \#(R) \\ &\therefore d(a_1, B) + d(a_2, B) - K_A \end{aligned}$$

$$= 2M_s - \#(R) = 2d(a_1b) - \#(R)$$

$$= 2d - \#(R)$$

$$\text{Similarly } W_B(A) = 2d - \#(R).$$

Q.E.D.

Remark: If one of the cubes reduces to a point (i.e., 0-cube), then $\#(R) = 0$. Otherwise, the 0-cube becomes a higher dimension cubes. And from theorem 19, a new parameter appears, that is the $\#(R)$ which gives the relative direction of two cubes, which can be observed on the map. Theorem 18, 19 are analogous to plane geometry theorems shown in Figure 2.4.1, and Figure 2.4.2. Other examples and demonstrative figures are given at the end of this section and in Chapters IV and V.

For Figure 2.4 We have:

$$\begin{aligned} \overline{Bx} \cdot \overline{Ax} &= \overline{Dx} \cdot \overline{Cx} \\ &= (D+d) \quad (d) \\ &= Dd + d^2 = \text{invariant} \\ &\text{for any } \overline{DCx} \text{ line} \\ &\text{through } x \text{ which} \\ &\text{intersects the} \\ &\text{circle} \end{aligned}$$

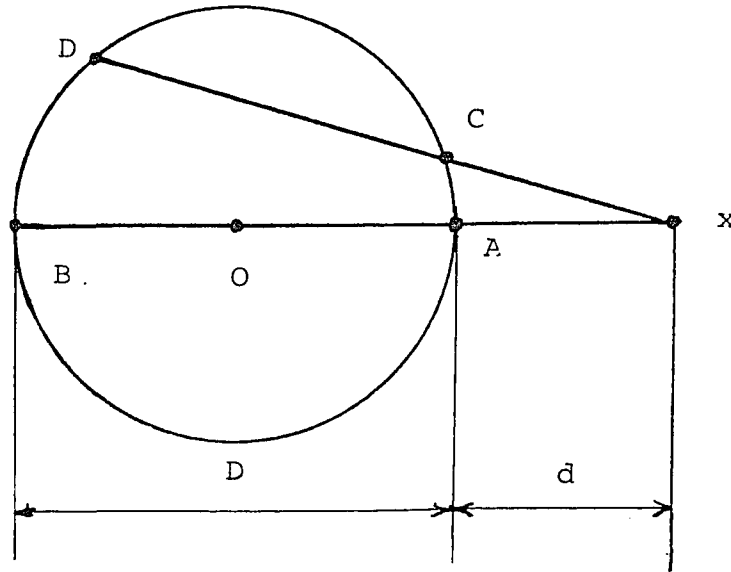


Figure 2.4 The Analogy of the Awayness of a Point to a Cube in Plane Geometry

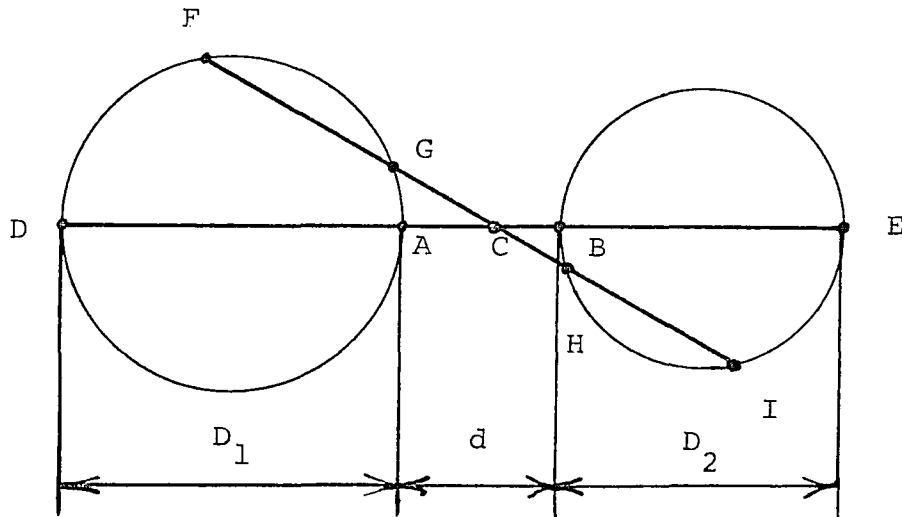


Figure 2.5 The Analogy of the Awayness of a Cube to a Cube in Plane Geometry

For Figure 2.5.

If, Let $AC = BC = \frac{1}{2} d$

$$\text{Then } GC.FC = \frac{1}{2}D_1d + \frac{1}{4}d^2 \dots \dots \dots (1)$$

$$IC.HC = \frac{1}{2}D_2d + \frac{1}{4}d^2 \dots \dots \dots (2)$$

$$Gc.Fc + Ic.Hc = \frac{d}{2} (D_1 + D_2) + \frac{d^2}{2}$$

$$GC.FC.IC.HC = (D_1 + \frac{d}{2}) (D_2 + \frac{d}{2}) \frac{d^2}{4}$$

For minimizaing minterms of a switching function we have the following theorems.

Theorem 20: If a point is j coverable and having awaysness of W_{ab} to cube $\langle a,b \rangle$ of D dimension, then the number of the common points between the cube covering of the point and the original cube is zero if W_{ab} is greater than j . If W_{ab} is smaller than j , then there are at most $2^{j-W_{ab}}$ points in common between the old cube and the new one.

Proof: If W_{ab} is greater than k , then x is aloof from the old cube; thus no common points exist. If W_{ab} is smaller than j , then there may be at most $2^{j-W_{ab}}$ common points.

It is clear that W_{ab} is a very important factor in determining the priority of the Multiple-coverable points.

Theorem 21: $j-W_{ab}$ is always smaller than D , the dimension of the chosen cube.

Proof: If $j-W_{ab}$ is larger than D , it is possible to find a cube that covers x and the whole cube $\langle a,b \rangle$ of dimension

j. But $\langle a, b \rangle$ is the largest cube to cover a and b. If the j cube does not intersect $\langle a, b \rangle$ then they are separated and there is no relation. This can also be extended to higher ordered cubes. Let us define the maximal distance of two cubes as the envelope of two cubes.

Definition 32: The envelope of two cubes A and B, $[e(A, B)]$ is defined as $e(A, B) = K_A + K_B - \#(R) + d(A, B)$

Theorem 22: The smallest cube covering both A and B has the dimension of $e(A, B) + \#(R)$.

Proof: From the proof of Theorem 19.

$$\begin{aligned} e(A, B) + \#(R) &= K_A + K_B - 2\#(R) + d(A, B) + \#(R) \\ &= \{(K_A - \#(R))\} + \{K_B - \#(R)\} + d(A, B) + \#(R) \end{aligned}$$

where $K_A - \#(R)$ = the number of coordinate positions only cube A is varying.

$K_B - \#(R)$ = the number of coordinate positions only cube B is varying

$\#(R)$ = The number of coordinate positions both A.B are varying

$d(A, B)$ = The number of coordinate positions neither A nor B are varying but where A and B differ.

Q.E.D.

Theorem 23: If cube A and cube B are of the same dimension R, and $d(A, B) = e(A, B)$ then all varying coordinates coincide.

Proof: $e(A, B) = K + K - 2\#(R) + d(A, B)$ (By Theorem 22)

$$\begin{aligned} \therefore 2K - 2\#(R) &= e(A, B) - d(A, B) \\ &= 0 \end{aligned}$$

$$\therefore K = \#(R)$$

Q.E.D.

Corollary: Of $e(A,B) = d(A,B) = 0$, $A=B$.

Theorem 24: If cube A and cube B have the same dimension, $d(A,B) = 1$, $e(A,B) = K+1$ then A and B form a $(k+1)$ -cube.

Proof:

$$\begin{aligned} e(A,B) &= K+K - 2\#(R) + d(A,B) \\ &= 2K - 2\#(R) + d(A,B) \end{aligned}$$

By Theorem 22, the smallest cube covering A and B has a dimension of

$$e(A,B) + \#(R) = 2K - 2\#(R) + 1 + \#(R) = 2k - \#(R) + 1$$

$$\#(R) = 2K + 1 - e(A,B) = 2K + 1 - K - 1$$

$$= K.$$

$$\therefore e(A,B) + \#(R) = 2K - \#(R) + 1 = 2K - K + 1$$

$$= K + 1.$$

Q.E.D.

Theorem 25: Let $f(A,B) = \max \{d(x,y) : x \in A, y \in B\}$.

Then $(A,B) = e(a,b)$ iff $\#(R) = 0$

1) if part,

Proof: $f(A,B) = \text{Max } (d(x,y) : x \in A, y \in B).$

$$= \text{Max } \{M_p + M_q + M_r + M_s.\}$$

Maximum of $m_p = K_A - \#(R)$ Maximum of $M_q = K_B - \#(R)$,

$M_r = \#(R)$, and $M_s = d(A,B)$

$$\begin{aligned} \text{Thus } f(A,B) &= K_A - \#(R) + K_B - \#(R) + \#(R) + d(A,B) \\ &= K_A + K_B - 2\#(R) + d(A,B) + \#(R) \\ &= e(A,B) + \#(R) \end{aligned}$$

$$\therefore \#(R) = 0$$

2) only if part.

$$\#(R) = 0. \text{ implies } M_p = K_A, M_q = K_B \quad M_r = 0$$

$$M_s = d(A, B)$$

$$\therefore f(a, b) = K_A + K_B + d(A, B).$$

$$e(A, B) = K_A + K_B - 0 + d(A, B) = K_A + K_B + d(A, B)$$

$$\therefore f(a, b) = e(A, B)$$

Corollary: $f(a, b) = e(A, B) + SK(A, B)$

Theorem 25: Given cubes A, B, and $d(A, B) = 0$, then

$$\text{the number of common points} = 2^{\#(R)} = 2^{SK(A, B)}$$

Proof: (1) If $d(A, B) = 0$, there are no coordinate positions

difference outside of those position where either A,

or B are varying. Thus the intersection of cube A

and cube B has a dimension of $\#(R) = SK(AB)$. and their

varying coordinate positions are exactly those $\#(R)$

positions where both A and B are varying.

Theorem 26: Given cubes A and B. and $d(A, B) = 1$.

$$e(A, B) = 1. \text{ Then } SK(AB) = \frac{K_A + K_B}{2}, f(a, b) = \frac{K_A + K_B + 2}{2}$$

Proof. $e(A, B) = K_A + K_B - 2SK(AB) + d(A, B)$

$$\therefore K_A + K_B - 2SK(A, B) = 0$$

$$\therefore K_A + K_B = 2SK$$

$$SK(AB) = \frac{K_A + K_B}{2}$$

$$f(a, b) = e(a, b) + SK(A, B) = 1 + \frac{K_A + K_B}{2} = \frac{K_A + K_B + 2}{2}$$

Q.E.D.

(c) Aloofness

Another measure of the merit in finding the minimum covering is how evenly can the largest cubes be arranged. Ideally, we would like to find a maximal-cube adjacent to the existing cube covering. For this, we shall introduce the concept of aloofness, which turns out to be a way of gauging how evenly we can arrange the maximal cubes. In order to define aloofness, we shall generalize the definition of adjacency among 0-cubes of any dimensions.

Definition 33: The "aloofness" is defined as the frequency that a point's awayness exceeds the highest dimensions, that is, the number of times when $W_{ab}(x) - d_x \geq 0$, d_x being the highest dimension of cubes that x can have for maximal-cube covering.

Definition 34: $W_A(B) - d_x \geq 0$ is the aloofness for cube B with regard to cube A.

It is evident that the selection of the next point should be based on smaller awayness and larger aloofness so that we can evenly spread out as will be demonstrated in chapter IV.

(d) Adjacency

This property is commonly used in literature, and refers to finding as many adjacent sub-cubes to the existing cube covering as possible.

Definition 35: Adjacency is defined between two cubes only if cubes A and B have a distance of 1, i.e. $d(A,B) = 1$. It is obvious that we want to choose the maximal cubes that are adjacent. The reason the 'map method is so powerful is that one can observe this by visual investigation. In the method given here, it is analytical, not a trial and error method. To achieve the indicated results, a merit value is assigned to each uncovered point or cube.

Definition 36: The degree of adjacency is defined as the number of distinct sub-cubes that are adjacent to the existing maximal cube covering.

e) Probability Overlapping.

In the process of choosing maximal cube covering, the existing maximal cube covering should be overlapped as little as possible.

Definition 37: The probability of sharing common parts of a new cube with the existing cube covering is: $(share) = D$, if the highest coverable maximal cube has a dimension of $J \geq D + W_{ab}$ or $= J - W_{ab}$, if $J \leq D + W_{ab}$.

If the awayness and the aloofness are the same, then the next criteria is the relative probability of sharing. The last check is to compare the real over-lap of common parts.

Definition 38: If $d(A,B) > 1$ for all $B \in F$, A is isolated.

Definition 39: Let $*$ = $1/2$ to be the don't care value
 $Q^* = \{0, 1, *\}$ for the value of the
 coordinates of the center of pseudo-
 circles.

f) Other related structure.

Lemma 7: If $d(A,B) = 0$, then A, B are coincident

Lemma 8: If $d(A,B) > 1$, then A, B are separated

CHAPTER III

BRANCHING GRAPH APPLICATION AND ILLUSTRATIVE EXAMPLES.

3.1 THE ALGORITHM

The minimization of two-level sum-of-products logic can be summarized as follows:

- Step 1: The distance Matrix A is constructed;
- Step 2: From Matrix A, the distance-frequency distribution Matrix B is obtained;
- Step 3: From Matrix A and Matrix B, the limit matrix L is obtained;
- Step 4: From the theorems and matrices A, B, and L and with possible core points as candidates, the essential maximal cubes (i.e, essential prime implicants) $\bigcup_{j=0}^P M_j^1$ are selected by starting with the highest dimensioned j-cubes down to the lowest, i.e., finding $\bigcup_{j=0}^P M_j^1(f) \supseteq \bigcup_{j=0}^P F_j^1$ (All of the essential maximal cubes have been found);
- Step 5: The core points are used as the initial nodes to find the branching graph for each set of connected j-coverable points, where $j=k, k-1, k-2, \dots, 3, 2, 1$ and k is the highest dimension of the maximal cube covering;
- Step 6: If there are terminal points of higher cover-

able points or branches containing higher coverable points, then they are compared with the cube graph of higher dimension, to see if any less-than cubes exist;

Step 7: A minimum covering of the branching graph, which the minimum non-redundant set of maximal cube covering, i.e., $\bigcup_{i=1}^{p,q} MMF_j^i$ is found.

3.2 Demonstrative Electives:

Example 1: To illustrate that a non-essential prime implicant will not be generated if it is already included in the set of essential cubes:

$$f(x) = \Sigma (0, 2, 3, 4, 5, 6, 7, 13).$$

Step 1: Construct Matrix A (Fig. 1). (Matrix A is symmetrical).

Step 2: Construct Matrix B (Fig. 2) from Matrix A.

Step 3: Construct Matrix L (Fig. 3) from Matrix B .

- 1) Since only four points (4, 5, 6, and 7) satisfy the necessary conditions for a 3-cube covering, no 3-cube covering exists. Multiple 2-covering may exist for these points.
- 2) Only points 0 and 3 are possible 2-coverable core points.
- 3) Point 13 is a possible 1-coverable core point.

- Step 4: Since 0 and 3 are of same order and degree, point 0 is arbitrarily chosen first. Only 2-cubes are being dealt with in this case, and therefore any distance over 3 is immaterial. All distances ≥ 3 or over are replaced by X. The table in Fig. 4 shows the distances for 0, 3, 5, and 6 taken from Matrix A. Since 3, 5, and 6 are points having a distance 2 to 0, they are possible candidates for diametrical extremes. Theorem 2 shows that only the $\langle 0,6 \rangle$ cube exists. ($\langle 0,6 \rangle = 0,2,4,6$). Now L is updated to L_1 , by changing the row representing 0 in L to $(-, -, e)$, indicating that point $0 \in F_2^1$, and 2, 4, and 6 belong to the don't care conditions (Fig. 5).
- Step 5: Continue to point 3, which is another candidate for a 2-coverable core point and find the $\langle 3,6 \rangle$ which is an essential 2-cube. Update L_1 to L_2 (Figs. 6 and 7).
- Step 6: Having exhausted the 2-coverable core points, the possible 1-coverable core points are considered. Thus, point 13 is to be covered by $\langle 13, 5 \rangle$. Update L_2 to L_3 (Figs. 8 and 9). Since all points are covered, the process is complete.

	0	2	3	4	5	6	7	13
0	0	1	2	1	2	2	3	3
2	1	0	1	2	3	1	2	4
3	2	1	0	3	2	2	1	3
4	1	2	3	0	1	1	2	2
5	2	3	2	1	0	2	1	1
6	2	1	2	1	2	0	1	3
7	3	2	1	2	1	1	0	2
13	3	4	3	2	1	3	2	0

Figure 3.1 Distance Matrix A of Example 1

	ρ_0	ρ_1	ρ_2	ρ_3	ρ_4	
	1	2	3	2	0	0
	1	3	2	1	1	2
	1	2	3	2	0	3
	1	3	3	1	0	4
	1	3	3	1	0	5
	1	3	3	1	0	6
	1	3	3	1	0	7
	1	1	2	3	1	13

Figure 3.2 Distance Frequency Matrix B of Example 1

	c_0	c_1	c_2	c_3	c_4
0	-	-	c	0	0
2	-	-	-	0	0
3	-	-	c	0	0
4	-	-	-	c	0
5	-	-	-	c	0
6	-	-	-	c	0
7	-	-	-	c	0
13	-	c	0	0	0

Figure 3.3 Matrix L of Example 1

	0	2	3	4	5	6	7	13
0	0	1	2	1	2	2	X	X
3	2	1	0	X	2	2	1	X
5	2	X	2	1	0	2	1	1
6	2	1	2	1	2	0	1	X

									N_2	C_2
0 + 3	2	2	2	X	X	X	X	X	3	NO
0 + 5	2	X	X	2	2	X	X	X	3	NO
0 + 6	2	2	X	2	X	2	X	X	4	YES

Figure 3.4 Table for Finding 2-Cubes Covering Point 0 in Example 1

	c_0	c_1	c_2
0	-	-	e
2	-	-	d
3	-	-	c
4	-	-	d
5	-	-	-
6	-	-	d
7	-	-	-
13	-	c	0

Figure 3.5 Matrix L_1 of Example 1

	0	2	3	4	5	6	7	13		
3	2	1	0	X	2	2	1	X		
5	2	X	2	1	0	2	1	1		
6	2	1	2	1	2	0	1	X		

									N_2	C_2
3 + 5	X	X	2	X	2	X	2	X	3	NO
3 + 6	X	2	2	X	X	2	2	X	4	YES

Figure 3.6 Table for Finding 2-Cubes Covering Point 3 in Example 1

	c_0	c_1	c_2
0	-	-	e
2	-	-	d
3	-	-	e
4	-	-	d
5	-	-	-
6	-	-	d
7	-	-	d
13	-	c	0

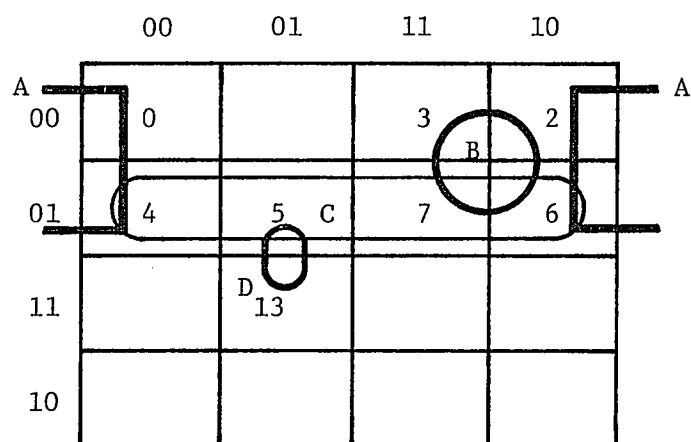
Figure 3.7 Matrix L_2 of Example 1

	0	2	3	4	5	6	7	13		
13	X	X	X	X	1	X	X	0		
5	X	X	X	1	0	X	1	1		
									N_1	C_1
13 + 5	X	X	X	X	1	X	X	1	2	YES

Figure 3.8 Table for Finding 1-Cubes Covering Point 13 in Example 1

	c_0	c_1	c_2
0	-	-	e
2	-	-	d
3	-	-	e
4	-	-	d
5	-	-	d
6	-	-	d
7	-	-	d
13	-	e	0

Figure 3.9 Matrix L_3 of Example 1



Example 3.10 Karnaugh Map for Example 1

Remarks: Comparing with the Karnaugh map method in Fig. 10, the result is obtained algebraically. The so-called "less than cube" (4,5,6,7) was not generated.

Example 2: Branching:

$$F = (0, 2, 3, 4, 5, 7).$$

Step 1: After obtaining matrices A, B, and L (Figs. 11, 12, and 13), it is observed that none of the points are possible core points, and that each is multiple 1-coverable. Starting arbitrarily with point 0, $\langle 0,2 \rangle$, and $\langle 0,4 \rangle$ are found. Points 2 and 4 are the next nodes, and L is updated to L_j (Figs. 14 and 15).

Step 2: From point 2, point 3 is the next node with $\langle 2, 3 \rangle$ the connecting cube. From 4, point 5 is the next node with $\langle 4,5 \rangle$ the connecting cube. L_1 is updated to L_2 (Figs. 16, 17 and 18).

Step 3: Point 3 goes to point 7 by cube $\langle 3,7 \rangle$ and point 5 goes to point 7 by cube $\langle 5,7 \rangle$. The next node, point 7, in both cases, is a terminal node (Figs. 19, 20 and 21).

Step 4: The branching graph for the connected j -coverable points and the cube-graph corresponding to it are developed in Fig. 22 (a)

	0	2	3	4	5	7
0	0	1	2	1	2	3
2	1	0	1	2	3	2
3	2	1	0	3	2	1
4	1	2	3	0	1	2
5	2	3	2	1	0	1
7	3	2	1	2	1	0

Figure 3.11 Matrix A of Example 2

	ρ_0	ρ_1	ρ_2	ρ_3
0	1	2	2	1
2	1	2	2	1
3	1	2	2	1
4	1	2	2	1
5	1	2	2	1
7	1	2	2	1

Figure 3.12 Matrix B of Example 2

	c_0	c_1	c_2
0	-	-	0
2	-	-	0
3	-	-	0
4	-	-	0
5	-	-	0
7	-	-	0

Figure 3.13 Matrix L of Example 2

	0	2	3	4	5	7
0	0	1	X	1	X	X
2	1	0	1	X	X	X
4	1	X	X	0	1	X

							N_2	C_2
0 + 2	1	1	X	X	X	X	2	YES
0 + 4	1	X	X	1	X	X	2	YES

Figure 3.14 Table for Finding 1-Cubes Covering Point 0 in Example 2

	c_0	c_1
0	-	d
2	-	n
3	-	-
4	-	n
5	-	-
7	-	-

Figure 3.15 Matrix L_1 of Example 2

	0	2	3	4	5	7			
2	1	0	1	X	X	X			
3	X	1	0	X	X	1			

							N_1	C_1
2 + 3	X	1	1	X	X	X	2	YES

Figure 3.16 Table for Finding 1-Cubes Covering Point 2 of Example 2

	0	2	3	4	5	7		
4	1	X	X	0	1	X		
5	X	X	X	1	0	1		

							N ₁	C ₁
4 + 5	X	X	X	1	1	X	2	YES

Figure 3.17 Table for Finding 1-Cubes Covering Point 4 of Example 2

	C_0	C_1
0	-	d
2	-	d
3	-	d
4	-	d
5	-	d
7	-	-

Figure 3.18 Matrix L_2 of Example 2

	0	2	3	4	5	7		
5	X	X	X	1	0	1		
7	X	X	1	X	1	0		

							N_1	C_1
5 + 7	X	X	X	X	1	1	2	YES

	0	2	3	4	5	7		
3	X	1	0	X	X	1		
7	X	X	1	X	1	0		

							N_1	C_1
3 + 7	X	X	1	X	X	1	2	YES

Figure 3.19 Table for Finding 1-Cubes Covering Points 3 and 5 of Example 2

	c_0	c_1
0	-	d
2	-	d
3	-	d
4	-	d
5	-	d
7	-	d

Figure 3.20 Matrix L_3 of Example 2

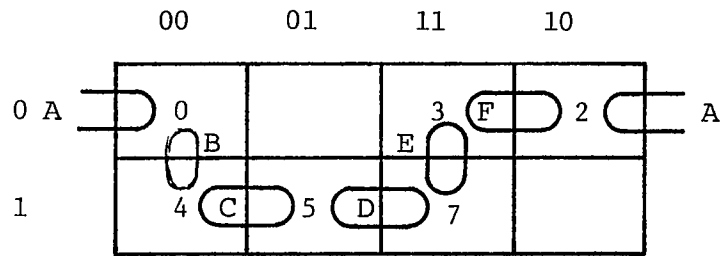
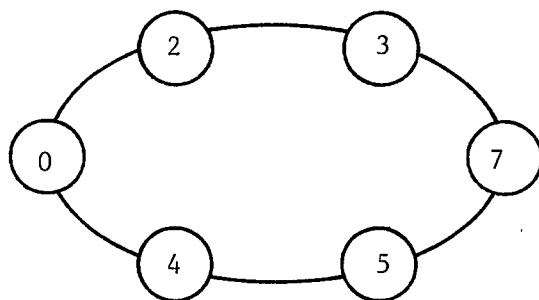
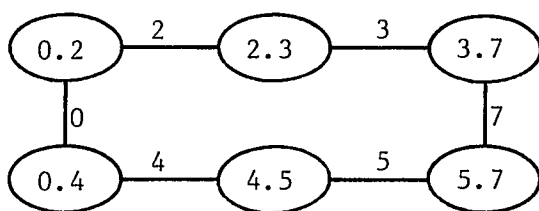


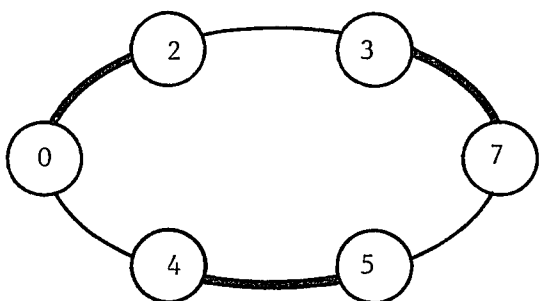
Figure 3.21 Karnaugh Map of Example 2



(a)



(b)



(c)

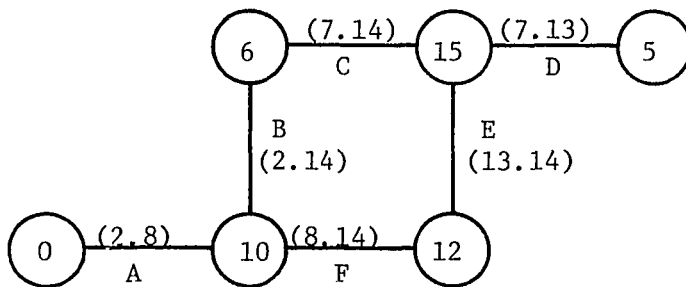
Figure 3.22 Branching Graph and Cube Graph of Example 2

and 22 (b) respectively. The minimum set includes every other branch in the branching graph (Fig. 22 (c)). (Theorem 7, Corollary 1).

Example 3: Covering of simple branching graph:

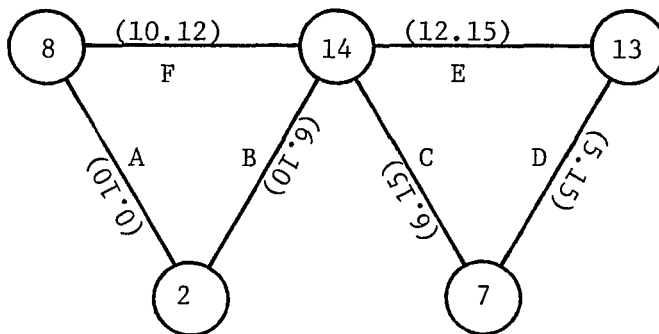
$$f(x) = \Sigma(0, 2, 5, 6, 7, 8, 10, 12, 13, 14, 15)$$

The two equivalent branching graphs for $k=2$ are shown in figures 23(a) and (b). Each of these will give the same answer. However, figure 23(a) is preferred because it has two terminal nodes of the first kind. From figure 23(a), it is clear that branches A and D are essential in any minimal set of non-redundant maximal cubes. As for the loop consisting of branches B, C, E, and F, either B and E, or C and F can be selected. If the branching graph shown in figure 23(b) is used, then there are two loops. For the loop on the left side which consists of branches A, B and F, A is chosen first. For the loop on the right side which consists of branches C, D and E, D is chosen first. Because node 14 has four branches connected to it while nodes 2, 8, 7 and 13 each have two branches connected to it, therefore either B or F in the left side loop is chosen. The last branch needed to complete the covering the loop on the right side depends on the previous selection of B or F. If B were chosen, E is required. If F were chosen, C is required. There is



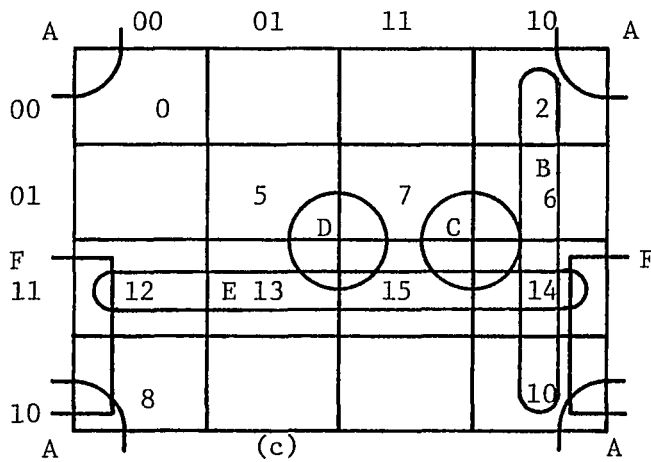
K = 2

(a)



K = 2

(b)



(c)

Figure 3.23 Branching Graphs and Karnaugh Map for Example 3

more chance for making erroneous selections in using the branching graph of figure 23(b). This does not pose any difficulty because if a core point is used as a starting node, a branching graph similar to that in figure 23(a) will result. The solution using figure 23(b) illustrates the equivalency between branching graphs. Figure 23(c) is the Karnaugh map for Example 3.

Example 4: Branching graph of different dimensions with inseparable linkage:

$$f(x) = \Sigma(0, 1, 2, 5, 6, 7, 8, 10, 12, 13, 14, 15)$$

The branching graph for 2-coverable points is the same as in Example 3. In Example 4, an additional branching graph for 1-coverable points is shown in Figure 24(a). Figure 24(b) depicts the linking of the branching graph for the 1-coverables to the branching graph for the 2-coverables. The thick lines indicate the chosen branches, i.e., branches A, C, E, and H. Branches C and E are chosen for the circuit B, C, E, and F. Applying Definition 15 to eliminate branch D, A is essential. Therefore we find B and F are not necessary for the circuit of B, C, E, and F. (Figure 24(c) is the Karnaugh map for Example 4).

Example 5: Branching graph covering with linking for different k but separable:

$$f(x) = (0, 1, 4, 5, 8, 9, 12, 13, 16, 17, 18, 19, 20, 21,$$

22,23,24,25,26,27,28,29,30,31,32,33,36,37,
 40,41,43,44,45,50,51,52,53,54,55,58,59,60,
 61,62,63)

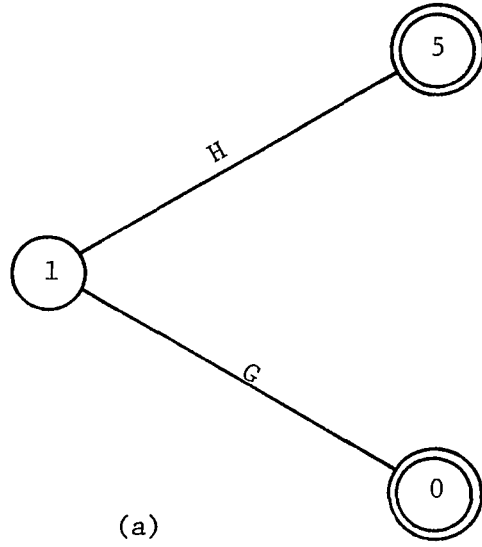
The branching graphs for $k = 4$ and $k = 1$ are shown in Figures 25(a) and (b). Although the branching graph is linkable, the linkage of the branching graph does not meet the requirement of Definition 15. Therefore, they are separate branching graphs. (Figure 26 is the linked branching graph. Figure 27 is the Karnaugh map for Example 5. Figure 28 is one of the equivalent branching graphs for 4-coverables in Example 5.)

Example 6: Having unspecified minterms:

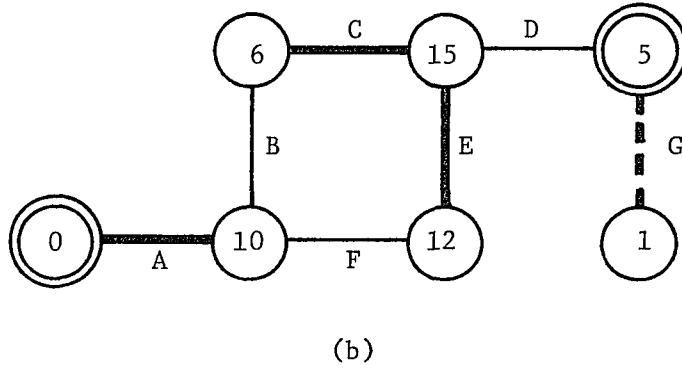
$$f(x) = \Sigma(2,10,12) + \Sigma_d(0,1,3,5,7,8,9,11,14,15)$$

Matrices A, B, and L are shown in Figures 29(a), (b) and (c) respectively. Since only the points in the on-set (2,10,12) are necessary, these representing don't cares do not appear in matrices A, B and L. These matrices are therefore rectangular. From Figures 29(d) and 29(g), cubes $\langle 2, 9 \rangle$ and $\langle 12,10 \rangle$ were found, and Figure 29(g) indicates that all on-set points are covered.

The differences in having don't cares are: (1) The rows representing don't cares do not appear in matrices



$K = 1$



(c)

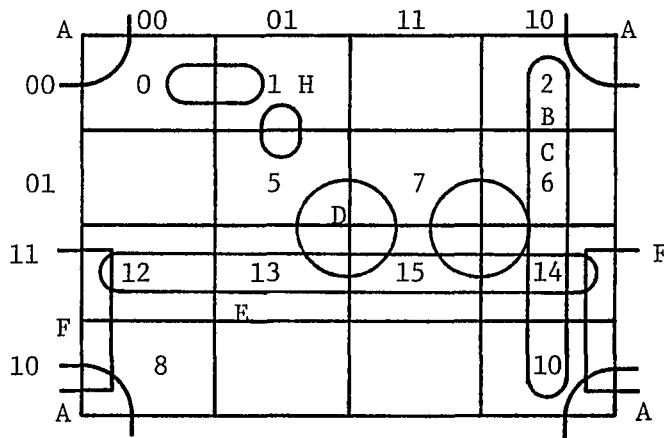
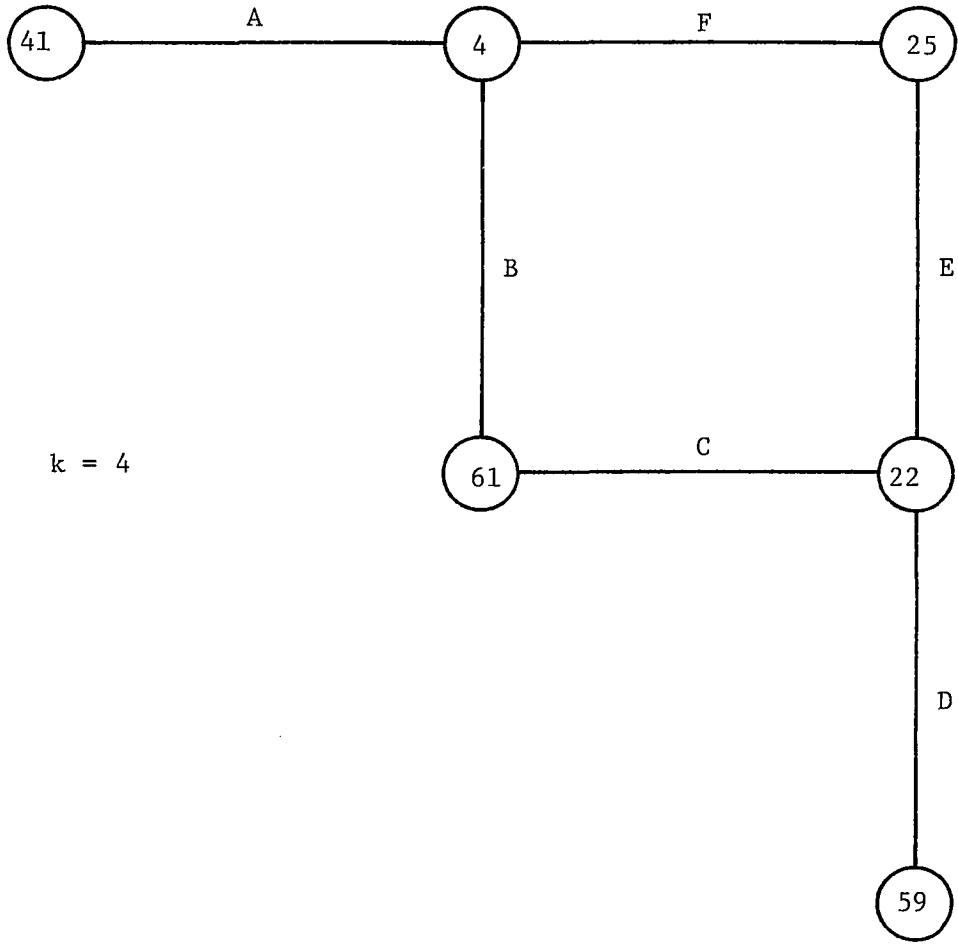
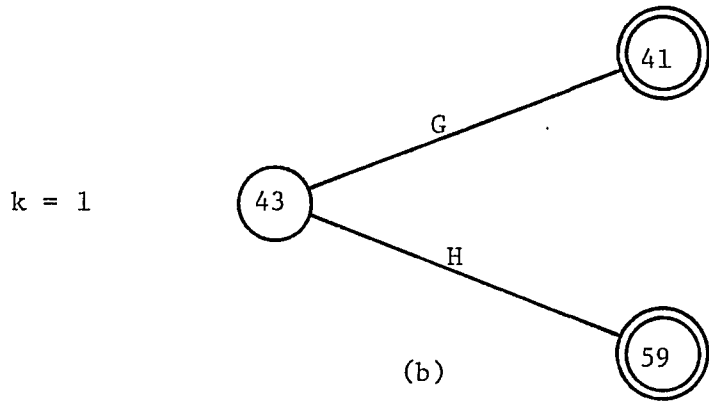


Figure 3.24 Branching Graphs and Karnaugh Map for Example 4



(a)



(b)

Figure 3.25 Branching Graph for Example 5

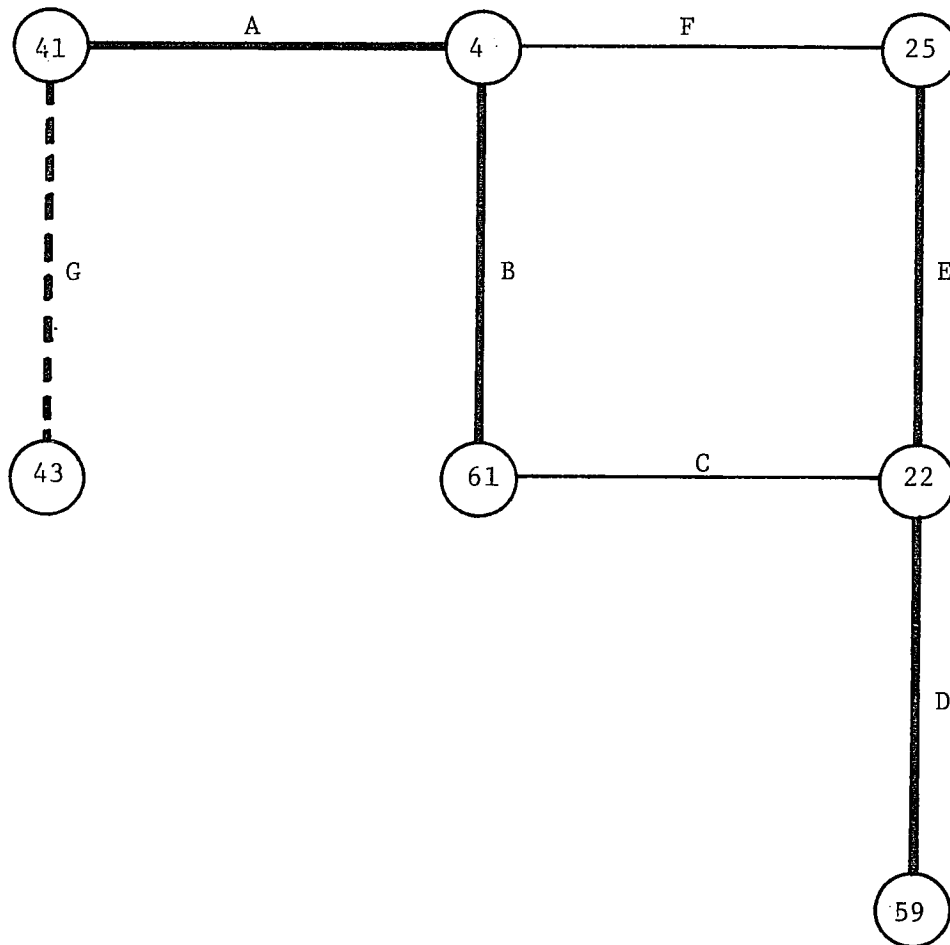


Figure 3.26 Branching Graph (Jointed and Covered)
for Example 5

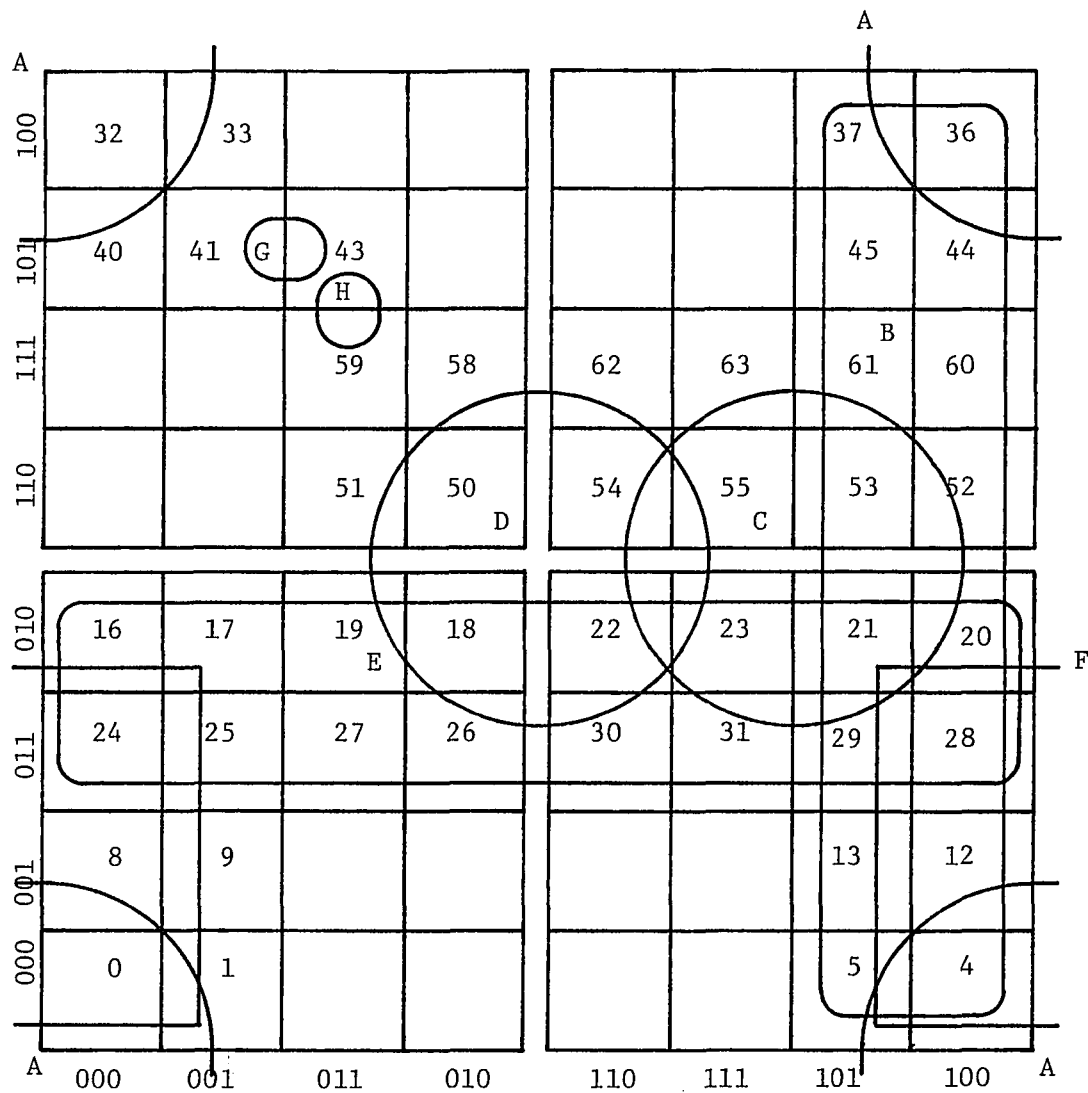


Figure 3.27 Karnaugh Map for Example 5

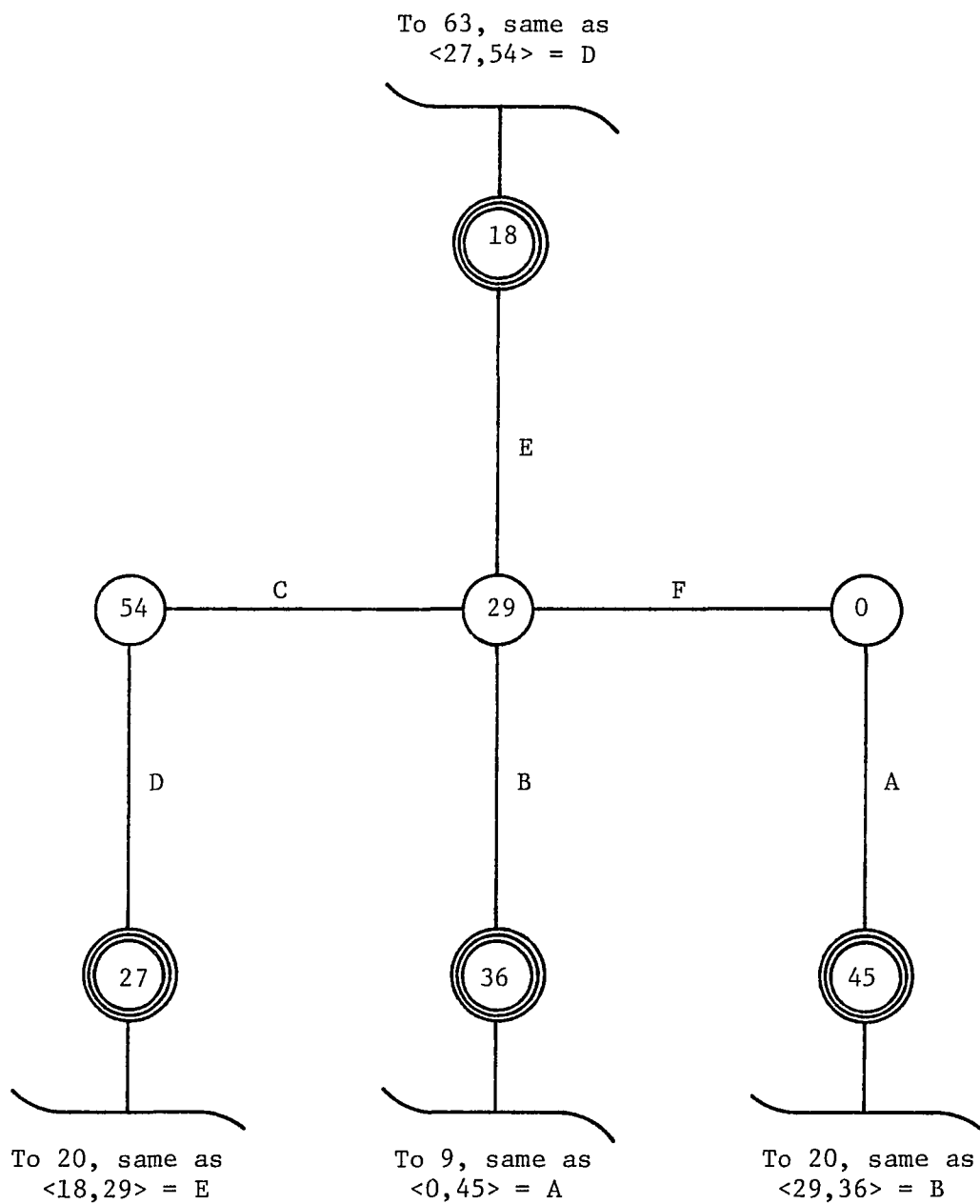


Figure 3.28 Branching Graph for Example 5
(Different but Isomorphic to Figure 3.25)

$$\begin{array}{c}
 2 \quad 10 \quad 12 \quad 0 \quad 1 \quad 3 \quad 5 \quad 7 \quad 8 \quad 9 \quad 11 \quad 14 \quad 15 \\
 \left(\begin{array}{cccccccccccc}
 0 & 1 & 3 & 1 & 2 & 1 & 3 & 2 & 2 & 3 & 2 & 2 & 3 \\
 1 & 0 & 2 & 2 & 3 & 2 & 4 & 3 & 1 & 2 & 1 & 1 & 2 \\
 3 & 2 & 0 & 2 & 3 & 4 & 2 & 3 & 1 & 2 & 3 & 1 & 2
 \end{array} \right)
 \end{array}$$

(a)

$$\begin{array}{c}
 \rho_0 \quad \rho_1 \quad \rho_2 \quad \rho_3 \quad \rho_4 \\
 \left(\begin{array}{ccccc}
 1 & 3 & 5 & 4 & 0 \\
 1 & 4 & 5 & 2 & 1 \\
 1 & 2 & 5 & 4 & 1
 \end{array} \right)
 \end{array}$$

(b)

$$\begin{array}{c}
 c_0 \quad c_1 \quad c_2 \quad c_3 \quad c_4 \\
 \left(\begin{array}{ccccc}
 - & - & - & c & 0 \\
 - & - & - & - & 0 \\
 - & - & c & 0 & 0
 \end{array} \right)
 \end{array}$$

(c)

Figure 3.29 Diagrams for Example 6

- (a) Matrix A
- (b) Matrix B
- (c) Matrix L

2 10 12 0 1 3 5 7 8 9 11 14 15

2	0	1	3	1	2	1	3	2	2	3	2	2	3
12	3	2	0	2	3	X	2	3	1	2	3	1	2
5	3	X	2	2	1	2	0	1	3	2	3	3	2
9	3	2	2	2	1	2	2	3	1	0	1	3	2
15	3	2	2	X	3	2	2	1	3	2	1	1	0

N_3 C_3

2 + 12	3	3	3	3	X	X	X	X	3	X	X	3	X	6	NO
2 + 5	3	X	X	3	3	3	3	3	X	X	X	X	X	6	NO
2 + 9	3	3	X	3	3	3	X	X	3	3	3	X	X	8	YES
2 + 15	3	3	X	X	X	3	X	3	X	X	3	3	3	7	NO

(d)

$$\begin{array}{c}
 C_0 \quad C_1 \quad C_2 \quad C_3 \\
 \left. \begin{array}{l}
 2 \\
 10 \\
 12
 \end{array} \right\} \begin{array}{l}
 - \quad - \quad - \quad e \\
 - \quad - \quad - \quad d \\
 - \quad - \quad c \quad 0
 \end{array}
 \end{array}$$

(e)

Figure 3.29 (cont.)

(d) Table for Finding 3-Cubes

(e) Matrix L_1

	2	10	12	0	1	3	5	7	8	9	11	14	15
12	X	2	0	2	X	X	2	X	1	2	X	1	2
10	X	0	2	2	X	X	X	X	1	2	X	1	2
0	X	2	2	0	X	X	2	X	1	2	X	X	X
5	X	X	2	2	X	X	0	X	X	2	X	X	2
9	X	2	2	2	X	X	2	X	1	0	X	X	2
15	X	2	2	X	X	X	2	X	X	2	X	1	0

Figure 3.29 (continued)

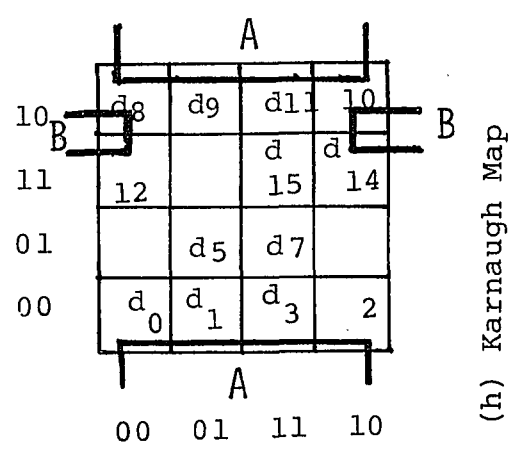
N_2 C_2

12 + 10	X	2	2	X	X	X	X	X	2	X	X	2	X	4	YES
12 + 0	X	X	2	2	X	X	X	X	2	X	X	X	X	3	NO
12 + 5	X	X	2	X	X	X	2	X	X	X	X	X	X	2	NO
12 + 9	X	X	2	X	X	X	X	X	2	2	X	X	X	3	NO
12 + 15	X	X	2	X	X	X	X	X	X	X	X	2	2	3	NO

(f) Table for Finding 2-Cubes

	C_0	C_1	C_2	C_3	
2	-	-	-	e	(g) Matrix L_2
10	-	-	-	d	
12	-	-	e	0	

(g)



(h) Karnaugh Map

A, B and L; (2) The distance row has to be established, only when the don't care point is a potential candidate for a diametrical extreme of a cube that covers an on-set point. Even then, only the relevant entries appear in the row. For example, in Figure 29(d), only the rows representing don't care points 5, 9 and 15 are needed. In Figure 29(f), although the don't care points 0,5,9 and 15 are needed, the entries of these rows under the columns of 2,1,3,7 and 11 are irrelevant. In this case, we are considering a 2-cube covering for point 12, yet all the distances from the points (2,1,3,7 and 11 to 12) are greater than 2. Therefore, a 3-cube covering cannot be formed. In the case where the majority of points are don't cares, then a covering for the on-set and the off-set would each be found and then combined by Nelson's method (N3) or McCluskey's method (Mc2).

3.3 CONCLUSION

The concept of distance matrices and distance-distribution matrices leads to the algorithm of finding the maximal cube covering starting with the largest possible cube first. This is the reverse procedure to those using conventional tabular methods. Obviously, when larger cubes exist, it is thus more effective. However, when there are numerous small cubes, the algorithm tends to increase the labor while the conventional methods tend to decrease it.

Initial work has not been effective in ascertaining the break even point. Presently, computer programming to minimize the effort is being studied.

Although similar in some respects to the map method, this method is more powerful and effective. Firstly, it does not depend upon the human faculty of recognizing geometric patterns; and secondly, it is systematic and precise and can be readily computerized.

The concept of the branching graph not only leads to solving the PI table covering problem more effectively, but also gives a definite way of finding all the maximal cubes.

Since the new approach breaks down the on-set of a switching function into connected j -coverable subsets, the maximal cube table (or prime-implicants table) evolves into an easier solution. In fact, the branching graph shows more interrelationships between the maximal cubes and the elements of the on-set of a switching function than the row dominance and the column dominance in a prime-implicant table. Such relations include the diametrical opposite points of a cube, the elimination of less-than cubes, the true non-essential maximal cubes, and the connectivity of the maximal cubes.

Because of the corresponding characteristics between the Karnaugh map and the new approach, it is conceivable

that many problems in switching theory and logic design can be solved with the new approach. These may include the multi-output problems, the switching functions with feedback and sequential digital circuits. The full impact of this method cannot be realized until more types of problems are studied.

CHAPTER IV

THE APPLICATION OF GEOMETRIC APPROACH

4.1 Algorithm for two level Logic Minimization.

Following the analysis in Chapter 2 Section 4, we can define the characteristics function for an array of the on-set and don't cares of a switching function as following:

Characteristic 1: The highest coverable dimension.

Characteristic 2: Core point candidacy.

Characteristic 3: Probability of sharing common parts.

Characteristic 4: The accumulative awayness.

The general algorithms as follows:

MAIN PROGRAM

- (1) Read in number of variables and number of don't care minterms, minterm zero is included or not, data card of all minterms in the on-set and don't care set.
- (2) Calculate the maximal cube dimension and binomial coefficients for that dimension.
- (3) Calculate the distance Matrix A for all the minterms.
- (4) Write matrix A into disc file.
- (5) Calculate the distance frequency distribution for each minterm. Then immediately calculate Char. 1 and Char. 2.
- (6) If $d(I, X) > 1$ for all X in the array, I is isolated.
- (7) Print out all isolated minterms.

- (8) Char (I,1) = Characteristic 1 of minterm I = the highest coverable dimension for satisfying the following two criteria.

$$(i) \quad l_{ij} = \underset{k=0}{\overset{k=j}{\text{Min}}} \binom{bik}{\binom{j}{k}} = \text{Char}(I,1) \quad (\text{Theorem 3})$$

- (ii) The number of minterms in the array meeting condition 1 is more than 2^1_{ij} .

- (9) If Frequency of distance = 1 or 2^j where j is the highest coverable cube covering, then the minterm in examination is a candidate of core point. Char. (I,2) = 1 or 2^j , otherwise Char. (I.2) = 5 or 999999, which means multiple covering.

- (10) Check if Maximal Cube dimension is valid or not. If not, lower Maximal-Cube dimension by 1. The criterion is \log_2 (number of maximal cube coverable minterms) \geq maximal cube dimension.

- (11) Characteristic 3 = Char (I,3) = 0 for on-set
= 999999 for don't care.

at the start of the program. It will change as the process of finding cube covering progresses. At the start the on-set are not common with any cube covering and the don't care minterms are in common with every cube covering. Therefore the probability of sharing any common part for the on-set minterms is

zero, for don't care minterms is 1, for which we use $999999/10^6$ to approximate and then conveniently drop the common denominator, 10^6 .

- (12) Test all the core point candidates as to whether there is a cube or not. If the cube exists, then this minterm is covered and all the minterms in this cube are also covered. Therefore their Characteristic 3 all become 99999. The probability of sharing is certain for future cube covering. If not, then this minterm being examined becomes a j-1 coverable minterm.

After a core point and core cube is established, the accumulative awayness is computed for the rest of the minterms in the array.

- (13) Use the criteria that a new cube should have the least awayness from and the least intersection with the don't care set or covered set. The rest of the array is gradually covered with the minimum maximal-cube covering for the whole array.

SUBROUTINE CORE CUBES

- (1) Pick any minterm that is a core point candidate; if all possible candidates have been checked stop.
- (2) Check characteristic 3 to single out don't cares.
- (3) Use characteristic 1 for dimension. Call subroutine CUBE to verify the existence of the core cube.

- (4) If core point candidate does not have a cube covering, lower the dimension and charge the char (3.2) to non core candidate the minterm being processed.
- (5) If indeed the cube exists, then print out the data for this cube.
- (6) If not finished, return to step 1.

SUBROUTINE MULTIPLE COVERABLE POINTS.

1. Select starting point.
 - a. Highest coverable by Char (I,1)
 - b. Smallest awayness
 - c. Largest aloofness.
 - d. Least probability to share common parts with existing maximal-cube covering.
2. Call subroutines cubes to verify the existence of all the highest maximal cubes.
3. Check the actual number of common parts between the cubes found in 2. and choose the one that has the fewest common parts.
4. Checking less-than-cubes by characteristic 1 to check actual coverable dimension.
5. Starting dimension 1 to dimension maximal.
6. Go through 1 through 3.

SUBROUTINE CUBES

1. Called by main program minterm = MTH, Dimension = D, Type = ZNTH.

2. Read from file the distance function for Mth.
3. If ZNTH = 1, core point candidate verify the existing of a cube and print out otherwise, go to 5.
4. Calculate accumulated awayness, relative probability of showing common parts go to 7.
5. Calculate probability of share verigy all D-cubes covering MTH. If awayness are the same, then compare the probability of share.
6. Calculate accumulative awayness and aloofness.
7. Return to main program. The detailed flow diagrams are shown in Figure 4.1 for the main program. Figure 4.2 shows the sub-routine of finding the core cubes. 4.3 is program for finding the remainder of the minimum maximal cube covering, and Figure 4.3 shows the sub-routine cube identifying or cube. The listing of the Fortran IV programs is given in the appendix. A detailed example is worked out in the next section to demonstrate the mechanism of the program.

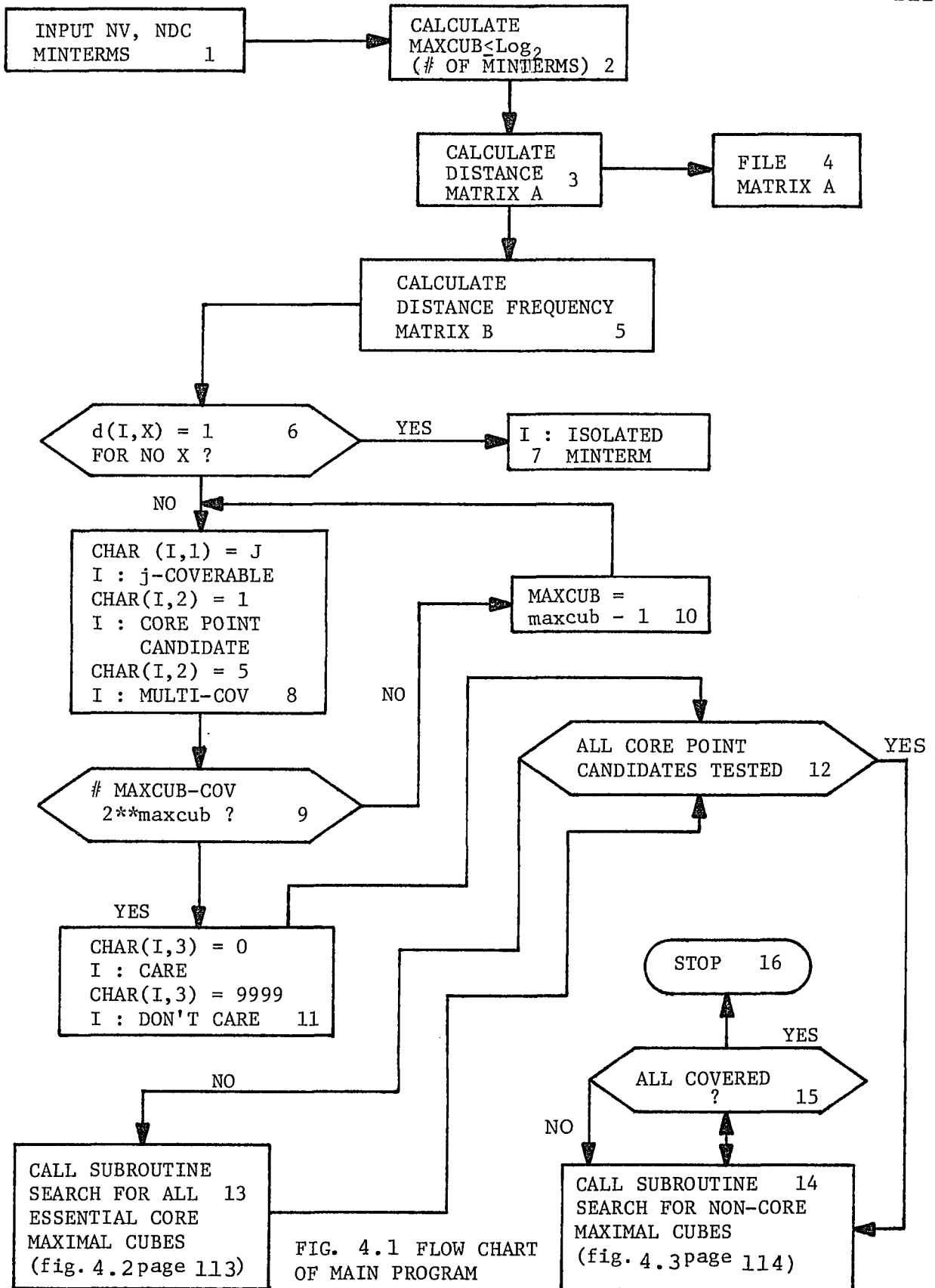


FIG. 4.1 FLOW CHART OF MAIN PROGRAM

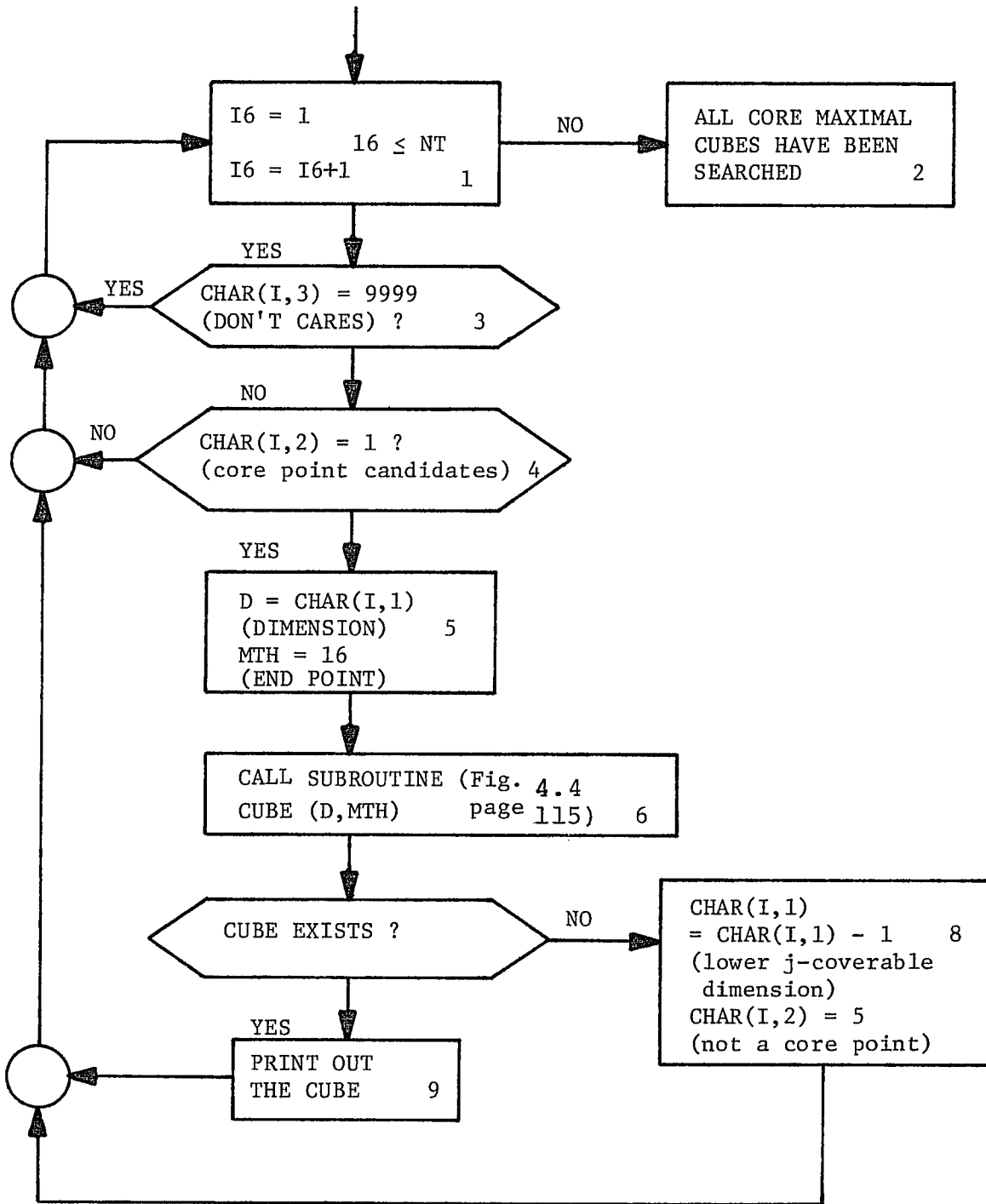


FIGURE 4.2 Flow Chart of Subroutine Search for All Essential Core Maximal Cubes (called by main program in Figure 4.1, Block 13)

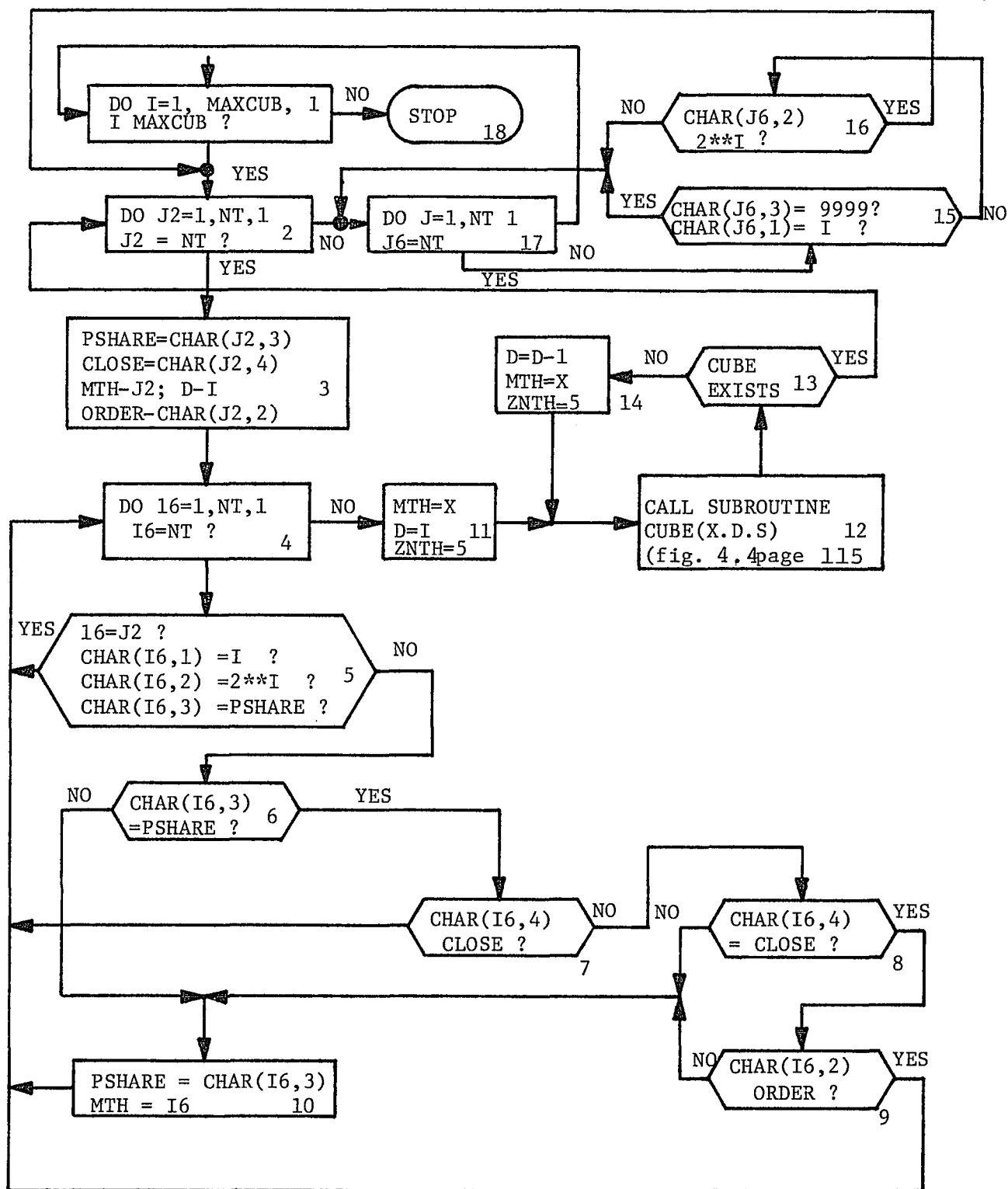


FIGURE 4.3 SUBROUTINE SEARCHING FOR NON-CORE MAXIMAL CUBES

(Called by main program fig.4.4 block 14, page 112)

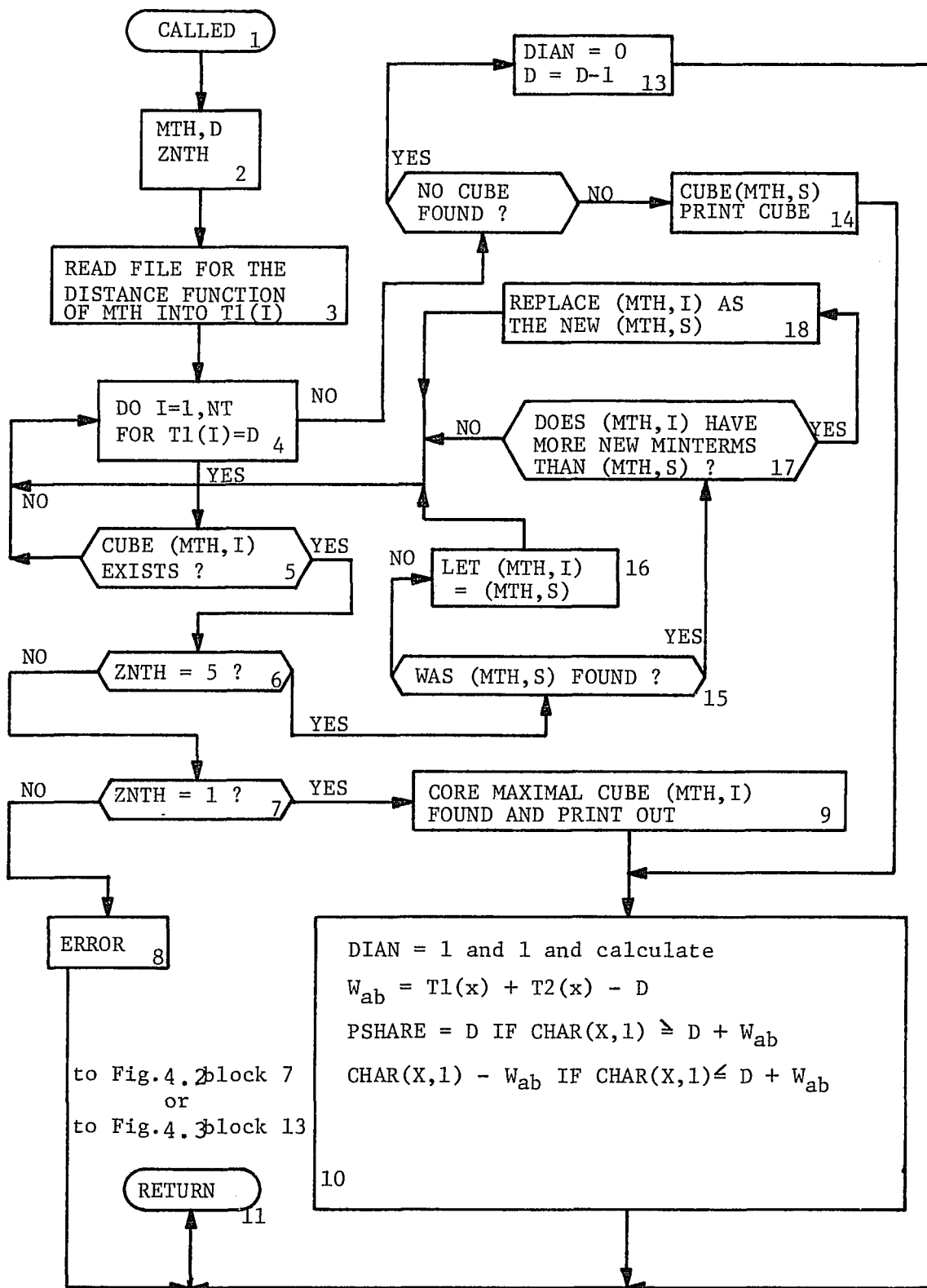


FIGURE 4.4 FLOW CHART OF SUBROUTINE CUBE
(Called by Fig.4.2 block 6, or Called by Fig.4.3 block 12)

4.2 Computer format and mechanism illustration. The computer program implemented in the Univac VMOS SYSTEM at New Jersey Institute of Technology can be illustrated as following:

1) Step 1. entering

the number of don't cares, and the number of variables.

Then enter all the minterms (on-set plus don't cares with don't cares after on-set as follow:

```
* 5      5      5
* 1      2      3      4      5      6      7      9      11      12      13
*14     15     17     20     23     24     25     27     28
*9
```

Step 2: The computer will calculate the distance function and store them in a disc file. Then calculate all the characteristic and point out

MAXIMUM CUBE DIMENSION = 4

```
1  0  0  0  2  1  0  0  3  3  1  0  4
6  4  1
```

THE CHARACTERISTICS:

```
4  1  2  1  4  1  4  1  4  1  4  1  4
5  4  1  4  1  4  1  4  1  3  1  4  1
2  1  2  1  1  1  2  1  3  5  2  1  3
1
```

Step 3, the computer will start to search all the

core cubes first; and point out, the relative probability of sharing (checking points), the accumulated awayness, the cube found and cube selected. When all the core cubes has been reached, it will then point out a statement of the fact and the maximum cube dimension of all the core cubes. A sample is as following:

CHECKING POINTS

```

  0 9999 9999    2    2 9999 9999    0    2    0    0    2
  2  0    0 9999 9999 9999 9999 9999

```

ACCUMULATED AWAYNESS:

```

  2  0  0  2  2  0  0  4  2  4  4  2
  2  4  4  0  0  0  0  0

```

CUBE (2 . 7) IS SELECTED.

CUBE (14 . 5) HAS BEEN FOUND.

CHECKING POINTS

```

  4 9999 9999 9999 9999 9999 9999    2    4 9999 9999 9999
  9999  0  1 9999 9999 9999 9999 9999

```

ACCUMULATED AWAYNESS:

```

  4  0  0  2  2  0  0  6  4  4  4  2
  2  8  6  0  0  0  0  0

```

CUBE (14 , 5) IS SELECTED

CUBE (17 , 9) HAS BEEN FOUND.

CHECKING POINTS

```

  9999 9999 9999 9999 9999 9999 9999 9999    6 9999 9999 9999
  9999 9999    1 9999 9999 9999 9999 9999

```

ACCUMULATED AWAYNESS :

4	0	0	2	2	0	0	6	6	4	4	2
2	8	10	0	0	0	0	0				

CUBE (17 , 9) IS SELECTED.

CUBE (20 , 12) HAS BEEN FOUND.

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	9999	6	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999				

ACCUMULATED AWAYNESS:

4	0	0	2	2	0	0	6	12	4	4
2	2	8	10	0	0	0	0	0		

CUBE (20 , 12) IS SELECTED

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS: 3

Step 4: after all the core cubes has been found, then it start to look for the rest of the minimum-set of maximal cube covering and prints out like this:

CUBE (11 , 5) HAS BEEN FOUND

CHECKING POINTS

```

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999

```

ACCUMULATED AWAYNESS:

4	0	0	2	2	0	0	6	12	4	4
2	2	8	10	0	0	0	0	0		

CUBE (11 , 5) IS SELECTED

**FORTRAN ** STOP

Notice the computer also prints out two up-dated value of characteristic. CHECKING POINTS are the relative probability of sharing. Accumalated Awayness is also printed out. This is for de-bugging and applications other than just the minimum set. For minimization only all this can be omitted. At this point, all the calculation is completed. And the results is output as follow:

```

Total Number of On-set and Dont*Cares = 20
ON-SET = 15  DONT-CARES = 5
NUMBERS OF VARIABLES = 5

```

MINT.# 1 = 1 MINT.# 2 = 2 MINT.# 3 = 3
 MINT.# 4 = 4 MINT.# 5 = 5 MINT.# 6 = 6
 MINT.# 7 = 7 MINT.# 8 = 9 MINT.# 9 = 11
 MINT.# 10 = 12 MINT.# 11 = 13 MINT.# 12 = 14
 MINT.# 13 = 15 MINT.# 14 = 17 MINT.# 15 = 20
 MINT.# 16 = 23 MINT.# 17 = 24 MINT.# 18 = 25
 MINT.# 19 = 27 MINT.# 20 = 28 MINT.#

THE CUBE COVERING LIST:

THE 1-TH CUBE COVERING IS:

2-CUBE = (2 , 7) = 00-1-

THIS CUBE COVERS THESE MINTERMS †

2 7 3 6

THIS MINTERM $MI(2) = 2$ IS A CORE POINT

THE 2-TH CUBE COVERING IS:

3-CUBE = (14 , 5) = 0-1--

THIS CUBE COVERS THESE MINTERMS:

14 5 4 6 7 12 13 15

THIS MINTERM $MI(12) = 14$ IS A CORE POINT

THE 3-TH CUBE COVERING IS:

2-CUBE = (17 , 9) = --001

THIS CUBE COVERS THESE MINTERMS:

17 9 1 25

THIS MINTERM $MI(14) = 17$ IS A CORE POINT

THE 4-TH CUBE COVERING IS:

2-CUBE = (20 , 12) = --100

THIS CUBE COVERS THESE MINTERMS:

20 12 4 28

THIS MINTERM $MI(15) = 20$ IS A CORE POINT

THE 5-TH CUBE COVERING IS:

3-CUBE = (11 , 5) = 0---1

THIS CUBE COVERS THESE MINTERMS:

11 5 1 3 7 9 13 15

THE FINAL CHECKING

```

205 101 205 205 205 205 205 205 105 205 205 101
205 101 101 109 109 209 109 209

```

The final checking indicates what status of the maximal cube coverings and also checks whether all on-set is covered.

A second completed illustrative example is reproduced as following to show other features.

In the appendices, a collection of various example problems has been selected with the out-put result print out. For reference purpose and better understanding of the mechanics of the computer program. Those checking and debug out-put print out are also enclosed following some of the examples results out-put print.

Illustrative Examples 2.

FORTRAN IV PROGRAM LARGEE STARTED ---

```

6 9
 0  1  2  5  8
.13 18 19 28 34 37 39 47 56 57 61
 4 12 16 20 23 38 49 59 629

```

MAXIMUM CUBE DIMENSION = 4

```

1 0 0 0 2 1 0 0 3 3 1 0 4 6
4 1

```

THE CHARACTERISTICS:

```

3 5 2 1 3 1 3 5 2 1 2 1 3 1 2 1 2 1 2 1
3 1 1 1 1 1 1 5 1 1 4 1 3 5 3 1 3 1 1 1 2 1
1 1 1 1 0 0

```

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 3

CUBE (1 , 4) HAS BEEN FOUND.

CHECKING POINTS

```

9999 9999      2 9999      1  1  0  0  0  0  1  0
  0  0      0  0 9999 9999 9999 9999 9999 9999 9999
9999

```

ACCUMULATED AWAYNESS:

```

  0  0  2  0  2  2  4  4  4  4  2  4
  6  6  6  6  0  0  0  0  0  0  0  0
  0

```

CUBE (1 , 4) IS SELECTED.

CUBE (8 , 4) HAS BEEN FOUND.

CHECKING POINTS

```

9999 9999      3 9999 9999      2  0  0  1  0  1  0
  0  0      0  0 9999 9999 9999 9999 9999 9999 9999
9999

```

ACCUMULATED AWAYNESS:

```

  0  0  4  0  2  4  8  10  6  8  6  10
 12 10 12 12  0  0  0  0  0  0  0  0
  0

```

CUBE (8 , 4) IS SELECTED.

CUBE (13 , 4) HAS BEEN FOUND.

CHECKING POINTS

```

9999 9999      3 9999 9999 9999      0  0  2  0  2  0
  0  0      0  0 9999 9999 9999 9999 9999 9999 9999
9999

```

ACCUMULATED AWAYNESS:

```

  0  0  8  0  2  4  14  16  8  14  8  14
 16 16 18 16  0  0  0  0  0  0  0  0
  0

```

CUBE (13 , 4) IS SELECTED.

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

CUBE (28 , 4) HAS BEEN FOUND

CHECKING POINTS

```

9999 9999 3 9999 9999 9999 0 0 9999 0 2
0 0 0 0 0 9999 9999 9999 9999 9999 9999
9999 9999 9999

```

ACCUMULATED AWAYNESS:

```

0 0 12 0 2 4 18 22 8 20 12
20 22 20 24 20 0 0 0 0 0 0
0 0 0

```

CUBE (28 , 4) IS SELECTED.

CUBE (47 , 39) HAS BEEN FOUND.

CHECKING POINTS

```

9999 9999 3 9999 9999 9999 0 0 9999 0 2
9999 9999 0 0 0 9999 9999 9999 9999 9999 9999
9999 9999 9999

```

ACCUMULATED AWAYNESS:

```

0 0 18 0 2 4 26 28 8 24 14
20 22 28 30 24 0 0 0 0 0 0
0 0 0

```

CUBE (47 , 39) IS SELECTED.

CUBE (56 , 57) HAS BEEN FOUND.

CHECKING POINTS

```

9999 9999 3 9999 9999 9999 0 0 9999 0 2
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999

```

ACCUMULATED AWAYNESS:

0	0	36	0	2	4	40	40	8	38	24
20	22	28	30	26	0	0	0	0	0	0
0	0	0								

CHECKING POINTS

9999	9999	3	9999	9999	9999	0	0	9999	0	2
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999								

ACCUMULATED AWAYNESS:

0	0	36	0	2	4	40	40	8	38	24
20	22	28	30	26	0	0	0	0	0	0
0	0	0								

CUBE (61 , 57) IS SELECTED.

ALL CORE CUBES HAVE BEEN SEARCHED,

AND THE MAXIMUM CUBE DIMENSION IS : 2

CUBE (18 , 0) HAS BEEN FOUND.

CHECKING POINTS

```

9999 9999 9999 9999 9999 9999 9999 0 9999 0 2 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999

```

ACCUMULATED AWAYNESS:

```

0 0 36 0 2 4 40 42 8 40 8 40
22 28 30 26 0 0 0 0 0 0 0 0
30 20
0 0
0

```

```

CUBE ( 18 , 0 ) IS SELECTED
CUBE ( 34 , 2 ) HAS BEEN FOUND.
CUBE ( 34 , 38 ) HAS BEEN FOUND.
WATCH FOR 2 AND 38

```

CHECKING POINTS

```

9999 9999 9999 9999 9999 9999 9999 0 9999 9999 2 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999

```

ACCUMULATED AWAYNESS:

```

0 0 36 0 2 4 40 48 8 40 40 20
22 28 30 26 0 0 0 0 0 0 0 0

```

```

CUBE ( 19 , 23 ) IS SELECTED
CUBE ( 37 , 5 ) HAS BEEN FOUND.
CUBE ( 37 , 39 ) HAS BEEN FOUND.

```

CHECKING POINTS

```

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999

```

ACCUMULATED AWAYNESS:

```

0 0 36 0 2 4 40 48 8 40 40 20
22 28 30 26 0 0 0 0 0 0 0 0
0

```

CUBE (37 , 5) IS SELECTED.

**FORTRAN ** STOP

TOTAL NUMBER OF ON-SET AND DONT*CARES = 25

ON-SET = 16 DONT-CARES = 9

NUMBERS OF VARIABLES - 6

MINT.# 1 = 0	MINT.# 2 = 1	MINT.# 3 = 2
MINT.# 4 = 5	MINT.# 5 = 8	MINT.# 6 = 13
MINT.# 7 = 18	MINT.# 8 = 19	MINT.# 9 = 28
MINT.# 10 = 34	MINT.# 11 = 37	MINT.# 12 = 39
MINT.# 13 = 47	MINT.# 14 = 56	MINT.# 15 = 57
MINT.# 16 = 61	MINT.# 17 = 4	MINT.# 18 = 12
MINT.# 19 = 16	MINT.# 20 = 20	MINT.# 21 = 23
MINT.# 22 = 38	MINT.# 23 = 49	MINT.# 24 = 59
MINT.# 25 = 62	MINT.#	

MINTERM # 25 = 62 IS ISOLATED AND ESSENTIAL

THE CUBE COVERING LIST:

THE 1-TH CUBE COVERING IS :
 2-CUBE = (1 , 4) = 000-0-

THIS CUBE COVERS THESE MINITERMS :
 1 4 0 5

THIS MINITERM MI(2) = 1 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
 2-CUBE = (8 , 4) = 00--00

THIS CUBE COVERS THESE MINITERMS :
 8 4 0 12

THIS MINITERM MI(5) = 8 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
 2-CUBE = (13 , 4) = 00-10-

THIS CUBE COVERS THESE MINITERMS :
 13 4 5 12

THIS MINITERM MI(6) = 13 IS A CORE POINT

THE 4-TH CUBE COVERING IS :
 2-CUBE = (28 , 4) = 0--100

THIS CUBE COVERS THESE MINITERMS :
 28 4 12 20

THIS MINITERM MI(9) = 28 IS A CORE POINT

THE 5-TH CUBE COVERING IS :
 1-CUBE = (47 , 39) = 10-111

THIS CUBE COVERS THESE MINITERMS :
 47 39

THIS MINITERM MI(13) = 47 IS A CORE POINT

THE 6-TH CUBE COVERING IS :
 1-CUBE = (56 , 57) = 11100-

THIS CUBE COVERS THESE MINTERMS :
56 57

THIS MINTERM $MI(14) = 56$ IS A CORE POINT

THE 7-TH CUBE COVERING IS :
1-CUBE = (61 , 57) = 111-01

THIS CUBE COVERS THESE MINTERMS :
61 57

THIS MINTERM $MI(16) = 61$ IS A CORE POINT

THE 8-TH CUBE COVERING IS :
2-CUBE = (18 , 0) = 0-00-0

THIS CUBE COVERS THESE MINTERMS :
18 0 2 16

THE 9-TH CUBE COVERING IS :
1-CUBE = (34 , 38) = 100-10

THIS CUBE COVERS THESE MINTERMS :
34 38

THE 10-TH CUBE COVERING IS :
1-CUBE = (19 , 23) = 010-11

THIS CUBE COVERS THESE MINTERMS :
19 23

THE 11-TH CUBE COVERING IS :
1-CUBE = (37 , 5) = 00101

THIS CUBE COVERS THESE MINTERMS :
37 5

THE FINAL CHECKING

```

105 101 105 105 101 101 105 105 101 105 105 105
101 101 205 101 105 105 105 105 209 209 109 109
109

```

I) The tracing of cyclic cubes:

The computer is capable to trace out the cyclic cubes as shown in the following example.

```

TOTAL NUMBER OF ON-SET AND DONT-CARES = 14
ON-SET = 14 DONT-CARES = 0
NUMBERS OF VARIABLES = 5

```

```

MINT.# 1 = 1 MINT.# 2 = 4 MINT.# 3 = 9
MINT.# 4 = 11 MINT.# 5 = 12 MINT.# 6 = 27
MINT.# 7 = 26 MINT.# 8 = 29 MINT.# 9 = 28
MINT.# 10 = 18 MINT.# 11 = 22 MINT.# 12 = 23
MINT.# 13 = 21 MINT.# 14 = 0 MINT.#

```

THE CUBE COVERING LIST:

THE 1-TH CUBE COVERING IS :
 1-CUBE = (0 , 4) = 00-00

THIS CUBE COVERS THESE MINTERMS :
 4 0

THE 2-TH CUBE COVERING IS :
 1-CUBE = (12 , 28) = --1100

THIS CUBE COVERS THESE MINTERMS :
 12 28

THE 3-TH CUBE COVERING IS :
 1-CUBE = (21 , 29) = 1-101

THE 4-TH CUBE COVERING IS :
 1-CUBE = (22 , 23) = 1011-

THIS CUBE COVERS THESE MINTERMS :
 22 23

THE 5-TH CUBE COVERING IS :
 1-CUBE = (18 , 26) = 1-010

THIS CUBE COVERS THESE MINTERMS :
 26 18

THE 6-TH CUBE COVERING IS :
 1-CUBE = (27 , 11) = -1011

THIS CUBE COVERS THESE MINTERMS :
 11 27

THE 7-TH CUBE COVERING IS :
 1-CUBE = (1 , 9) = 0-001

THIS CUBE COVERS THESE MINITERMS :
 1 9

THE FINAL CHECKING

105 105 105 105 105 105 105 105 105 105 105 105
 105 105

The way it trace out of the cyclic cubes is shown in the checking
 print out of the computer as follow:

* 5
 * 1 4 9 11 12 27 26 29 28 18 22 23 21
 *9

MAXIMUM CUBE DIMENSION = 3

1 0 0 2 1 0 3 3 1

THE CHARACTERISTICS :

2 1 2 1 2 1 2 1 2 1 2 1 2 1 2 1 2 1 2 1
 2 1 2 1 2 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 1

ALL CORE CUBES HAVE BEEN SEARCHED,
 AND THE MAXIMUM CUBE DIMENSION IS : 1

CUBE (0 , 1) HAS BEEN FOUND.

CUBE (0 , 4) HAS BEEN FOUND.

WATCH FOR 1 AND 4

CHECKING POINTS

0 9999 0 0 0 0 0 0 0 0 0 0
0 9999

ACCUMULATED AWAYNESS :

2 0 4 6 2 8 6 6 4 4 4 6
4 0

CUBE (0 , 4) IS SELECTED.

CUBE (12 , 4) HAS BEEN FOUND.

CUBE (12 , 28) HAS BEEN FOUND.

CHECKING POINTS

0 9999 0 0 9999 0 0 0 9999 0 0 0
0 9999

ACCUMULATED AWAYNESS :

8 2 8 12 2 14 10 8 4 10 8 12
8 4

CUBE (12 , 28) IS SELECTED.

CUBE (21 , 29) HAS BEEN FOUND.

CUBE (21 , 23) HAS BEEN FOUND.

WATCH FOR 29 AND 23

CHECKING POINTS

0 0000 0 0 9999 0 0 9999 9999 0 0 0
9999 9999

ACCUMULATED AWAYNESS :

12 6 12 18 6 18 16 8 6 16 12 14
8 10

CUBE (21 , 29) IS SELECTED.

CUBE (22 , 18) HAS BEEN FOUND.

CUBE (22 , 23) HAS BEEN FOUND.

WATCH FOR 18 AND 23

CHECKING POINTS

0 9999 0 0 9999 0 0 9999 9999 0 9999 9999
9999 9999

ACCUMULATED AWAYNESS :

24 16 26 28 18 24 20 18 14 18 14 18
16 20

CUBE (18 , 26) IS SELECTED.

CUBE (27 , 11) HAS BEEN FOUND.

CUBE (27 , 26) HAS BEEN FOUND.

CHECKING POINTS

0 9999 0 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999

ACCUMULATED AWAYNESS :

28 24 28 28 24 24 22 22 20 22 20 22
 22 26

CUBE (27 , 11) IS SELECTED.

CUBE (1 , 9) HAS BEEN FOUND.

CUBE (1 , 0) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999

ACCUMULATED AWAYNESS :

28 28 28 30 28 28 28 26 26 28 28 28
 26 28

CUBE (1 , 9) IS SELECTED.

2) To find out isolated minterms :

```

-4
0  3  5  6 12 15  9 .109

```

MAXIMUM CUBE DIMENSION = 3

```

1  0  0  2  1  0  3  3  1

```

THE CHARACTERISTICS :

```

0  0  0  0  0  0  0  0  0  0  0  0  0  0  0  0

```

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 0

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 0

*9

**FORTRAN ** STOP

TOTAL NUMBER OF ON-SET AND DONT*CARES = 8

ON-SET = 8 DONT=CARES = 0

NUMBERS OF VARIABLES = 4

```

MINT.# 1 = 0 MINT.# 2 = 3 MINT.# 3 = 5
MINT.# 4 = 6 MINT.# 5 = 12 MINT.# 6 = 15
MINT.# 7 = 9 MINT.# 8 = 10 MINT.#

```

MINTERM # 1 = 0 IS ISOLATED AND ESSENTIAL

MINTERM # 2 = 3 IS ISOLATED AND ESSENTIAL

MINTERM # 3 = 5 IS ISOLATED AND ESSENTIAL

MINTERM # 4 = 6 IS ISOLATED AND ESSENTIAL

MINTERM # 5 = 12 IS ISOLATED AND ESSENTIAL

MINTERM # 6 = 15 IS ISOLATED AND ESSENTIAL

MINTERM # 7 = 9 IS ISOLATED AND ESSENTIAL

MINTERM # 8 = 10 IS ISOLATED AND ESSENTIAL

4.3 CONCLUSION:

The computer programming was done on a relatively slow and limited storage teaching computer to demonstrate the theory of minimal maximal cube coverings. It is possible to improve the efficiency and enlarge the scope of the programs with a larger and faster computer. It is also better implemented with fortran 77, in which case the two assembly subroutines can be replaced with fortran 77 subroutines.

The many extra features in the programming demonstrate the applicability in other areas of switching theory. The examples selected in appendix II, are typical to its size of on-set number of don't cares, and number of variables. The time used on CPU is not optimized because of the need of checking out all the steps and understanding the computer with various real examples. Further research is directed to develop a whole series of canned programs to do all the logic design on computers by optimizing the programming.

CHAPTER V

EXTENSION OF THE MINIMUM MAXIMAL-CUBE THEORY TO HIGHER CUBES AND OTHER KIND OF SWITCHING FUNCTIONS

5.1 EXTENSION TO HIGHER DIMENSIONED CUBES

In chapter II, section 4, we had discussed the distance of cubes higher than 0. (Ref. to definition 24 and theorem 19). It is natural for one try to generalizing the theory into higher dimensioned cubes and to try the minimization or decomposition of functions having different dimensioned cubes. In chapter II, we studied the interrelationship between cubes and found that not only the distance between cubes played an important part but also the skewness of cubes. The reason we did not use the skewness in setting up distance matrix for 0-cubes is because the awayness theorems (theorem 18 and theorem 19) indicated that skewness is always zero for 0-cubes. To handle higher cubes, one has to set up a complex matrix C with entries of a pair (d,s) instead of a single number, where d = the distance between cubes, and s = the skewness between cubes. This was followed by a complex matrix D with entries of a pair (f_d, f_s) where f_d = the frequency of distances and f_s = the frequency of skewness. After setting up matrices C and D , one then can establish the kind of characteristics one needs for minimization or decomposition. It appears to be very promising for further research work.

5.2 EXTENSION TO OTHER KIND OF SWITCHING FUNCTIONS

Since the whole theory started from a metric space, it is obvious to consider whatever the space and metric of a particular switching functions would be able to render a similar kind of theory like what we have developed in chapter II. In other words, the space Q^n can be the cartesian product of a prime number or the power of a prime number, thus almost all kind of multi-valued logic or Galois switching functions are included. However, the metric of the space other than clean is usually not unique and simple, and therefore quite cumbersome to work with. Recent research indicates that ternarylogics and Boolean Galois switching functions are among the most promising kind other than binary (or Boolean) logic. The metric of these functions has been explored extensively by coding theoreticians. We speculate from this vantage point that the maximal-cube covering approach will play important role in the coding theory. It is conceivable to discover some new codes or to re-interpret some old codes. In any event, the only limitation of extending the theory of maximal-cube covering is one's own imagination.

5.3 DISCUSSION

It is interesting to note that although the original

purpose of developing the algebraic structure of maximal-cube covering was to manipulate Boolean functions, however it approaches its own mathematical form and structure that warrants the study of its own. Further study in the interest of switching theory or coding theory or just abstract algebra are being pursued by the author of this dissertation; the result reported in this dissertation representing the original goal of the authors research subject.

CHAPTER VISUMMARY, CONCLUSIONS AND RECOMMENDATIONS6.1 Summary

In this dissertation, we have systematically developed the Minimum Maximal-Cube Covering approach to handle switching circuits and its extended application into other areas. The work can be summarized as follows:

- (a) We have looked into the intrinsic characteristics of switching circuits by chronologically surveying previous work and comparing and analyzing various approaches thoroughly. The results have led to a defined set of most desirable characteristics for simplifying and realizing switching functions with the most economical circuits.
- (b) We have developed the algebraic structure to facilitate the Minimum Maximal-Cube Covering approach to switching theory.
- (c) The application of a branching graph to two-level minimization of switching circuits was fully illustrated with detailed examples. Other switching problems were also treated with the branching graph method.
- (d) An alternative method of using algebraic-geometric properties of the Minimum Maximum-Cube Covering was also developed. The application of this method into

switching circuits was illustrated with various examples. Since this method turns out to be better to implement with a computer program, demonstrative computer programming was carried out in detail to verify the validity of the intrinsic merits of computerization. Because of limited resources, we were able to use only the Univac Spectra 70 VMOS at New Jersey Institute of Technology. No attempt was made to optimize programming.

- (e) Exploratory work has been carried into other fields of interest in Chapter Five.

6.2 Recommendations

- (a) Due to lack of resources, we were not able to fully develop the complete set of computer programs for digital system usage in industry and for educational purposes in universities. We would urgently recommend this development using the theory and technique developed in this dissertation.
- (b) The extension of the Minimum Maximal-Cube Covering into other fields of interest was only lightly explored. Further work should also be attempted.

6.3 Conclusions

- (a) The mathematical properties of the Minimum Maximal-Cube Covering, in addition to their original purpose of facilitating handling of switching circuits and

digital systems, are quite interesting and rich in their own right. In the area of Boolean algebra and the area of lattice theory, Minimum Maximal-Cube Covering seems to offer a very fresh, powerful new tool.

- (b) With the ever-increasing availability of microprocessors and LSI's, hardware realized instruments and machines using Minimum Maximal-Covering principle are quite conceivable.

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APPENDIX A COMPUTER PROGRAMS

1. PROGRAM LARGE
2. SUBROUTINE CUBESD
3. PROGRAM DISTBL
4. PROGRAM DIGBIT
5. ASSEMBLY SUBROUTINE NDIST
6. ASSEMBLY SUBROUTINE BICHA

FORTRAN IV (VER 45) SOURCE LISTING:

```

1      PROGRAM LARGE
2      COMMON CHAR(520,4),MT(520),T1(520),T2(520),OUTA(8),NV4,NT,RW
3      %,CUBEN,DIAN,MTH,ZNTH
4      DIMENSION B(9),BINO(9,9),TEMT(1040)
5      EQUIVALENCE (TEMT(1),T1(1))
6      DEFINE FILE 28(520,520,E,RW)
7      IMPLICIT INTEGER (A-Z)
8      1 READ(5,2,END=1000)NV,DONT,ZERO
9      2 FORMAT(3I3)
10     IF(NV.GE.900)GO TO 1000
11     NV4=NV/4.0*0.8
12     CALL NVARIS(NV)
13     CALL BICHA(NV)
14     DO 502 I=1,1040
15     502 TEMT(I)=0
16     DO 4 U=1,65
17     VI=(U-1)*16
18     VA=VI+1
19     VB=VI+16
20     READ(5,3)(TEMT(V),V=VA,VB)
21     3 FORMAT(8I10)
22     DO 4 V=VA,VB
23     IF(TEMT(V).LT.2147483640)GO TO 4
24     NTEM=V
25     GO TO 5
26     4 CONTINUE
27     5 K=1
28     DO 9 J=1,NTEM
29     MT(K)=TEMT(J)
30     IF(ZERO.EQ.0)GO TO 6
31     IF(MT(K).EQ.0)GO TO 9
32     6 J6=K-1
33     IF(K.EQ.1)GO TO 8
34     DO 7 I=1,J6
35     IF(MT(I).EQ.MT(K))GO TO 9
36     7 CONTINUE
37     8 K=K+1
38     9 CONTINUE
39     10 NT=K-2
40     CT=NT-DONT
41     CUBEN=0
42     PRINT 11,NT,CT,DONT,NV
43     11 FORMAT('1TOATAL NUMBER OF ON-SET AND DONT-CARES = ',
44     %,15,/, ' ON-SET = ',15,5X, ' DONT-CARES = ',15
45     %,/, ' NUMBERS OF VARIABLES = ',15)
46     IF(CT.LE.0)GO TO 1
47     RW=1
48     MAXCUB=1
49     DO 12 I=1,10
50     MAXCUB=MAXCUB*2

```

FORTRAN IV (VER 45) SOURCE LISTING: LARGE PROGRAM

```

51      IF(MAXCUB.GT.NT)GO TO 14
52      12 CONTINUE
53      14 MAXCUB=I-1
54      BINO(1,1)=1
55      BINO(2,1)=2
56      BINO(2,2)=1
57      DO 142 I=2,MAXCUB
58      I1=I-1
59      BINO(I,1)=I
60      DO 141 J=2,I1
61      J1=J-1
62      141 BINO(I,J)=BINO(I1,J1)+BINO(I1,J)
63      142 BINO(I,I)=1
64      WRITE(2,145)MAXCUB,((BINO(I,J),J=1,MAXCUB),I=1,MAXCUB)
65      145 FORMAT('1 MAXIMUM CUBE DIMENSION = ',I3,/,,(14I5))
66      PRINT 30,((I,MT(I)),I=1,NT)
67      30 FORMAT(/,2(5X,'MINT.#',I4,' = ',I10))
68      DO 45 I=1,NT
69      CHAR(I,1)=0
70      CHAR(I,2)=0
71      CHAR(I,3)=0
72      CHAR(I,4)=0
73      45 CONTINUE
74      ZEN=0
75      DO 100 I2=1,NT
76      DO 51 I=1,9
77      51 B(I)=0
78      DO 60 I1=1,NT
79      IF(I1.NE.I2)GO TO 56
80      MTD=0
81      GO TO 60
82      56 CALL NDIST(MTD,MT(I1),MT(I2))
83      IF(MTD.LE.MAXCUB)B(MTD)=B(MTD)+1
84      IF(MTD.LE.9)GO TO 60
85      MTD=9
86      60 T1(I1)=MTD
87      WRITE(28'RW,61)(T1(I),I=1,NT)
88      61 FORMAT(200I1,200I1,120I1)
89      IF(B(1).EQ.0)GO TO 85
90      B1=MAXCUB
91      B2=B1+1
92      DO 80 J11=1,B1
93      J1=B2-J11
94      DO 70 J2=1,J1
95      IF(B(J2).LT.BINO(J1,J2))GO TO 80
96      70 CONTINUE
97      CHAR(I2,1)=J1
98      IF(B(1).NE.BINO(J1,1))GO TO 72
99      CHAR(I2,2)=1
100     GO TO 100

```

FORTRAN IV (VER 45) SOURCE LISTING: LARGE PROGRAM

```

101 72 CHAR(I2,2)=5
102 GO TO 100
103 80 CONTINUE
104 85 MINT=MT(I2)
105 PRINT 90,I2,MINT
106 90 FORMAT('0MINTERM #',I4,5X,'=',I10,' IS ISOLATED AND ESSEN'
107 100 CONTINUE
108 PRINT 98
109 98 FORMAT('1THE CUBE COVERING LIST :',/)
110 WRITE(2,101){(CHAR(I,J),J=1,2),I=1,NT)
111 101 FORMAT('0THE CHARACTERISTICS :',/,(5X,22I3))
112 CHECK=0
113 IF(CT,GE,NT)GO TO 102
114 DO 103 I=1,DONT
115 X=CT+I
116 CHAR(X,3)=9999
117 CHAR(X,2)=109
118 103 CONTINUE
119 102 MAX=2*MAXCUB
120 KOUNT=0
121 DO 105 I=1,NT
122 IF(CHAR(I,1).NE.MAXCUB)GO TO 105
123 KOUNT=KOUNT+1
124 105 CONTINUE
125 IF(KOUNT,GE,MAX)GO TO 107
126 DO 106 I=1,NT
127 IF(CHAR(I,1).NE.MAXCUB)GO TO 106
128 CHAR(I,1)=CHAR(I,1)-1
129 CHAR(I,2)=5
130 106 CONTINUE
131 MAXCUB=MAXCUB-1
132 WRITE(2,108)MAXCUB
133 108 FORMAT('0MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO:',I10)
134 GO TO 102
135 107 CONTINUE
136 IF(ZEN.EQ.1)GO TO 125
137 DO 120 I6=1,NT
138 IF(CHAR(I6,2).NE.1)GO TO 120
139 D=CHAR(I6,1)
140 MTH=I6
141 ZNTH=1
142 CALL CUBESD(ZNTH,MTH,D)
143 IF(DIAN,NE.0)GO TO 119
144 CHAR(MTH,1)=D-1
145 CHAR(MTH,2)=5
146 IF(D.NE.MAXCUB)GO TO 120
147 IF(CHECK.NE.0)GO TO 120
148 GO TO 102
149 119 IF(D.NE.MAXCUB)GO TO 120
150 IF(CHECK.NE.0)GO TO 120

```

FORTRAN IV (VER 45) SOURCE LISTING: LARGE PROGRAM

```

151      CHECK=1
152      120 CONTINUE
153      ZEN=1
154      WRITE(2,122)MAXCUB
155      122 FORMAT('0ALL CORE CUBES HAVE BEEN SEARCHED, '
156      %,/, ' AND THE MAXIMUM CUBE DIMENSION IS :',I5)
157      125 B1=MAXCUB
158      B2=B1+1
159      DO 280 J1=1,B1
160      140 J1=B2-J11
161      DO 180 J2=1,NT
162      IF(CHAR(J2,1).NE.J1)GO TO 180
163      IF(CHAR(J2,2).NE.5)GO TO 180
164      PSHARE=CHAR(J2,3)
165      CLOSE=CHAR(J2,4)
166      MTH=J2
167      D=J1
168      DO 170 I6=1,NT
169      IF(I6.EQ.J2)GO TO 170
170      IF(CHAR(I6,1).NE.J1)GO TO 170
171      IF(CHAR(I6,2).NE.5)GO TO 170
172      IF(CHAR(I6,3).GT.PSHARE)GO TO 170
173      IF(CHAR(I6,3).EQ.PSHARE)GO TO 150
174      160 PSHARE=CHAR(I6,3)
175      MTH=I6
176      GO TO 170
177      150 CONTINUE
178      IF(CHAR(I6,4).GT.CLOSE)GO TO 170
179      CLOSE=CHAR(I6,4)
180      GO TO 160
181      170 CONTINUE
182      ZNTH=5
183      CALL CUBESD(ZNTH,MTH,D)
184      IF(DIAN.NE.0)GO TO 175
185      CHAR(MTH,1)=D=1
186      IF(CHECK.NE.0)GO TO 180
187      IF(D.NE.MAXCUB)GO TO 180
188      CHECK=CHECK+1
189      GO TO 102
190      175 IF(D.NE.MAXCUB)GO TO 180
191      IF(CHECK.NE.0)GO TO 180
192      CHECK=1
193      180 CONTINUE
194      DO 280 J6=1,NT
195      IF(CHAR(J6,1).NE.J1)GO TO 280
196      IF(CHAR(J6,2).EQ.5)GO TO 140
197      280 CONTINUE
198      PRINT 300,(CHAR(I,2),I=1,NT)
199      300 FORMAT('0THE FINAL CHECKING ',/, (10X,12I5))
200      GO TO 1

```

FORTRAN IV (VER 45) SOURCE LISTING: LARGE PROGRAM

```
201 1000 STOP  
202      END
```

FORTRAN IV (VER 45) SOURCE LISTING: CUBESD SUBROUTINE 12/16/77 20:20

```

1      SUBROUTINE CUBESD(ZNTH,MTH,D)
2      IMPLICIT INTEGER (A-Z)
3      COMMON CHAR(520,4),MT(520),T1(520),T2(520),OUTA(8),NV4,NT,RW
4      %,CUBEN,DIAN,MTH,ZNTH
5      DIMENSION T(520),S(520)
6      DEFINE FILE 28(520,520,E,RW)
7      CHK=2**D
8      RW=MTH
9      READ(28,RW,125)(T1(I),I=1,NT)
10     125 FORMAT(200I1,200I1,120I1)
11     MINT=MT(MTH)
12     IF(D.NE.9)GO TO 140
13     DO 130 I=1,NT
14     IF(T1(I).LT.9)GO TO 130
15     CALL NDIS(T(MT),MINT,MT(I))
16     T1(I)=MT
17     130 CONTINUE
18     140 IR=MTH
19     DIAN=0
20     RAWMT=0
21     S3=NT+1
22     CHAR(S3,4)=9000
23     T(1)=IR
24     DO 200 J3=1,NT
25     IF(T1(J3).NE.D)GO TO 200
26     IF(CHAR(J3,1).LT.D)GO TO 200
27     RW=J3
28     READ(28,RW,125)(T2(I),I=1,NT)
29     T(2)=J3
30     IT=2
31     DO 160 J4=1,NT
32     IF(T1(J4).GE.D)GO TO 160
33     IF(T2(J4).GE.D)GO TO 160
34     LD=T1(J4)+T2(J4)
35     IF(LD.NE.D)GO TO 160
36     IT=IT+1
37     T(IT)=J4
38     160 CONTINUE
39     IF(CHK.NE.IT)GO TO 200
40     WRITE(2,162)MT(IR),MT(J3)
41     162 FORMAT('0CUBE (' ,I10,' ,',I10,' ) HAS BEEN FOUND.')
```

FORTRAN IV (VER 45) SOURCE LISTING: CUBESD SUBROUTINE

```

51      COUNT=COUNT+1
52      170 CONTINUE
53      IF(COUNT.LT.RAWMT)GO TO 200
54      IF(COUNT.EQ.RAWMT)GO TO 175
55      172 S3=J3
56      RAWMT=COUNT
57      DO 180 I=1,IS
58      S(I)=T(I)
59      180 CONTINUE
60      GO TO 200
61      175 CONTINUE
62      WRITE(2,173)MT(S3),MT(J3)
63      173 FORMAT('O WATCH FOR ',I10,' AND ',I10)
64      IF(CHAR(J3,4).GT.CHAR(S3,4))GO TO 200
65      GO TO 172
66      200 CONTINUE
67      IF(DIAN.EQ.0)GO TO 700
68      205 RW=S3
69      READ(28'RW,125)(T2(I),I=1,NT)
70      DO 210 I=1,IS
71      TF=S(I)
72      CHAR(TF,2)=CHAR(TF,2)+100
73      CHAR(TF,3)=9999
74      210 CONTINUE
75      DO 220 I=1,NT
76      IF(CHAR(I,3).EQ.9999)GO TO 220
77      WAB=T1(I)+T2(I)-D
78      CHAR(I,4)=WAB+CHAR(I,4)
79      DA=CHAR(I,1)
80      IF(DA.LT.WAB)GO TO 220
81      G=DA-WAB
82      G1=G
83      IF(G.GE.D)G1=D
84      SHARE=2**G1
85      CHAR(I,3)=CHAR(I,3)+SHARE
86      220 CONTINUE
87      WRITE(2,212)(CHAR(I,3),I=1,NT)
88      212 FORMAT('O CHECKING POINTS ',/, (10X,12I5))
89      WRITE(2,216)(CHAR(I,4),I=1,NT)
90      216 FORMAT('O ACCUMULATED AWAYNESS ',/, (10X,12I5))
91      CUBEN=CUBEN+1
92      PRINT 228,CUBEN
93      228 FORMAT(///,'O THE ',IS,'-TH CUBE COVERING IS :')
94      CALL BICHA2(MT(IR),MT(S3),OUTA)
95      PRINT 230,D,MT(IR),MT(S3),(OUTA(0),O=1,NV4)
96      230 FORMAT(5X,I4,'-CUBE = (',I10,', ',I10,' ) =',5X,B4)
97      WRITE(2,222)MT(IR),MT(S3)
98      222 FORMAT('O CUBE (',I10,', ',I10,' ) IS SELECTED.')
99      DO 232 I=1,IS
100     STEM=S(I)

```

FORTRAN IV (VER 45) SOURCE LISTING: CUBESD SUBROUTINE

```
101      S(I)=MT(STEM)
102  232  CONTINUE
103      PRINT 235,(S(I),I=1,IS)
104  235  FORMAT('0THIS CUBE COVERS THESE MINTERMS : ',/,(5X,6I10))
105      IF(ZNTH.EQ.5)GO TO 700
106      IF(ZNTH.NE.1)GO TO 680
107      PRINT 240,IR,MT(IR)
108  240  FORMAT('0THIS MINTERM MT(',I4,' ) = ',I10,5X,' IS A CORE PO
109      GO TO 700
110  300  DO 320 I=1,IS
111      S(I)=T(I)
112  320  CONTINUE
113      GO TO 205
114  680  PRINT 685
115  685  FORMAT('0THERE IS AN ERROR IN THE PROGRAM!!!')
116  700  RETURN
117      END
```


FORTRAN IV (VER 45) SOURCE LISTING:

```

1      PROGRAM DISTBL
2      COMMON MT(520),T1(520),T2(520),OUTA(8),NV4,NT,RW
3      %,CUBEN,DIAN,MTH,ZNTH
4      DIMENSION B(9),TEMT(1040)
5      EQUIVALENCE (TEMT(1),T1(1))
6      IMPLICIT INTEGER (A-Z)
7      1 READ(5,2,END=1000)NV,DONT,ZERO
8      2 FORMAT(3I3)
9      IF(NV.GE.900)GO TO 1000
10     NV4=NV/4.0+0.8
11     CALL NVARIS(NV)
12     CALL BYCHA(NV)
13     DO 502 I=1,1040
14     502 TEMT(I)=0
15     DO 4 U=1,65
16     VI=(U-1)*16
17     VA=VI+1
18     VB=VI+16
19     READ(5,3)(TEMT(V),V=VA,VB)
20     3 FORMAT(16I5)
21     DO 4 V=VA,VB
22     IF(TEMT(V).LT.90000,GO TO 4
23     NTEM=V
24     GO TO 5
25     4 CONTINUE
26     5 K=1
27     DO 9 J=1,NTEM
28     MT(K)=TEMT(J)
29     IF(ZERO.EQ.0)GO TO 6
30     IF(MT(K).EQ.0)GO TO 9
31     6 J6=K-1
32     IF(K.EQ.1)GO TO 8
33     DO 7 I=1,J6
34     IF(MT(I).EQ.MT(K))GO TO 9
35     7 CONTINUE
36     8 K=K+1
37     9 CONTINUE
38     10 NT=K-2
39     CT=NT-DONT
40     CUBEN=0
41     PRINT 11,NT,CT,DONT,NV
42     11 FORMAT('1TOTAL NUMBER OF ON-SET AND DONT-CARES = '
43     %,15,/, ' ON-SET = ',15,5X, ' DONT-CARES = ',15
44     %,/, ' NUMBERS OF VARIABLES = ',15)
45     IF(CT.LE.0)GO TO 1
46     DO 100 I2=1,NT
47     PRINT 51,I2,MT(I2)
48     51 FORMAT('0THE DISTANCE FUNCTION FOR MINTERM #',I5,' = ',I5
49     %, ' IS :',//)
50     DO 52 I=1,9

```

FORTRAN IV (VER 45) SOURCE LISTING: DISTBL PROGRAM

```
51      52 B(I)=0
52      DO 60 I1=1,NT
53      IF(I1.NE.I2)GO TO 56
54      MTD=0
55      GO TO 60
56      56 CALL NDIST(MTD,MT(I1),MT(I2))
57      B(MTD)=B(MTD)+1
58      IF(MTD.LE.9)GO TO 60
59      MTD=9
60      60 T1(I1)=MTD
61      PRINT 61,((I,T1(I)),I=I2,NT)
62      61 FORMAT(5(2X,'T(',I4,') = ',I1),/)
63      PRINT 62,(B(I),I=1,9)
64      62 FORMAT(//,' THE DISTANCE DISTRIBUTION ARE :',/,'915',/)//)
65      100 CONTINUE
66      GO TO 1
67      1000 STOP
68      END
```

FORTRAN IV (VER 45) SOURCE LISTING:

```

1      PROGRAM DIGBIT
2      COMMON CHAR(520,4),MT(520),T1(520),T2(520),OUTA(8),NV4,NT,PW
3      %,CUBEN,DIAN,MTH,ZNTH
4      DIMENSION B(9),BIND(9,9),TEMT(1040)
5      EQUIVALENCE (TEMT(1),T1(1))
6      IMPLICIT INTEGER (A-Z)
7      1 READ(5,2,END=1000)NV,DONT,ZERO
8      2 FORMAT(3I3)
9      IF(NV.GE.900)GO TO 1000
10     NV4=NV/4.0+0.8
11     CALL NVARIS(NV)
12     CALL BICHA(NV)
13     DO 502 I=1,1040
14     502 TEMT(I)=0
15     DO 4 U=1,65
16     VI=(U-1)*16
17     VA=VI+1
18     VB=VI+16
19     READ(5,3)(TEMT(V),V=VA,VB)
20     3 FORMAT(16I5)
21     DO 4 V=VA,VB
22     IF(TEMT(V).LT.90000)GO TO 4
23     NTEM=V
24     GO TO 5
25     4 CONTINUE
26     5 K=1
27     DO 9 J=1,NTEM
28     1 MT(K)=TEMT(J)
29     IF(ZERO.EQ.0)GO TO 6
30     IF(MT(K).EQ.0)GO TO 9
31     6 J6=K-1
32     IF(K.EQ.1)GO TO 8
33     DO 7 I=1,J6
34     IF(MT(I).EQ.MT(K))GO TO 9
35     7 CONTINUE
36     8 K=K+1
37     9 CONTINUE
38     10 NT=K-2
39     CT=NT-DONT
40     CUBEN=0
41     PRINT 11,NT,CT,DONT,NV
42     11 FORMAT('TOTAL NUMBER OF ON-SET AND DONT-CARES = '
43     %,15,/, ' ON-SET = ',15,5X, ' DONT-CARES = ',15
44     %,/, ' NUMBERS OF VARIABLES = ',15)
45     IF(CT.LE.0)GO TO 1
46     DO 21 I=1,NT
47     CALL BICHA1(MT(I),OUTA)
48     PRINT 20,I,MT(I),(OUTA(J),J=1,NV4)
49     20 FORMAT(5X,'MINTERM #',14,' = ',14,' IS ',5X,8A4)
50     21 CONTINUE

```

FORTRAN IV (VER 45) SOURCE LISTING: DIGBIT PROGRAM

```
51      GO TO 1
52 1000 STOP
53      END
```

TITLE 'NDIST--A BINARY DISTANCE FINDING FUNCTION--OCT. 75'

*

* FOR FORTRAN PROGRAM, FOLLOW THE FOLLOWING TWO STEPS:

* 1. GIVING THE # OF VARIABLES, N, BY CALL NVARIS(N)

* (NOTE: THIS N WILL BE USED UNTIL NEXT CALL NVARIS. DEFAULT N=3;

* 2. FINDING THE DISTANCE OF 2 VECTORS, I, J, BY FUNCTION NDIST(I, J)

*

NVARIS	SPACE	3
	CSECT	
	USING	NVARIS,15
	L	1,0(0,1)
	L	1,0(0,1)
	ST	1,NV
	LNR	1,1
	A	1,N32
	STH	1,PRESHF+2
	MVI	12(13),X'FF'
	BR	14
	SPACE	3
NDIST	ENTRY	NDIST
	USING	NDIST,15
	STM	1,5,24(13)
	LM	3,4,4(1)
	L	5,0(0,1)
	L	2,0(0,3)
	L	3,0(0,4)
PRESHF,	XR	3,2
	SLL	3,0
	SLR	1,1
	L	4,NV
LP	SLR	2,2
	SLDL	2,1
	AR	1,2
	BCT	4,LP
	ST	1,0(0,5)
	LM	1,5,24(13)
	MVI	12(13),X'FF'
	BR	14
	SPACE	3
N32	DC	F'32'
NV	DC	F'32'
	END	
/EDN		

BICHA	TITLE	CSECT	USING	
	'BICHA -- BITS TO CHARACTERS (FOR OUTPUT)'		BICHA,15	
		L	1,0(0,1)	
		L	1,0(0,1)	
		ST	1,NV	
		STC	1,SETN1+1	
		STC	1,SETN2+1	
		LNR	1,1	
		A	1,N32	
		STH	1,PRESHF1+2	
		STH	1,PRESHF2+2	
		STH	1,PRESHF3+2	
		MVI	12(13),X'FF'	
		BR	14	
		SPACE	3	
		ENTRY	BICHA1	
BICHA1		USING	BICHA1,15	
		STM	1,4,24(13)	
		L	3,0(0,1)	
		L	3,0(0,3)	
		L	1,4(0,1)	
SETN1		MVC	0(32,1),BIANK	
		L	4,NV	
PRESHF1		SLL	3,0	
LP1		SLR	2,2	
		SLDL	2,1	
		AH	2,ZERO	
		STC	2,0(0,1)	
		A	1,ONE	
		BCT	4,LP1	
		MVI	12(13),X'FF'	
		LM	1,4,24(13)	
		BR	14	
		SPACE	3	
		ENTRY	BICHA2	
BICHA2		USING	BICHA2,15	
		STM	1,6,24(13)	
		LM	3,4,0(1)	
		L	3,0(0,3)	
		L	5,0(0,4)	
		L	1,8(0,1)	
		L	6,NV	
SETN2		MVC	0(32,1),BIANK	
		XR	5,3	
PRESHF2		SLL	3,0	
PRESHF3		SLL	5,0	
LP2		SLR	2,2	
		SLDL	2,1	
		SLR	4,4	
		SLDL	4,1	
		ALR	4,4	
		BC	5,DASH	
		AH	2,ZERO	

	STC	2,0(0,1)
	B	ADDR
DASH	MVI	0(1),C,-
ADDR	A	1,ONE
	BCT	6,LP2
	LM	1,6,24(13)
	MVI	12(13),X'FF'
	BR	14
ZERO	DC	X'00F0'
ONE	DC	F'1'
NV	DC	F'32'
N32	DC	F'32'
BLANK	DC	CL32'
	END	

APPENDIX B

EXAMPLE OF 4 VARIABLES

TOATAL NUMBER OF ON-SET AND DONT*CARES = 8
 ON-SET = 8 DONT-CARES = 0
 NUMBERS OF VARIABLES = 4

MINT.#	1	=	0	MINT.#	2	=	1	MINT.#	3	=	5
MINT.#	4	=	7	MINT.#	5	=	15	MINT.#	6	=	14
MINT.#	7	=	10	MINT.#	8	=	8	MINT.#		=	

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 1-CUBE = (8 , 10) = 10-0

THIS CUBE COVERS THESE MINTERMS :
 8 10

THE 2-TH CUBE COVERING IS :
 1-CUBE = (14 , 15) = 111-

THIS CUBE COVERS THESE MINTERMS :
 14 15

THE 3-TH CUBE COVERING IS :
 1-CUBE = (7 , 5) = 01-1

THIS CUBE COVERS THESE MINTERMS :
 7 5

THE 4-TH CUBE COVERING IS :
 1-CUBE = (1 , 0) = 000-

THIS CUBE COVERS THESE MINTERMS :
 1 0

THE FINAL CHECKING
 105 105 105 105 105 105 105 105

TOATAL NUMBER OF ON-SET AND DONT*CARES = 9
 ON-SET = 9 DONT-CARES = 0
 NUMBERS OF VARIABLES = 4

MINT.#	1	=	0	MINT.#	2	=	1	MINT.#	3	=	4
MINT.#	4	=	5	MINT.#	5	=	7	MINT.#	6	=	8
MINT.#	7	=	10	MINT.#	8	=	14	MINT.#	9	=	15

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 1-CUBE = (1 , 4) = 0-0-

THIS CUBE COVERS THESE MINTERMS :
 1 4 0 5

THIS MINTERM MT(2) = 1 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
 1-CUBE = (10 , 8) = 10-0

THIS CUBE COVERS THESE MINTERMS :
 10 8

THE 3-TH CUBE COVERING IS :
 1-CUBE = (14 , 15) = 111-

THIS CUBE COVERS THESE MINTERMS ;
 14 15

THE 4-TH CUBE COVERING IS :
 1-CUBE = (7 , 5) = 01-1

THIS CUBE COVERS THESE MINTERMS ;
 7 5

THE FINAL CHECKING
 105 101 101 205 105 105 105 105 105

TOATAL NUMBER OF ON-SET AND DONT-CARES = 14
 ON-SET = 14 DONT-CARES = 0
 NUMBERS OF VARIABLES = 4

MINT.#	1	=	1	MINT.#	2	=	2	MINT.#	3	=	3
MINT.#	4	=	4	MINT.#	5	=	5	MINT.#	6	=	6
MINT.#	7	=	7	MINT.#	8	=	8	MINT.#	9	=	9
MINT.#	10	=	10	MINT.#	11	=	11	MINT.#	12	=	12
MINT.#	13	=	13	MINT.#	14	=	14	MINT.#		=	

THE CUBE COVERING LIST :

[Empty box]

THE 1-TH CUBE COVERING IS :
 2-CUBE = (14 , 8) = 1--0

THIS CUBE COVERS THESE MINTERMS :
 14 8 10 12

THE 2-TH CUBE COVERING IS :
 2-CUBE = (1 , 13) = --01

THIS CUBE COVERS THESE MINTERMS :
 1 13 5 9

[Empty box]

THE 3-TH CUBE COVERING IS :
 2-CUBE = (7 , 2) = 0-1-

THIS CUBE COVERS THESE MINTERMS :
 7 2 3 6

THE 4-TH CUBE COVERING IS :
 2-CUBE = (11 , 8) = 10--

THIS CUBE COVERS THESE MINTERMS :
 11 8 9 10

[Empty box]

THE 5-TH CUBE COVERING IS :
2-CUBE = (4 , 14) = -1-0

THIS CUBE COVERS THESE MINTERMS :
4 14 6 12

THE FINAL CHECKING
105 105 105 105 105 205 105 205 205 205 105 20
105 205

TOTAL NUMBER OF ON-SET AND DONT-CARES = 9
ON-SET = 9 DONT-CARES = 0
NUMBERS OF VARIABLES = 4

MINT.#	1	=	1	MINT.#	2	=	2	MINT.#	3	=	3
MINT.#	4	=	5	MINT.#	5	=	15	MINT.#	6	=	14
MINT.#	7	=	8	MINT.#	8	=	9	MINT.#	9	=	11

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
1-CUBE = (2 , 3) = 001-

THIS CUBE COVERS THESE MINTERMS :
2 3

THIS MINTERM MT(2) = 2 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
1-CUBE = (5 , 1) = 0-01

THIS CUBE COVERS THESE MINTERMS :
5 1

THIS MINTERM MT(4) = 5 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
 1-CUBE = (14 , 15) = 111-

THIS CUBE COVERS THESE MINTERMS :
 14 15

THIS MINTERM MT(6) = 14 IS A CORE POINT

THE 4-TH CUBE COVERING IS :
 1-CUBE = (8 , 9) = 100-

THIS CUBE COVERS THESE MINTERMS :
 8 9

THIS MINTERM MT(7) = 8 IS A CORE POINT

THE 5-TH CUBE COVERING IS :
 2-CUBE = (11 , 1) = -0-1

THIS CUBE COVERS THESE MINTERMS :
 11 1 3 9

THE FINAL CHECKING ✓
 205 101 205 101 105 101 101 205 105

TOATAL NUMBER OF ON-SET AND DONT#CARES = 12
 ON-SET = 12 DONT-CARES = 0
 NUMBERS OF VARIABLES = 5

MINT.# 1 = 0	MINT.# 2 = 1	MINT.# 3 =
MINT.# 4 = 5	MINT.# 5 = 6	MINT.# 6 =
MINT.# 7 = 24	MINT.# 8 = 25	MINT.# 9 = 2
MINT.# 10 = 28	MINT.# 11 = 30	MINT.# 12 = 3

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 1-CUBE = (31 , 30) = 1111-

THIS CUBE COVERS THESE MINTERMS :
 31 30

THE 2-TH CUBE COVERING IS :
 1-CUBE = (28 , 24) = 11-00

THIS CUBE COVERS THESE MINTERMS :
 28 24

THE 3-TH CUBE COVERING IS :
 1-CUBE = (27 , 25) = 110-1

THIS CUBE COVERS THESE MINTERMS :
 27 25

THE 4-TH CUBE COVERING IS :
 1-CUBE = (7 , 6) = 0011-

THIS CUBE COVERS THESE MINTERMS :
 7 6

THE 5-TH CUBE COVERING IS :
 1-CUBE = (5 , 1) = 00-01

THIS CUBE COVERS THESE MINTERMS :
 5 1

THE 6-TH CUBE COVERING IS :
 1-CUBE = (0 , 2) = 000-0

THIS CUBE COVERS THESE MINTERMS :
 0 2

THE FINAL CHECKING
 105 105 105 105 105 105 105 105 105 105 105 105 10

TOTAL NUMBER OF ON-SET AND DONT-CARES = 16
 ON-SET = 4 DONT-CARES = 12
 NUMBERS OF VARIABLES = 5

MINT.#	1	=	8	MINT.#	2	=	14	MINT.#	3	=	16
MINT.#	4	=	19	MINT.#	5	=	20	MINT.#	6	=	21
MINT.#	7	=	22	MINT.#	8	=	23	MINT.#	9	=	24
MINT.#	10	=	25	MINT.#	11	=	26	MINT.#	12	=	27
MINT.#	13	=	28	MINT.#	14	=	29	MINT.#	15	=	30
MINT.#	16	=	31	MINT.#		=				=	

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 1-CUBE = (8 , 24) = -1000

THIS CUBE COVERS THESE MINTERMS :
 8 24

THIS MINTERM MT(1) = 8 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
 1-CUBE = (14 , 30) = -1110

THIS CUBE COVERS THESE MINTERMS :
 14 30

THIS MINTERM MT(2) = 14 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
 2-CUBE = (16 , 28) = 1--00

THIS CUBE COVERS THESE MINTERMS :
 16 28 20 24

THIS MINTERM MT(3) = 16 IS A CORE POINT

THE 4-TH CUBE COVERING IS :
 2-CUBE = (19 , 31) = 1--11

THIS CUBE COVERS THESE MINTERMS :
 19 31 23 27

THIS MINTERM MT(4) = 19 IS A CORE POINT

THE FINAL CHECKING

101 101 101 101 209 109 109 209 309 109 109 20
 209 109 209 200

TOATAL NUMBER OF ON-SET AND DONT-CARES = 17
 ON-SET = 10 DONT-CARES = 7
 NUMBERS OF VARIABLES = 5

MINT.#	1	=	0	MINT.#	2	=	3	MINT.#	3	=	4
MINT.#	4	=	6	MINT.#	5	=	7	MINT.#	6	=	15
MINT.#	7	=	21	MINT.#	8	=	23	MINT.#	9	=	26
MINT.#	10	=	28	MINT.#	11	=	2	MINT.#	12	=	8
MINT.#	13	=	12	MINT.#	14	=	14	MINT.#	15	=	24
MINT.#	16	=	31	MINT.#	17	=	17	MINT.#		=	

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 2-CUBE = (3 , 6) = 00-1-

THIS CUBE COVERS THESE MINTERMS ;
 3 6 7 2

THIS MINTERM MT(2) = 3 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
 1-CUBE = (26 , 24) = 110-0

THIS CUBE COVERS THESE MINTERMS :
 26 24

THIS MINTERM MT(9) = 26 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
 2-CUBE = (28 , 8) = -1-00

THIS CUBE COVERS THESE MINTERMS :
 28 8 12 24

THIS MINTERM MT(10) = 28 IS A CORE POINT

THE 4-TH CUBE COVERING IS :
 2-CUBE = (0 , 6) = 00--0

THIS CUBE COVERS THESE MINTERMS :
 0 6 4 2

THE 5-TH CUBE COVERING IS :
 2-CUBE = (15 , 6) = 0-11-

THIS CUBE COVERS THESE MINTERMS :
 15 6 7 14

THE 6-TH CUBE COVERING IS :
 2-CUBE = (23 , 15) = --111

THIS CUBE COVERS THESE MINTERMS :
 23 15 7 31

THE 7-TH CUBE COVERING IS :
 1-CUBE = (21 , 17) = 10-01

THIS CUBE COVERS THESE MINTERMS :
 21 17

THE FINAL CHECKING

105	101	105	205	205	205	105	105	101	101	105	101
105	105	205	209	209							

TOATAL NUMBER OF ON-SET AND DONT*CARES = 14
 ON-SET = 14 DONT-CARES = 0
 NUMBERS OF VARIABLES = 5

MINT.#	1	=	0	MINT.#	2	=	2	MINT.#	3	=	3
MINT.#	4	=	4	MINT.#	5	=	8	MINT.#	6	=	10
MINT.#	7	=	11	MINT.#	8	=	12	MINT.#	9	=	16
MINT.#	10	=	20	MINT.#	11	=	24	MINT.#	12	=	26
MINT.#	13	=	27	MINT.#	14	=	28	MINT.#		=	

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 2-CUBE = (3 , 10) = 0-01-

THIS CUBE COVERS THESE MINTERMS :
 3 10 2 11

THIS MINTERM MT(3) = 3 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
 3-CUBE = (4 , 24) = ---00

THIS CUBE COVERS THESE MINTERMS :
 4 24 0 8 12 16 20 28

THIS MINTERM MT(4) = 4 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
 2-CUBE = (27 , 10) = -101-

THIS CUBE COVERS THESE MINTERMS :
 27 10 11 26

THIS MINTERM MT(13) = 27 IS A CORE POINT

THE FINAL CHECKING

105 105 101 101 105 205 201 101 101 101 105 10
 101 101

MAXIMUM CUBE DIMENSION = 3

1 0 0 2 1 0 3 3 1

THE CHARACTERISTICS :

2 1 2 1 2 1 2 1 2 1 2 1 2 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 1

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 1

CUBE (8 , 0) HAS BEEN FOUND.

CUBE (8 , 10) HAS BEEN FOUND.

WATCH FOR 0 AND 10

CHECKING POINTS

0 0 0 0 0 0 9999 9999

ACCUMULATED AWAYNESS :

2 4 6 6 4 2 0 0

CUBE (18 , 10) IS SELECTED.

CUBE (14 , 15) HAS BEEN FOUND.

CUBE (14 , 10) HAS BEEN FOUND.

CHECKING POINTS

0 0 0 0 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

8 10 10 8 4 2 0 0

CUBE (14 , 15) IS SELECTED.

CUBE (7 , 5) HAS BEEN FOUND.

CUBE (7 , 15) HAS BEEN FOUND.

CHECKING POINTS

0 0 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

12 12 10 8 4 2 0 0

CUBE (7 , 5) IS SELECTED.

CUBE (1 , 0) HAS BEEN FOUND.

CUBE (1 , 5) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

12 12 10 8 4 2 0 0

CUBE (1 , 0) IS SELECTED.

MAXIMUM CUBE DIMENSION = 3

1 0 0 2 1 0 3 3 1

THE CHARACTERISTICS :

2 5 2 1 2 1 2 5 2 1 2 1 2 1 2 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

CUBE (1 , 4) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 1 1 0 0 0

ACCUMULATED AWAYNESS :

0 0 0 0 2 2 4 4 4

CUBE (1 , 4) IS SELECTED.

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 2

CUBE (10 , 8) HAS BEEN FOUND.

CUBE (10 , 14) HAS BEEN FOUND.

WATCH FOR 8 AND 14

CHECKING POINTS

9999 9999 9999 9999 1 9999 9999 0 0

ACCUMULATED AWAYNESS :

0 0 0 0 8 2 4 6 8

CUBE (10 , 8) IS SELECTED.

CUBE (14 , 10) HAS BEEN FOUND.

CUBE (14 , 15) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 1 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

0 0 0 0 10 2 4 6 8

CUBE (14 , 15) IS SELECTED.

CUBE (7 , 5) HAS BEEN FOUND.

CUBE (7 , 15) HAS BEEN FOUND.

WATCH FOR 5 AND 15

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

0 0 0 0 10 2 4 6 8

CUBE (7 , 5) IS SELECTED.

MAXIMUM CUBE DIMENSION = 3

1 0 2 1 0 3 3 1

THE CHARACTERISTICS :

3 1 3 1 3 5 3 1 3 5 3 5 3 1 3 1 3 5 3 5 3 1
3 5 3 1 3 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 2

CUBE (14 , 2) HAS BEEN FOUND.

CUBE (14 , 4) HAS BEEN FOUND.

WATCH FOR 2 AND 4

CUBE (14 , 8) HAS BEEN FOUND.

WATCH FOR 4 AND 8

CHECKING POINTS

0 1 0 1 0 1 0 9999 1 9999 1 9999
1 9999

ACCUMULATED AWAYNESS :

4 2 4 2 4 2 4 0 2 0 2 0
2 0

CUBE (14 , 8) IS SELECTED.

CUBE (1 , 7) HAS BEEN FOUND.

CUBE (1 , 11) HAS BEEN FOUND.

WATCH FOR 7 AND 11

CUBE (1 , 13) HAS BEEN FOUND.

WATCH FOR 11 AND 13

CUBE (1 , 13) IS SELECTED.

CUBE (7 , 1) HAS BEEN FOUND.

CUBE (7 , 2) HAS BEEN FOUND.

CUBE (7 , 4) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 3 9999 9999 9999 9999 9999 9999 3 9999
9999 9999

ACCUMULATED AWAYNESS :

4 6 6 6 4 6 6 0 2 0 6 0
2 0

CUBE (7 , 2) IS SELECTED.

CUBE (11 , 1) HAS BEEN FOUND.

CUBE (11 , 2) HAS BEEN FOUND.

WATCH FOR 1 AND 2

CUBE (11 , 8) HAS BEEN FOUND.

WATCH FOR 1 AND 8

CHECKING POINTS

9999 9999 9999 3 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999

ACCUMULATED AWAYNESS :

4 6 6 10 4 6 6 0 2 0 6 0
2 0

CUBE (11 , 8) IS SELECTED.

CUBE (4 , 7) HAS BEEN FOUND.

CUBE (4 , 13) HAS BEEN FOUND.

WATCH FOR 7 AND 13

CUBE (4 , 14) HAS BEEN FOUND.

WATCH FOR 13 AND 14

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999

ACCUMULATED AWAYNESS :

4 6 6 10 4 6 6 0 2 0 6 0
2 0

CUBE (4 , 14) IS SELECTED.

MAXIMUM CUBE DIMENSION = 3

1 3 2 1 3 3 1

THE CHARACTERISTICS :

3 1 1 1 3 1 1 1 2 1 1 1 1 1 3 1 3 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

CUBE (2 , 3) HAS BEEN FOUND.

CHECKING POINTS

1 9999 9999 0 0 0 0 0 0 1

ACCUMULATED AWAYNESS :

2 0 0 4 4 4 4 4 2

CUBE (2 , 3) IS SELECTED.

CUBE (5 , 1) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 0 0 0 1 1

ACCUMULATED AWAYNESS :

2 0 0 4 8 10 8 6 6

CUBE (5 , 1) IS SELECTED.

CHECKING POINTS

9999 1 1 2 9999 1 1 9999 9999 9999 2 9999
 9999 9999

ACCUMULATED AWAYNESS :

4 6 6 4 4 6 6 0 2 0 4 n
 2 0

CUBE (1 , 13) IS SELECTED,

CUBE (7 , 1) HAS BEEN FOUND,

CUBE (7 , 2) HAS BEEN FOUND,

CUBE (7 , 4) HAS BEEN FOUND,

CHECKING POINTS

9999 9999 9999 3 9999 9999 9999 9999 9999 9999 3 9999
 9999 9999

ACCUMULATED AWAYNESS :

4 6 6 6 4 6 6 0 2 0 6 n
 2 0

CUBE (7 , 2) IS SELECTED,

CUBE (11 , 1) HAS BEEN FOUND,

CUBE (11 , 2) HAS BEEN FOUND,

WATCH FOR 1 AND 2

CUBE (11 , 6) HAS BEEN FOUND,

WATCH FOR 1 AND 8

CHECKING POINTS

9999 9999 9999 3 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999

ACCUMULATED AWAYNESS :

4 6 6 10 4 6 6 0 2 0 6 n
 2 0

CUBE (11 , 8) IS SELECTED,

CUBE (4 , 7) HAS BEEN FOUND,

CUBE (4 , 13) HAS BEEN FOUND,

CUBE (14 , 15) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 0 1 2

ACCUMULATED AWAYNESS :

2 0 0 4 8 10 12 10 8

CUBE (14 , 15) IS SELECTED.

CUBE (8 , 9) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 3

ACCUMULATED AWAYNESS :

2 0 0 4 8 10 12 10 10

CUBE (6 , 9) IS SELECTED.

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 2

CUBE (11 , 1) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

2 0 0 4 8 10 12 10 10

CUBE (11 , 1) IS SELECTED.

MAXIMUM CUBE DIMENSION = 3

1 0 2 1 0 3 3 1

THE CHARACTERISTICS :

2 1 2 1 2 1 2 1 2 1 2 1 2 1 2 1 2 1 2 1
2 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 1

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 1

CUBE (31 , 27) HAS BEEN FOUND.

CUBE (31 , 30) HAS BEEN FOUND.

WATCH FOR 27 AND 30

CHECKING POINTS

0 0 0 0 0 0 0 0 0 0 0 9999 9999

ACCUMULATED AWAYNESS :

8 8 6 6 4 4 4 4 2 2 0 0

CUBE (31 , 30) IS SELECTED.

CUBE (28 , 24) HAS BEEN FOUND.

CUBE (28 , 30) HAS BEEN FOUND.

CHECKING POINTS

0 0 0 0 0 0 9999 0 0 9999 9999 9999

ACCUMULATED AWAYNESS :

12 14 12 12 10 12 4 6 6 2 0 0

CUBE (28 , 24) IS SELECTED.

CUBE (27 , 25) HAS BEEN FOUND.

CUBE (27 , 31) HAS BEEN FOUND.

CHECKING POINTS

0 0 0 0 0 0 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

18 18 18 18 18 18 4 6 6 2 0 0

CUBE (27 , 25) IS SELECTED.

CUBE (7 , 5) HAS BEEN FOUND.

CUBE (7 , 6) HAS BEEN FOUND.

WATCH FOR 5 AND 6

CHECKING POINTS

0 0 0 0 9999 9999 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

22 22 20 20 18 18 4 6 6 2 0 0

CUBE (7 , 6) IS SELECTED.

CUBE (5 , 1) HAS BEEN FOUND.

CUBE (5 , 7) HAS BEEN FOUND.

CHECKING POINTS

0 9999 0 9999 9999 9999 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

24 22 24 20 18 18 4 6 6 2 0 0

CUBE (5 , 1) IS SELECTED.

CUBE (0 , 1) HAS BEEN FOUND.

CUBE (0 , 2) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

24 22 24 20 18 18 4 6 6 2 0 0

CUBE (0 , 2) IS SELECTED.

MAXIMUM CUBE DIMENSION = 4

1 0 0 0 2 1 0 0 3 3 1 0 4 6
4 1

THE CHARACTERISTICS :

1 1 1 1 2 1 2 1 3 5 3 1 3 1 3 5 3 5 3 1 3 1
3 5 3 5 3 5 3 5 3 5

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 3

CUBE (6 , 24) HAS BEEN FOUND.

CHECKING POINTS

9999 0 1 0 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

0 4 2 6 0 0 0 0 0 0 0 0
 0 0 0 0

CUBE (8 , 24) IS SELECTED.

CUBE (14 , 30) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 1 0 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

0 4 8 12 0 0 0 0 0 0 0 0
 0 0 0 0

CUBE (14 , 30) IS SELECTED.

CUBE (16 , 28) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 0 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

0 4 8 16 0 0 0 0 0 0 0 0
 0 0 0 0

CUBE (16 , 28) IS SELECTED.

CUBE (19 , 31) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

0 4 8 16 0 0 0 0 0 0 0 0
 0 0 0 0

CUBE (19 , 31) IS SELECTED.

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 3

MAXIMUM CUBE DIMENSION = 4

1 0 0 0 2 1 0 0 3 3 1 0 4 6
4 1

THE CHARACTERISTICS :

3 1 2 1 3 1 3 5 3 5 3 1 2 1 3 1 1 1 2 1 3 1
3 1 3 5 3 1 3 1 2 1 1 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 3

CUBE (3 , 6) HAS BEEN FOUND.

CHECKING POINTS

1 9999 2 9999 9999 2 0 2 0 0 9999 9999
9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

2 0 2 0 0 2 4 2 4 6 0 0
0 0 0 0 0

CUBE (3 , 6) IS SELECTED.

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

CUBE (26 , 24) HAS BEEN FOUND.

CHECKING POINTS

1 9999 2 9999 9999 2 0 2 9999 1 9999 9999
9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

6 0 8 0 0 8 10 8 4 8 0 0
0 0 0 0 0

CUBE (26 , 24) IS SELECTED.

CUBE (28 , 8) HAS BEEN FOUND.

CHECKING POINTS

2 9999 3 9999 9999 2 0 2 9999 9999 9999 9999
9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

8 0 10 0 0 12 14 14 4 8 0 0
0 0 0 0 0

CUBE (28 , 8) IS SELECTED.

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 2

CUBE (0 , 6) HAS BEEN FOUND.

CUBE (0 , 12) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 2 0 2 9999 9999 9999 9999
9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

8 0 10 0 0 16 18 18 4 8 0 0
0 0 0 0 0

CUBE (0 , 6) IS SELECTED.

CUBE (15 , 6) HAS BEEN FOUND.

CUBE (15 , 23) HAS BEEN FOUND.

WATCH FOR 6 AND 23

CHECKING POINTS

9999 9999 9999 9999 9999 9999 0 3 9999 9999 9999 9999
9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

8 0 10 0 0 16 22 20 4 8 0 0
0 0 0 0 0

CUBE (15 , 6) IS SELECTED.

CUBE (23 , 15) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 0 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

8 0 10 0 0 16 24 20 4 8 0 0
0 0 0 0 0

CUBE (23 , 15) IS SELECTED.

CUBE (21 , 23) HAS BEEN FOUND.

CUBE (21 , 17) HAS BEEN FOUND.

WATCH FOR 23 AND 17

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

8 0 10 0 0 16 24 20 4 8 0 0
0 0 0 0 0

CUBE (21 , 17) IS SELECTED.

MAXIMUM CUBE DIMENSION = 3

1 0 0 2 1 0 3 3 1

THE CHARACTERISTICS :

3 5 3 1 2 1 3 1 3 5 3 5 3 1 3 1 3 1 3 1 3 5
3 1 2 1 3 1

CUBE (3 , 10) HAS BEEN FOUND.

CHECKING POINTS

2 9999 9999 0 2 9999 9999 0 0 0 0 2
1 0

ACCUMULATED AWAYNESS :

2 0 0 4 2 0 0 4 4 6 4 2
2 6

CUBE (3 , 10) IS SELECTED.

CUBE (4 , 24) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 4
1 9999

ACCUMULATED AWAYNESS :

2 0 0 4 2 0 0 4 4 6 4 4
6 6

CUBE (4 , 24) IS SELECTED.

CUBE (27 , 10) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999

ACCUMULATED AWAYNESS :

2 0 0 4 2 0 0 4 4 6 4 4
6 6

CUBE (27 , 10) IS SELECTED.

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 3

**FORTRAN ** STOP
/LOGOFF

C E420 LOGOFF AT 2318 ON 11/30 FOR TSN 2390,

C E421 CPU TIME USED: 000090.4681 SECONDS.

APPENDIX C

EXAMPLE OF 6 VARIABLES, 18 MINTERMS

TOTAL NUMBER OF ON-SET AND DONT-CARES = 18
 ON-SET = 18 DONT-CARES = 0
 NUMBERS OF VARIABLES = 6

MINT.#	1	=	6	MINT.#	2	=	4	MINT.#	3	=	8
MINT.#	4	=	9	MINT.#	5	=	11	MINT.#	6	=	14
MINT.#	7	=	27	MINT.#	8	=	30	MINT.#	9	=	29
MINT.#	10	=	28	MINT.#	11	=	16	MINT.#	12	=	17
MINT.#	13	=	19	MINT.#	14	=	61	MINT.#	15	=	45
MINT.#	16	=	0	MINT.#	17	=	44	MINT.#	18	=	36

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS ;
 1-CUBE = (36 , 44) = 10-100

THIS CUBE COVERS THESE MINTERMS ;
 44 36

THE 2-TH CUBE COVERING IS ;
 1-CUBE = (45 , 61) = 1-1101

THIS CUBE COVERS THESE MINTERMS ;
 61 45

THE 3-TH CUBE COVERING IS ;
 1-CUBE = (28 , 29) = 01110-

THIS CUBE COVERS THESE MINTERMS ;
 29 28

THE 4-TH CUBE COVERING IS ;
 1-CUBE = (4 , 0) = 000-00

THIS CUBE COVERS THESE MINTERMS ;
 4 0

THE 5-TH CUBE COVERING IS :
 1-CUBE = (8 , 9) = 00100-

THIS CUBE COVERS THESE MINTERMS :
 8 9

THE 6-TH CUBE COVERING IS :
 1-CUBE = (14 , 30) = 0-1110

THIS CUBE COVERS THESE MINTERMS :
 14 30

THE 7-TH CUBE COVERING IS :
 1-CUBE = (6 , 4) = 0001-0

THIS CUBE COVERS THESE MINTERMS :
 6 4

THE 8-TH CUBE COVERING IS :
 1-CUBE = (16 , 17) = 01000-

THIS CUBE COVERS THESE MINTERMS :
 16 17

THE 9-TH CUBE COVERING IS :
 1-CUBE = (11 , 27) = 0-1011

THIS CUBE COVERS THESE MINTERMS :
 11 27

THE 10-TH CUBE COVERING IS :
 1-CUBE = (19 , 17) = 0100-1

THIS CUBE COVERS THESE MINTERMS :
 17 19

THE FINAL CHECKING

105	205	105	105	105	105	105	105	105	105	105	105	20
105	105	105	105	105	105							

MAXIMUM CUBE DIMENSION = 4

1 0 0 0 2 1 0 0 3 3 1 0 4 8
4 1

THE CHARACTERISTICS :

2 1 3 1 2 1 2 1 2 1 2 1 2 1 2 1 2 1 2 1
2 1 2 1 2 1 2 1 3 1 2 1 2 1

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 3

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 1

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 1

CUBE (36 , 4) HAS BEEN FOUND.

CUBE (36 , 44) HAS BEEN FOUND.

WATCH FOR 4 AND 44

CHECKING POINTS

0 0 0 0 0 0 0 0 0 0 0 0
0 0 0 0 9999 9999

ACCUMULATED AWAYNESS :

4 2 4 6 8 4 10 6 6 4 6 8
10 4 2 4 0 0

CUBE (36 , 44) IS SELECTED.

CUBE (45 , 61) HAS BEEN FOUND.

CUBE (45 , 44) HAS BEEN FOUND.

CHECKING POINTS

0 0 0 0 0 0 0 0 0 0 0 0
0 9999 9999 0 9999 9999

ACCUMULATED AWAYNESS :

12 8 10 10 14 10 16 12 8 8 14 14
18 4 2 12 2 4

CUBE (45 , 61) IS SELECTED.

CUBE (28 , 30) HAS BEEN FOUND.

CUBE (28 , 29) HAS BEEN FOUND.

WATCH FOR 30 AND 29

CHECKING POINTS

0 0 0 0 0 0 0 0 9999 9999 0 0
 0 9999 9999 0 9999 9999

ACCUMULATED AWAYNESS :

18 12 14 14 20 14 20 14 8 8 18 18
 24 6 6 18 6 10

CUBE (28 , 29) IS SELECTED.

CUBE (4 , 6) HAS BEEN FOUND.

CUBE (4 , 0) HAS BEEN FOUND.

WATCH FOR 6 AND 0

CUBE (4 , 36) HAS BEEN FOUND.

CHECKING POINTS

0 9999 0 0 0 0 0 0 9999 9999 0 0
 0 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

20 12 16 18 26 18 28 20 14 12 20 22
 30 14 12 18 10 12

CUBE (4 , 0) IS SELECTED.

CUBE (8 , 9) HAS BEEN FOUND.

CUBE (8 , 0) HAS BEEN FOUND.

CHECKING POINTS

0 9999 9999 9999 0 0 0 0 9999 9999 0 0
 0 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

26 16 16 18 28 22 32 26 18 16 24 26
 36 20 16 20 14 18

CUBE (8 , 9) IS SELECTED.

CUBE (14 , 6) HAS BEEN FOUND.

CUBE (14 , 30) HAS BEEN FOUND.

WATCH FOR 6 AND 30

ACCUMULATED AWAYNESS :

40	32	32	36	44	36	48	40	36	32	40	44
52	46	42	36	38	42						

CUBE (11 , 27) IS SELECTED.

CUBE (19 , 27) HAS BEEN FOUND.

CUBE (19 , 17) HAS BEEN FOUND.

WATCH FOR 27 AND 17

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999						

ACCUMULATED AWAYNESS :

46	38	38	40	48	44	50	46	40	38	42	44
52	52	50	40	48	50						

CUBE (19 , 17) IS SELECTED.

APPENDIX D

EXAMPLE OF 5 VARIABLES, 20 MINTERMS

TOATAL NUMBER OF ON-SET AND DONT-CARES = 20
 ON-SET = 20 DONT-CARES = 0
 NUMBERS OF VARIABLES = 5

MINT.# 1 = 1	MINT.# 2 = 3	MINT.# 3 = 5
MINT.# 4 = 4	MINT.# 5 = 8	MINT.# 6 = 9
MINT.# 7 = 11	MINT.# 8 = 10	MINT.# 9 = 25
MINT.# 10 = 27	MINT.# 11 = 26	MINT.# 12 = 30
MINT.# 13 = 29	MINT.# 14 = 28	MINT.# 15 = 19
MINT.# 16 = 0	MINT.# 17 = 18	MINT.# 18 = 22
MINT.# 19 = 21	MINT.# 20 = 20	MINT.#

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 2-CUBE = (20 , 29) = 1-10-

THIS CUBE COVERS THESE MINTERMS :
 29 28 21 20

THE 2-TH CUBE COVERING IS :
 2-CUBE = (18 , 30) = 1--10

THIS CUBE COVERS THESE MINTERMS :
 26 30 18 22

THE 3-TH CUBE COVERING IS :
 2-CUBE = (0 , 5) = 00-0-

THIS CUBE COVERS THESE MINTERMS :
 1 5 4 0

THE 4-TH CUBE COVERING IS :
 2-CUBE = (11 , 19) = --011

THIS CUBE COVERS THESE MINTERMS :
 3 11 27 19

THE 5-TH CUBE COVERING IS :
 2-CUBE = (8 , 11) = 010--

THIS CUBE COVERS THESE MINTERMS :
 8 9 11 10

THE 6-TH CUBE COVERING IS :
 2-CUBE = (25 , 11) = -10-1

THIS CUBE COVERS THESE MINTERMS :
 9 11 25 27

THE FINAL CHECKING

105	105	105	105	105	205	305	105	105	205	105	105
105	105	105	105	105	105	105	105				

```
* 5
* 1 3 5 4 8 9 11 10 25 27 26 30 29 28 19
* 18 22 21 209
```

MAXIMUM CUBE DIMENSION = 4

```
1 0 0 0 2 1 0 0 3 3 1 0 4 6
4 1
```

THE CHARACTERISTICS :

```
4 1 3 1 3 1 3 1 3 1 4 1 3 5 3 1 3 1 4 1 4 1
3 1 3 1 3 1 3 1 3 1 3 1 3 1 3 1 3 5
```

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 3

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 2

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 2

CUBE (20 , 5) HAS BEEN FOUND.

CUBE (20 , 30) HAS BEEN FOUND.

WATCH FOR 5 AND 30

CUBE (20 , 29) HAS BEEN FOUND.

WATCH FOR 30 AND 29

CHECKING POINTS

```
0 0 1 1 0 0 0 0 1 0 0 1
9999 9999 0 0 0 1 9999 9999
```

ACCUMULATED AWAYNESS :

```
4 6 2 2 4 4 6 6 2 4 4 2
0 0 4 4 4 2 0 0
```

CUBE (20 , 29) IS SELECTED.

CUBE (18 , 27) HAS BEEN FOUND.

CUBE (18 , 30) HAS BEEN FOUND.

WATCH FOR 27 AND 30

CHECKING POINTS

```
0 0 1 1 0 0 0 1 1 1 9999 9999
9999 9999 1 0 9999 9999 9999 9999
```

ACCUMULATED AWAYNESS :

10	10	8	6	8	10	10	8	6	6	4	2
4	2	6	8	4	2	4	2				

CUBE (18 , 30) IS SELECTED.

CUBE (0 , 5) HAS BEEN FOUND.

CUBE (0 , 9) HAS BEEN FOUND.

WATCH FOR 5 AND 9

CHECKING POINTS

9999	1	9999	9999	1	1	0	1	1	1	9999	9999
9999	9999	1	9999	9999	9999	9999	9999				

ACCUMULATED AWAYNESS :

10	12	8	6	10	12	14	12	10	12	10	8
8	6	10	8	8	6	6	4				

CUBE (0 , 5) IS SELECTED.

CUBE (11 , 1) HAS BEEN FOUND.

CUBE (11 , 8) HAS BEEN FOUND.

CUBE (11 , 25) HAS BEEN FOUND.

WATCH FOR 8 AND 25

CUBE (11 , 26) HAS BEEN FOUND.

CUBE (11 , 19) HAS BEEN FOUND.

WATCH FOR 25 AND 19

CHECKING POINTS

9999	9999	9999	9999	1	2	9999	2	2	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999				

ACCUMULATED AWAYNESS :

12	12	12	12	14	14	14	14	12	12	12	12
12	12	10	12	10	10	10	10				

CUBE (11 , 19) IS SELECTED.

CUBE (8 , 1) HAS BEEN FOUND.

CUBE (8 , 11) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 3 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

14	14	16	16	14	14	14	14	14	14	14	16
16	16	14	14	14	16	16	16				

CUBE (8 , 11) IS SELECTED.

CUBE (25 , 11) HAS BEEN FOUND.

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

16	16	20	22	16	14	14	16	14	14	16	20
18	20	16	18	18	22	20	22				

CUBE (25 , 11) IS SELECTED.

APPENDIX E

EXAMPLE OF 7 VARIABLES, 53 MINTERMS

TOATAL NUMBER OF ON-SET AND DONT-CARES = 63
 ON-SET = 53 DONT-CARES = 10
 NUMBERS OF VARIABLES = 7

MINT.# 1 = 5	MINT.# 2 = 10	MINT.# 3 = 13
MINT.# 4 = 14	MINT.# 5 = 21	MINT.# 6 = 22
MINT.# 7 = 29	MINT.# 8 = 30	MINT.# 9 = 34
MINT.# 10 = 37	MINT.# 11 = 42	MINT.# 12 = 45
MINT.# 13 = 50	MINT.# 14 = 53	MINT.# 15 = 61
MINT.# 16 = 64	MINT.# 17 = 65	MINT.# 18 = 66
MINT.# 19 = 67	MINT.# 20 = 68	MINT.# 21 = 69
MINT.# 22 = 70	MINT.# 23 = 71	MINT.# 24 = 72
MINT.# 25 = 74	MINT.# 26 = 80	MINT.# 27 = 81
MINT.# 28 = 82	MINT.# 29 = 83	MINT.# 30 = 84
MINT.# 31 = 85	MINT.# 32 = 86	MINT.# 33 = 87
MINT.# 34 = 88	MINT.# 35 = 89	MINT.# 36 = 90
MINT.# 37 = 91	MINT.# 38 = 92	MINT.# 39 = 93
MINT.# 40 = 94	MINT.# 41 = 95	MINT.# 42 = 96
MINT.# 43 = 97	MINT.# 44 = 98	MINT.# 45 = 99
MINT.# 46 = 100	MINT.# 47 = 101	MINT.# 48 = 102
MINT.# 49 = 103	MINT.# 50 = 104	MINT.# 51 = 105
MINT.# 52 = 106	MINT.# 53 = 107	MINT.# 54 = 108
MINT.# 55 = 109	MINT.# 56 = 110	MINT.# 57 = 111
MINT.# 58 = 112	MINT.# 59 = 113	MINT.# 60 = 114
MINT.# 61 = 115	MINT.# 62 = 116	MINT.# 63 = 117

THE CUBE COVERING LIST:

THE 1-TH CUBE COVERING IS :
 3-CUBE = (13 , 53) = 0---101

THIS CUBE COVERS THESE MINTERMS :
 13 53 5 21 29 37 45 61

THIS MINTERM MT(3) = 13 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
 2-CUBE = (22 , 94) = -01-110

THIS CUBE COVERS THESE MINTERMS :
 22 94 30 86

THIS MINTERM MT(6) = 22 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
 2-CUBE = (50 , 98) = -1-0010

THIS CUBE COVERS THESE MINTERMS :
 50 98 34 114

THIS MINTERM MT(13) = 50 IS A CORE POINT

THE 4-TH CUBE COVERING IS :
 4-CUBE = (89 , 86) = 101----

THIS CUBE COVERS THESE MINTERMS :
 89 86 80 81 82 83 84 85 87 88 90 91 92 93 94 95

THIS MINTERM MT(35) = 89 IS A CORE POINT

THE 5-TH CUBE COVERING IS :
 4-CUBE = (105 , 102) = 110----

THIS CUBE COVERS THESE MINTERMS :
 105 102 96 97 98 99 100 101 103 104 106 107 108 109 110 111

THIS MINTERM MT(51) = 105 IS A CORE POINT

THE 6-TH CUBE COVERING IS :
 4-CUBE = (67 , 84) = 10-0---

THIS CUBE COVERS THESE MINTERMS :
 67 84 64 65 66 68 69 70 71 80 81 82 83 85 86 87

THE 7-TH CUBE COVERING IS :
 4-CUBE = (117 , 64) = 1--0-0-

THIS CUBE COVERS THESE MINTERMS :
 117 64 65 68 69 80 81 84 85 96 97 100 101 112 113 116

THE 8-TH CUBE COVERING IS :
 4-CUBE = (114 , 65) = 1--00--

THIS CUBE COVERS THESE MINTERMS :
 114 65 64 66 67 80 81 82 83 96 97 98 99 112 113 115

CUBE (22 , 94) IS SELECTED.

CUBE (51 , 98) HAS BEEN FOUND.

CHECKING POINTS

9999	0	9999	0	9999	9999	9999	9999	9999	9999	1	9999
9999	9999	9999	2	2	6	2	4	8	6	4	0
4	4	2	2	4	6	10	9999	6	2	1	6
1	5	10	9999	5	4	4	9999	4	4	8	6
4	2	1	4	1	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999									

ACCUMULATED AWAYNESS :

0	14	0	12	0	4	0	4	12	0	14	0
18	0	0	16	18	14	16	14	16	12	14	18
16	14	16	12	14	12	14	6	12	16	18	14
16	14	16	6	14	16	18	14	16	14	16	12
14	18	20	14	18	0	0	0	0	0	0	0
0	0	0									

CUBE (51 , 98) IS SELECTED.

MAXIMUM CUBE DIMENSION HAS BEEN LOWERED TO: 4

CUBE (89 , 86) HAS BEEN FOUND.

CHECKING POINTS

9999	0	9999	0	9999	9999	9999	9999	9999	9999	1	9999
9999	9999	9999	6	6	10	6	8	12	10	8	2
8	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	5	5	9999	5	5	9	7
5	3	2	5	2	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999									

ACCUMULATED AWAYNESS :

0	18	0	16	0	4	0	4	12	0	20	0
18	0	0	18	20	16	18	16	18	14	16	20
18	14	16	12	14	12	14	6	12	16	18	14
16	14	16	6	14	20	22	14	20	18	20	16
18	22	24	20	22	0	0	0	0	0	0	0
0	0	0									

CUBE (89 , 86) IS SELECTED.

CUBE (105 , 102) HAS BEEN FOUND.

CHECKING POINTS

9999	0	9999	0	9999	9999	9999	9999	9999	9999	2	9999
9999	9999	9999	10	10	14	10	12	16	14	12	4
12	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999									

ACCUMULATED AWAYNESS :

0	22	0	20	0	4	0	4	12	0	22	0
10	0	0	20	22	18	20	18	20	16	18	22
20	14	16	12	14	12	14	6	12	16	18	14
16	14	16	6	14	20	22	14	20	18	20	16
18	22	24	20	22	0	0	0	0	0	0	0
	0	0									

CUBE (105 , 102) IS SELECTED.

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 4

CUBE (67 , 84) HAS BEEN FOUND.

CUBE (67 , 100) HAS BEEN FOUND.

WATCH FOR 84 AND 100

CUBE (67 , 112) HAS BEEN FOUND.

CHECKING POINTS

9999	0	9999	0	9999	9999	9999	9999	9999	9999	2	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	6
16	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999									

ACCUMULATED AWAYNESS :

0	26	0	24	0	4	0	4	12	0	28	0
10	0	0	20	22	18	20	18	20	16	18	24
22	14	16	12	14	12	14	6	12	16	18	14
16	14	16	6	14	20	22	14	20	18	20	16
18	22	24	20	22	0	0	0	0	0	0	0
	0	0									

CUBE (67 , 84) IS SELECTED.

CUBE (117 , 64) HAS BEEN FOUND.

ACCUMULATED AWAYNESS :

0	38	0	42	0	4	0	4	12	0	42	0
10	0	0	20	22	18	20	18	20	16	18	28
28	14	16	12	14	12	14	6	12	16	18	14
16	14	16	6	14	20	22	14	20	18	20	16
18	22	24	20	22	0	0	0	0	0	0	0
0	0	0									

CUBE (14 , 30) IS SELECTED.

**FORTRAN ** STOP
/LOGOFF

C E420 LOGOFF AT 1129 ON 11/14 , FOR TSN 4173.

C E421 CPU TIME USED: 000101.8603 SECONDS.

APPENDIX F

EXAMPLE OF 8 VARIABLES, 126 MINTERMS

TOATAL NUMBER OF ON-SET AND DONT*CARES = 126
 ON-SET = 126 DONT*CARES = 0
 NUMBERS OF VARIABLES = 8

MINT.# 1 = 0	MINT.# 2 = 1	MINT.# 3 = 2
MINT.# 4 = 3	MINT.# 5 = 4	MINT.# 6 = 5
MINT.# 7 = 6	MINT.# 8 = 7	MINT.# 9 = 8
MINT.# 10 = 9	MINT.# 11 = 10	MINT.# 12 = 11
MINT.# 13 = 12	MINT.# 14 = 13	MINT.# 15 = 14
MINT.# 16 = 15	MINT.# 17 = 16	MINT.# 18 = 17
MINT.# 19 = 18	MINT.# 20 = 19	MINT.# 21 = 20
MINT.# 22 = 21	MINT.# 23 = 22	MINT.# 24 = 23
MINT.# 25 = 24	MINT.# 26 = 25	MINT.# 27 = 26
MINT.# 28 = 27	MINT.# 29 = 28	MINT.# 30 = 29
MINT.# 31 = 30	MINT.# 32 = 31	MINT.# 33 = 32
MINT.# 34 = 33	MINT.# 35 = 34	MINT.# 36 = 35
MINT.# 37 = 36	MINT.# 38 = 37	MINT.# 39 = 38
MINT.# 40 = 39	MINT.# 41 = 40	MINT.# 42 = 41
MINT.# 43 = 42	MINT.# 44 = 43	MINT.# 45 = 44
MINT.# 46 = 45	MINT.# 47 = 46	MINT.# 48 = 47
MINT.# 49 = 48	MINT.# 50 = 49	MINT.# 51 = 50
MINT.# 52 = 51	MINT.# 53 = 52	MINT.# 54 = 53
MINT.# 55 = 54	MINT.# 56 = 55	MINT.# 57 = 56
MINT.# 58 = 57	MINT.# 59 = 58	MINT.# 60 = 59
MINT.# 61 = 60	MINT.# 62 = 61	MINT.# 63 = 62
MINT.# 64 = 63	MINT.# 65 = 64	MINT.# 66 = 65
MINT.# 67 = 66	MINT.# 68 = 67	MINT.# 69 = 68
MINT.# 70 = 69	MINT.# 71 = 70	MINT.# 72 = 71
MINT.# 73 = 72	MINT.# 74 = 74	MINT.# 75 = 75
MINT.# 76 = 76	MINT.# 77 = 77	MINT.# 78 = 78
MINT.# 79 = 79	MINT.# 80 = 80	MINT.# 81 = 81
MINT.# 82 = 82	MINT.# 83 = 83	MINT.# 84 = 84
MINT.# 85 = 85	MINT.# 86 = 86	MINT.# 87 = 87
MINT.# 88 = 88	MINT.# 89 = 89	MINT.# 90 = 90
MINT.# 91 = 91	MINT.# 92 = 92	MINT.# 93 = 93
MINT.# 94 = 94	MINT.# 95 = 95	MINT.# 96 = 96
MINT.# 97 = 97	MINT.# 98 = 98	MINT.# 99 = 99
MINT.# 100 = 100	MINT.# 101 = 101	MINT.# 102 = 103
MINT.# 103 = 104	MINT.# 104 = 105	MINT.# 105 = 106
MINT.# 106 = 107	MINT.# 107 = 108	MINT.# 108 = 109
MINT.# 109 = 110	MINT.# 110 = 111	MINT.# 111 = 112
MINT.# 112 = 113	MINT.# 113 = 114	MINT.# 114 = 115
MINT.# 115 = 116	MINT.# 116 = 117	MINT.# 117 = 118
MINT.# 118 = 119	MINT.# 119 = 120	MINT.# 120 = 121
MINT.# 121 = 122	MINT.# 122 = 123	MINT.# 123 = 124
MINT.# 124 = 125	MINT.# 125 = 126	MINT.# 126 = 127

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
6-CUBE = (9 , 54) = 00-----

THIS CUBE COVERS THESE MINTERMS :

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
32	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47
48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63

THIS MINTERM MT(10) = 9 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
6-CUBE = (89 , 54) = 0--1----

THIS CUBE COVERS THESE MINTERMS :

16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
48	49	50	51	52	53	54	55	56	57	58	59	60	61	62	63
80	81	82	83	84	85	86	87	88	89	90	91	92	93	94	95
112	113	114	115	116	117	118	119	120	121	122	123	124	125	126	127

THIS MINTERM MT(89) = 89 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
5-CUBE = (77 , 62) = 0---11--

THIS CUBE COVERS THESE MINTERMS :

12	13	14	15	28	29	30	31	44	45	46	47	60	61	62	63
76	77	78	79	92	93	94	95	108	109	110	111	124	125	126	127

THE 4-TH CUBE COVERING IS :
5-CUBE = (65 , 55) = 0---0--1

THIS CUBE COVERS THESE MINTERMS :

1	3	5	7	17	19	21	23	33	35	37	39	49	51	53	55
65	67	69	71	81	83	85	87	97	99	101	103	113	115	117	119

THE 5-TH CUBE COVERING IS :
 5-CUBE = (72 , 50) = 0----0-0

THIS CUBE COVERS THESE MINTERMS :

0	2	8	10	16	18	24	26	32	34	40	42	48	50	56	58
64	66	72	74	80	82	88	90	96	98	104	106	112	114	120	122

THE 6-TH CUBE COVERING IS :
 5-CUBE = (75 , 62) = 0---1-1-

THIS CUBE COVERS THESE MINTERMS :

10	11	14	15	26	27	30	31	42	43	46	47	58	59	62	63
74	75	78	79	90	91	94	95	106	107	110	111	122	123	126	127

THE 7-TH CUBE COVERING IS :
 5-CUBE = (70 , 29) = 0-0--1--

THIS CUBE COVERS THESE MINTERMS :

4	5	6	7	12	13	14	15	20	21	22	23	28	29	30	31
68	69	70	71	76	77	78	79	84	85	86	87	92	93	94	95

THE 8-TH CUBE COVERING IS :
 5-CUBE = (105 , 52) = 0-1---0-

THIS CUBE COVERS THESE MINTERMS :

32	33	36	37	40	41	44	45	48	49	52	53	56	57	60	61
96	97	100	101	104	105	108	109	112	113	116	117	120	121	124	125

THE FINAL CHECKING

205	205	205	205	205	305	205	305	205	101	305	205
305	305	405	405	305	305	305	305	305	405	305	405
305	205	405	305	405	405	505	505	305	305	205	205
205	305	101	205	305	205	305	205	305	305	305	305
405	405	305	305	305	405	205	305	405	305	405	305
405	405	405	405	105	105	105	105	105	205	105	205
105	205	105	205	205	305	305	205	205	205	205	205
305	205	305	205	101	305	205	305	305	405	405	205
205	105	105	105	205	105	205	105	205	105	205	205
205	205	305	305	205	205	205	305	101	205	305	205
305	205	305	305	305	305	305					

MAXIMUM CUBE DIMENSION = 6

1	0	0	0	0	0	2	1	0	0	0	0	3	3
1	0	0	0	4	6	4	1	0	0	5	10	10	5
1	0	6	15	20	15	6	1						

THE CHARACTERISTICS :

6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	1	6	5
6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5
6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5
6	5	6	5	6	5	6	5	6	5	6	1	6	5	6	5	6	5	6	5	6	5
6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5
6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	1
6	5	6	5	6	5	6	5	6	5	6	1	6	5	6	5	6	5	6	5	6	1
6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5
6	1	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5
6	1	6	5	6	1	6	5	6	5	6	1	6	5	6	5	6	5	6	5	6	5
6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5	6	5
6	5	6	5	6	5	6	5	6	5	6	5	6	1	6	5	6	5	6	5	6	5

CUBE (9 , 54) HAS BEEN FOUND.

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	16	16	16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16	16	16	16	16	16	16
16	16	16	16	16	16	16	16	16	16	16	16	16	16

ACCUMULATED AWAYNESS :

0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2	2

CUBE (9 , 54) IS SELECTED.

CUBE (89 , 54) HAS BEEN FOUND.

ACCUMULATED AWAYNESS :

0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	10	8	10	8	8	6	8	6	6
10	10	8	4	4	4	4	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	10
8	10	8	8	6	6	10	8	10	8	4	4	4
4	4	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2						

- CUBE (65 , 55) IS SELECTED.
- CUBE (72 , 22) HAS BEEN FOUND.
- CUBE (72 , 50) HAS BEEN FOUND.
- CUBE (72 , 52) HAS BEEN FOUND.
- CUBE (72 , 62) HAS BEEN FOUND.

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	56	9999	48	9999	9999
9999	9999	48	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	56	9999	9999	9999	56	9999	56	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	10	8	10	8	10	6	10	6	6
10	10	10	4	4	4	4	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	10
8	10	8	10	6	6	10	10	10	10	4	4	4
4	4	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2						

- CUBE (72 , 50) IS SELECTED.

CUBE (75 , 22) HAS BEEN FOUND.

CUBE (75 , 50) HAS BEEN FOUND.

WATCH FOR 22 AND 50

CUBE (75 , 55) HAS BEEN FOUND.

WATCH FOR 50 AND 55

CUBE (75 , 62) HAS BEEN FOUND.

WATCH FOR 55 AND 62

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	58	9999	56	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	58	9999	9999	9999	64	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	10	8	10	8	14	6	12	6	6
10	10	10	4	4	4	4	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	10
8	10	8	14	6	6	10	12	10	10	4	4	4
4	4	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	2

CUBE (75 , 62) IS SELECTED.

CUBE (70 , 17) HAS BEEN FOUND.

CUBE (70 , 24) HAS BEEN FOUND.

WATCH FOR 17 AND 24

CUBE (70 , 27) HAS BEEN FOUND.

CUBE (70 , 29) HAS BEEN FOUND.

ACCUMULATED AWAYNESS :

0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	10	8	10	8	14	6	12	6	6
10	10	10	4	4	4	4	2	2	2	2	2	2
2	2	2	2	2	2	2	2	2	2	2	2	10
8	10	8	16	6	6	10	16	10	10	4	4	4
4	4	2	2	2	2	2	2	2	2	2	2	2
2	2	2	2	2	2	2						

CUBE (105 , 52) IS SELECTED.

*9

**FORTRAN ** STOP

/LOGOFF BUT

%C E420 LOGOFF AT 1125 ON 11/09/77, FOR TSN 0951.

%C E421 CPU TIME USED: 000212.3728 SECONDS.

APPENDIX G

EXAMPLE OF 7 VARIABLES, 92 MINTERMS

TOATAL NUMBER OF ON-SET AND DONT-CARES = 92
 ON-SET = 92 DONT-CARES = 0
 NUMBERS OF VARIABLES = 7

MINT.# 1 = 34	MINT.# 2 = 35	MINT.# 3 = 37
MINT.# 4 = 39	MINT.# 5 = 44	MINT.# 6 = 47
MINT.# 7 = 48	MINT.# 8 = 49	MINT.# 9 = 50
MINT.# 10 = 51	MINT.# 11 = 52	MINT.# 12 = 53
MINT.# 13 = 54	MINT.# 14 = 55	MINT.# 15 = 56
MINT.# 16 = 57	MINT.# 17 = 58	MINT.# 18 = 61
MINT.# 19 = 62	MINT.# 20 = 63	MINT.# 21 = 64
MINT.# 22 = 83	MINT.# 23 = 84	MINT.# 24 = 85
MINT.# 25 = 86	MINT.# 26 = 87	MINT.# 27 = 88
MINT.# 28 = 89	MINT.# 29 = 90	MINT.# 30 = 95
MINT.# 31 = 1	MINT.# 32 = 2	MINT.# 33 = 3
MINT.# 34 = 4	MINT.# 35 = 5	MINT.# 36 = 6
MINT.# 37 = 7	MINT.# 38 = 8	MINT.# 39 = 9
MINT.# 40 = 10	MINT.# 41 = 11	MINT.# 42 = 14
MINT.# 43 = 15	MINT.# 44 = 16	MINT.# 45 = 17
MINT.# 46 = 18	MINT.# 47 = 19	MINT.# 48 = 20
MINT.# 49 = 21	MINT.# 50 = 22	MINT.# 51 = 23
MINT.# 52 = 24	MINT.# 53 = 25	MINT.# 54 = 26
MINT.# 55 = 27	MINT.# 56 = 28	MINT.# 57 = 29
MINT.# 58 = 30	MINT.# 59 = 31	MINT.# 60 = 32
MINT.# 61 = 96	MINT.# 62 = 97	MINT.# 63 = 98
MINT.# 64 = 99	MINT.# 65 = 100	MINT.# 66 = 101
MINT.# 67 = 102	MINT.# 68 = 103	MINT.# 69 = 104
MINT.# 70 = 105	MINT.# 71 = 106	MINT.# 72 = 107
MINT.# 73 = 108	MINT.# 74 = 109	MINT.# 75 = 110
MINT.# 76 = 111	MINT.# 77 = 112	MINT.# 78 = 113
MINT.# 79 = 114	MINT.# 80 = 115	MINT.# 81 = 116
MINT.# 82 = 117	MINT.# 83 = 118	MINT.# 84 = 119
MINT.# 85 = 120	MINT.# 86 = 121	MINT.# 87 = 122
MINT.# 88 = 123	MINT.# 89 = 124	MINT.# 90 = 125
MINT.# 91 = 126	MINT.# 92 = 127	MINT.#

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 1-CUBE = (44 , 108) = -101100

THIS CUBE COVERS THESE MINTERMS :
 44 108

THIS MINTERM MT(5) = 44 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
 1-CUBE = (64 , 96) = 1-00000

THIS CUBE COVERS THESE MINTERMS :
 64 96

THIS MINTERM MT(21) = 64 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
 3-CUBE = (83 , 55) = --10-11

THIS CUBE COVERS THESE MINTERMS :
 83 55 51 87 19 23 115 119

THIS MINTERM MT(22) = 83 IS A CORE POINT

THE 4-TH CUBE COVERING IS :
 4-CUBE = (84 , 55) = --101--

THIS CUBE COVERS THESE MINTERMS :
 84 55 52 53 54 85 86 87 20 21 22 23
 116 117 118 119

THIS MINTERM MT(23) = 84 IS A CORE POINT

THE 5-TH CUBE COVERING IS :
 3-CUBE = (89 , 56) = --1100-

THIS CUBE COVERS THESE MINTERMS :
 89 56 57 88 24 25 120 121

THIS MINTERM MT(28) = 89 IS A CORE POINT

THE 6-TH CUBE COVERING IS :
 3-CUBE = (90 , 56) = --110-0

THIS CUBE COVERS THESE MINTERMS :
 90 56 58 88 24 26 120 122

THIS MINTERM MT(29) = 90 IS A CORE POINT

THE 7-TH CUBE COVERING IS :
 3-CUBE = (95 , 55) = --1-111

THIS CUBE COVERS THESE MINTERMS :
 95 55 63 87 23 31 119 127

THIS MINTERM MT(30) = 95 IS A CORE POINT

THE 8-TH CUBE COVERING IS :
 3-CUBE = (4 , 23) = 00-01--

THIS CUBE COVERS THESE MINTERMS :
 4 23 5 6 7 20 21 22

THIS MINTERM MT(34) = 4 IS A CORE POINT

THE 9-TH CUBE COVERING IS :
 3-CUBE = (8 , 27) = 00-10--

THIS CUBE COVERS THESE MINTERMS :
 8 27 9 10 11 24 25 26

THIS MINTERM MT(38) = 8 IS A CORE POINT

THE 10-TH CUBE COVERING IS :
 4-CUBE = (14 , 19) = 00---1-

THIS CUBE COVERS THESE MINTERMS :
 14 19 2 3 6 7 10 11 15 18 22 23
 26 27 30 31

THIS MINTERM MT(42) = 14 IS A CORE POINT

THE 11-TH CUBE COVERING IS :
 4-CUBE = (28 , 19) = 001----

THIS CUBE COVERS THESE MINTERMS :
 28 19 16 17 18 20 21 22 23 24 25 26
 27 29 30 31

THIS MINTERM MT(56) = 28 IS A CORE POINT

THE 12-TH CUBE COVERING IS :
 3-CUBE = (32 , 114) = -1-00-0

THIS CUBE COVERS THESE MINTERMS :
 32 114 34 48 50 96 98 112

THIS MINTERM MT(60) = 32 IS A CORE POINT

THE 13-TH CUBE COVERING IS :
 5-CUBE = (97 , 126) = 11-----

THIS CUBE COVERS THESE MINTERMS :
 97 126 96 98 99 100 101 102 103 104 105 106
 107 108 109 110 111 112 113 114 115 116 117 118
 119 120 121 122 123 124 125 127

THIS MINTERM MT(62) = 97 IS A CORE POINT

THE 14-TH CUBE COVERING IS :
 4-CUBE = (49 , 118) = -110---

THIS CUBE COVERS THESE MINTERMS :
 49 118 48 50 51 52 53 54 55 112 113 114
 115 116 117 119

THE 15-TH CUBE COVERING IS :
 3-CUBE = (47 , 119) = -1--111

THIS CUBE COVERS THESE MINTERMS :
 47 119 39 55 63 103 111 127

THE 16-TH CUBE COVERING IS :
 3-CUBE = (37 , 119) = -1-01-1

THIS CUBE COVERS THESE MINTERMS :
 37 119 39 53 55 101 103 117

THE 17-TH CUBE COVERING IS :
 3-CUBE = (1 , 23) = 00-0--1

THIS CUBE COVERS THESE MINTERMS :
 1 23 3 5 7 17 19 21

THE 18-TH CUBE COVERING IS :
 3-CUBE = (35 , 119) = -1-0-11

THIS CUBE COVERS THESE MINTERMS :
 35 119 39 51 55 99 103 115

THE 19-TH CUBE COVERING IS :
 3-CUBE = (61 , 119) = -11-1-1

THIS CUBE COVERS THESE MINTERMS :
 61 119 53 55 63 117 125 127

THE 20-TH CUBE COVERING IS :
 3-CUBE = (62 , 119) = -11-11-

THIS CUBE COVERS THESE MINTERMS :
 62 119 54 55 63 118 126 127

THE FINAL CHECKING

105	105	105	305	101	105	205	105	205	305	205	405
305	905	205	105	105	105	105	405	101	101	101	101
101	305	205	101	101	101	105	105	205	101	201	201
305	101	101	201	201	101	101	105	205	205	405	305
405	405	705	405	305	405	301	101	101	205	305	101
305	101	205	205	101	205	101	405	101	101	101	101
205	101	101	205	305	205	305	405	305	505	405	1005
305	205	205	101	101	205	205	505				

0	0	0	0	0	0	0	1	0	0	0	2
9999	2	2	2	4	4	4	0	4	4	4	0
9999	2	2	2	2	2	2	0	4	0	0	0
4	0	0	0	2	2	2	0				0

ACCUMULATED AWAYNESS :

10	14	10	14	0	14	10	14	14	16	10	14
14	16	10	14	14	14	14	16	6	16	10	14
14	16	10	14	14	16	12	12	16	6	12	12
16	6	12	12	16	12	16	12	16	16	20	12
16	16	20	12	16	16	20	12	16	16	20	6
4	6	6	12	4	6	6	12	4	6	6	12
0	6	6	12	6	12	12	16	6	12	12	16
6	12	12	16	6	12	12	16				

CUBE (64 , 96) IS SELECTED.

CUBE (83 , 55) HAS BEEN FOUND.

CHECKING POINTS

2	4	0	6	9999	1	4	6	6	9999	3	6
6	9999	2	1	1	3	3	6	9999	9999	2	4
4	9999	1	0	0	2	2	4	6	0	2	2
6	0	0	0	2	1	2	4	6	6	9999	2
6	6	9999	0	2	2	6	1	2	2	6	2
9999	4	4	10	4	6	6	6	4	4	4	2
9999	2	2	4	4	10	10	9999	6	6	6	9999
1	4	2	2	6	2	4	6				

ACCUMULATED AWAYNESS :

14	16	14	16	0	16	14	16	16	16	14	16
16	16	16	16	16	16	16	20	6	16	14	16
16	16	16	16	16	20	16	16	16	14	16	16
16	16	16	16	20	16	20	16	16	16	20	16
16	16	20	16	20	20	22	16	20	20	22	12
4	12	12	14	10	12	12	14	12	14	14	16
0	14	14	16	12	14	14	16	12	14	14	16
14	16	16	16	14	16	16	16				

CUBE (83 , 55) IS SELECTED.

CUBE (84 , 55) HAS BEEN FOUND.

CHECKING POINTS

3	5	2	16	9999	1	12	16	16	9999	9999	9999
9999	9999	3	2	2	7	7	16	9999	9999	9999	9999
9999	9999	2	0	0	4	3	6	10	2	10	10
16	0	0	0	2	2	4	12	16	16	9999	9999
9999	9999	9999	2	4	4	10	5	10	10	16	2

9999	6	6	12	12	14	14	16	4	4	4	2
9999	4	4	6	12	18	18	9999	9999	9999	9999	9999
6	4	4	10	10	12	12	16				

ACCUMULATED AWAYNESS :

18	20	16	18	0	22	16	18	18	18	14	16
16	18	20	22	22	20	20	22	6	18	14	16
16	18	20	22	22	22	20	20	22	16	18	18
20	22	24	24	26	22	24	18	20	20	20	16
18	18	20	22	24	24	26	20	22	22	24	16
4	16	16	18	12	14	14	16	18	20	20	22
0	18	18	20	14	16	16	16	12	14	14	16
18	20	20	22	16	18	18	20				

CUBE (84 , 55) IS SELECTED.

CUBE (89 , 56) HAS BEEN FOUND.

CHECKING POINTS

3	5	2	16	9999	1	20	24	18	9999	9999	9999
9999	9999	9999	9999	6	11	8	18	9999	9999	9999	9999
9999	9999	9999	9999	2	4	4	6	10	2	10	10
16	2	4	2	4	2	4	20	24	18	9999	9999
9999	9999	9999	9999	9999	12	18	9	18	12	18	2
9999	8	6	12	12	14	14	16	12	12	6	4
9999	6	4	6	20	26	20	9999	9999	9999	9999	9999
9999	9999	12	18	18	20	14	18				

ACCUMULATED AWAYNESS :

24	26	22	26	0	28	18	20	22	18	14	16
16	18	20	22	24	22	24	26	6	18	14	16
16	18	20	22	24	26	24	26	28	22	24	26
28	24	26	28	30	28	30	20	22	24	20	16
18	18	20	22	24	26	28	22	24	26	28	20
4	20	22	24	18	20	22	24	20	22	24	26
0	22	24	26	16	18	20	16	12	14	14	16
18	20	22	24	18	20	22	24				

CUBE (89 , 56) IS SELECTED.

CUBE (90 , 56) HAS BEEN FOUND.

CHECKING POINTS

4	5	2	16	9999	1	28	26	26	9999	9999	9999
9999	9999	9999	9999	9999	12	12	20	9999	9999	9999	9999
9999	9999	9999	9999	9999	4	4	8	10	2	10	10
16	4	5	10	6	3	4	28	26	26	9999	9999
9999	9999	9999	9999	9999	9999	26	13	20	20	20	2
9999	8	8	12	12	14	14	16	20	14	14	6

9999 6 6 6 26 26 26 9999 9999 9999 9999 9999
 9999 9999 9999 26 26 22 22 20

ACCUMULATED AWAYNESS :

26 32 30 34 0 34 20 24 24 18 14 16
 16 18 20 22 24 26 26 30 6 18 14 16
 16 18 20 22 24 30 30 30 34 26 32 32
 36 26 30 30 34 32 36 22 26 26 20 16
 18 18 20 22 24 26 30 24 26 26 32 24
 4 26 26 30 24 26 26 32 22 26 26 30
 0 26 26 32 18 22 22 16 12 14 14 16
 18 20 22 26 20 24 24 26

CUBE (90 , 56) IS SELECTED.

CUBE (95 , 55) HAS BEEN FOUND.

CHECKING POINTS

4 6 2 24 9999 3 26 26 26 9999 9999 9999
 9999 9999 9999 9999 9999 16 16 9999 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 4 8 12 2 12 12
 24 4 5 10 8 4 12 26 26 26 9999 9999
 9999 9999 9999 9999 9999 9999 34 14 26 26 9999 2
 9999 8 8 14 12 16 16 24 20 14 14 8
 9999 8 8 14 26 30 30 9999 9999 9999 9999 9999
 9999 9999 9999 34 26 30 30 9999

ACCUMULATED AWAYNESS :

34 36 34 36 0 36 26 26 26 18 14 16
 16 18 20 22 24 26 26 30 6 18 14 16
 16 18 20 22 24 30 36 36 36 34 36 36
 38 34 36 36 38 36 38 26 30 30 20 16
 18 18 20 22 24 26 32 26 30 30 32 32
 4 32 32 34 30 32 32 34 30 32 32 34
 0 32 32 34 24 26 26 16 12 14 14 16
 18 20 22 26 24 26 26 26

CUBE (95 , 55) IS SELECTED.

CUBE (4 , 23) HAS BEEN FOUND.

CHECKING POINTS

5 7 4 32 9999 3 30 30 30 9999 9999 9999
 9999 9999 9999 9999 9999 17 17 9999 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 6 12 20 9999 9999 9999
 9999 4 6 12 10 8 20 36 36 36 9999 9999
 9999 9999 9999 9999 9999 9999 36 18 36 36 9999 2
 9999 8 8 14 14 18 18 26 20 14 14 8
 9999 8 8 14 26 30 30 9999 9999 9999 9999 9999
 9999 9999 9999 34 26 30 30 9999

ACCUMULATED AWAYNESS :

38	40	36	38	0	40	30	32	32	18	14	16
16	18	20	22	24	32	32	30	6	18	14	16
16	18	20	22	24	30	38	38	40	34	36	36
38	38	40	40	42	38	40	30	32	32	20	16
18	18	20	22	24	26	36	30	32	32	32	36
4	38	38	40	34	36	36	38	38	40	40	42
0	38	38	40	30	32	32	16	12	14	14	16
18	20	22	36	30	32	32	28				

CUBE (4 , 23) IS SELECTED.

CUBE (8 , 27) HAS BEEN FOUND.

CHECKING POINTS

6	8	4	32	9999	3	32	32	32	9999	9999	9999
9999	9999	9999	9999	9999	18	18	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	8	16	28	9999	9999	9999
9999	9999	9999	9999	9999	12	28	44	44	44	9999	9999
9999	9999	9999	9999	9999	9999	9999	22	44	44	9999	2
9999	8	8	14	14	18	18	26	22	16	16	10
9999	8	8	14	28	30	30	9999	9999	9999	9999	9999
9999	9999	9999	36	28	30	30	9999				

ACCUMULATED AWAYNESS :

42	44	42	44	0	44	34	36	36	18	14	16
16	18	20	22	24	36	36	30	6	18	14	16
16	18	20	22	24	30	40	40	42	34	36	36
38	38	40	40	42	40	42	32	34	34	20	16
18	18	20	22	24	26	36	32	34	34	32	40
4	44	44	46	42	44	44	46	42	44	44	46
0	44	44	46	36	38	38	16	12	14	14	16
18	20	22	40	36	38	38	28				

CUBE (8 , 27) IS SELECTED.

CUBE (14 , 19) HAS BEEN FOUND.

CHECKING POINTS

10	12	4	40	9999	3	34	34	40	9999	9999	9999
9999	9999	9999	9999	9999	19	22	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	10	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	32	32	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	26	32	9999	9999	2
9999	6	10	16	14	18	20	28	22	16	18	12
9999	6	10	16	28	30	32	9999	9999	9999	9999	9999
9999	9999	9999	38	28	30	32	9999				

ACCUMULATED AWAYNESS :

44	46	46	46	0	46	38	40	38	18	14	16
16	18	20	22	24	40	38	30	6	18	14	16
16	18	20	22	24	30	42	40	42	34	36	36
38	38	40	40	42	40	42	34	36	34	20	16
18	18	20	22	24	26	36	34	36	34	32	44
4	50	48	50	48	50	48	50	48	50	48	50
0	50	48	50	42	44	42	16	12	14	14	16
18	20	22	44	42	44	42	28				

CUBE (14 , 19) IS SELECTED.

CUBE (28 , 19) HAS BEEN FOUND.

CHECKING POINTS

11	13	4	42	9999	5	42	42	48	9999	9999	9999
9999	9999	9999	9999	9999	23	26	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	12	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	2
9999	8	10	16	14	18	20	28	22	16	18	12
9999	8	10	16	30	32	34	9999	9999	9999	9999	9999
9999	9999	9999	40	30	32	34	9999				

ACCUMULATED AWAYNESS :

48	50	50	50	0	50	40	42	40	18	14	16
16	18	20	22	24	42	40	30	6	18	14	16
16	18	20	22	24	30	44	40	42	34	36	36
38	38	40	40	42	40	42	34	36	34	20	16
18	18	20	22	24	26	36	34	36	34	32	48
4	56	54	56	54	56	54	56	54	56	54	56
0	56	54	56	46	48	46	16	12	14	14	16
18	20	22	48	46	48	46	28				

CUBE (28 , 19) IS SELECTED.

CUBE (32 , 114) HAS BEEN FOUND.

CHECKING POINTS

9999	17	4	44	9999	5	9999	50	9999	9999	9999	9999
9999	9999	9999	9999	9999	23	27	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	12	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	16	9999	24	22	20	28	30	30	18	26	14
9999	8	12	16	9999	40	9999	9999	9999	9999	9999	9999
9999	9999	9999	42	32	32	36	9999				

ACCUMULATED AWAYNESS :

48	52	54	54	0	56	40	44	40	18	14	16
16	18	20	22	24	48	44	30	6	18	14	16
16	18	20	22	24	30	48	40	42	34	36	36
38	38	40	40	42	40	42	34	36	34	20	16
18	18	20	22	24	26	36	34	36	34	32	48
4	58	54	58	56	60	56	60	56	60	56	60
0	62	58	62	46	50	46	16	12	14	14	16
18	20	22	52	50	54	50	28				

CUBE (32 , 114) IS SELECTED.

CUBE (97 , 126) HAS BEEN FOUND.

CHECKING POINTS

9999	21	6	52	9999	7	9999	58	9999	9999	9999	9999
9999	9999	9999	9999	9999	27	31	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	12	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

48	54	56	56	0	58	40	46	40	18	14	16
116	18	20	22	24	50	46	30	6	18	14	16
16	18	20	22	24	30	52	40	42	34	36	36
38	38	40	40	42	40	42	34	36	34	20	16
18	18	20	22	24	26	36	34	36	34	32	48
4	58	54	58	56	60	56	60	56	60	56	60
0	62	58	62	46	50	46	16	12	14	14	16
18	20	22	52	50	54	50	28				

CUBE (97 , 126) IS SELECTED.

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 5

CUBE (49 , 22) HAS BEEN FOUND.

CUBE (49 , 118) HAS BEEN FOUND.

WATCH FOR 22 AND 118

CHECKING POINTS

9999	23	6	54	9999	7	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	29	33	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	12	9999	9999	9999	9999	9999


```

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999

```

ACCUMULATED AWAYNESS :

```

48 56 58 58 0 62 40 46 40 18 14 16
16 18 20 22 24 52 48 30 6 18 14 16
16 18 20 22 24 30 56 40 42 34 36 36
38 38 40 40 42 40 42 34 36 34 20 16
18 18 20 22 24 26 36 34 36 34 32 48
4 58 54 58 56 60 56 60 56 60 56 60
0 62 58 62 46 50 46 16 12 14 14 16
18 20 22 52 50 54 50 28

```

CUBE (49 , 118) IS SELECTED.

CUBE (47 , 23) HAS BEEN FOUND.

CUBE (47 , 119) HAS BEEN FOUND.

WATCH FOR 23 AND 119

CHECKING POINTS

```

9999 25 10 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 31 30 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 12 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999

```

ACCUMULATED AWAYNESS :

```

48 56 60 58 0 62 40 46 40 18 14 16
16 18 20 22 24 54 50 30 6 18 14 16
16 18 20 22 24 30 62 40 42 34 36 36
38 38 40 40 42 40 42 34 36 34 20 16
18 18 20 22 24 26 36 34 36 34 32 48
4 58 54 58 56 60 56 60 56 60 56 60
0 62 58 62 46 50 46 16 12 14 14 16
18 20 22 52 50 54 50 28

```

CUBE (47 , 119) IS SELECTED.

CUBE (37 , 23) HAS BEEN FOUND.

CUBE (37 , 119) HAS BEEN FOUND.

WATCH FOR 23 AND 119

CHECKING POINTS

9999	27	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	33	35	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	12	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

48	60	60	58	0	62	40	46	40	18	14	16
16	18	20	22	24	56	54	30	6	18	14	16
16	18	20	22	24	30	66	40	42	34	36	36
36	36	40	40	42	40	42	34	36	34	20	16
18	18	20	22	24	26	36	34	36	34	32	48
4	58	54	58	56	60	56	60	56	60	56	60
0	62	58	62	46	50	46	16	12	14	14	16
18	20	22	52	50	54	50	28				

CUBE (37 , 119) IS SELECTED.

CUBE (1 , 23) HAS BEEN FOUND.

CUBE (1 , 27) HAS BEEN FOUND.

WATCH FOR 23 AND 27

CHECKING POINTS

9999	27	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	33	35	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

48	62	60	58	0	62	40	46	40	18	14	16
16	18	20	22	24	60	60	30	6	18	14	16
16	18	20	22	24	30	66	40	42	34	36	36
36	36	40	40	42	40	42	34	36	34	20	16
18	18	20	22	24	26	36	34	36	34	32	48
4	58	54	58	56	60	56	60	56	60	56	60
0	62	58	62	46	50	46	16	12	14	14	16
18	20	22	52	50	54	50	28				

CUBE (1 , 23) IS SELECTED.
 CUBE (35 , 18) HAS BEEN FOUND.
 CUBE (35 , 23) HAS BEEN FOUND.
 WATCH FOR 18 AND 23
 CUBE (35 , 114) HAS BEEN FOUND.
 WATCH FOR 23 AND 114
 CUBE (35 , 119) HAS BEEN FOUND.
 WATCH FOR 23 AND 119

CHECKING POINTS

```

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 33 35 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
    
```

ACCUMULATED AWAYNESS :

48	62	60	58	0	62	40	46	40	18	14	16
16	18	20	22	24	64	64	30	6	18	14	16
16	18	20	22	24	30	66	40	42	34	36	36
38	38	40	40	42	40	42	34	36	34	20	16
18	18	20	22	24	26	36	34	36	34	32	48
4	58	54	58	56	60	56	60	56	60	56	60
0	62	58	62	46	50	46	16	12	14	14	16
18	20	22	52	50	54	50	28				

CUBE (35 , 119) IS SELECTED.
 CUBE (61 , 17) HAS BEEN FOUND.
 CUBE (61 , 23) HAS BEEN FOUND.
 WATCH FOR 17 AND 23
 CUBE (61 , 113) HAS BEEN FOUND.
 WATCH FOR 23 AND 113
 CUBE (61 , 119) HAS BEEN FOUND.

ACCUMULATED AWAYNESS :

48	62	60	58	0	62	40	46	40	18	14	16
16	18	20	22	24	64	66	50	6	18	14	16
16	18	20	22	24	30	66	40	42	34	36	36
38	38	40	40	42	40	42	34	36	34	20	16
18	18	20	22	24	26	36	34	36	34	32	48
4	58	54	56	56	60	56	60	58	60	56	60
0	62	58	62	46	50	46	16	12	14	14	16
18	20	22	52	50	54	50	28				

CUBE (62 , 119) IS SELECTED.

*9

**FORTRAN ** STOP LOGOFF

/F ---

%CAN-

/LOGOFF BUT

%C E420 LOGOFF AT 2024 ON 08/09/78, FOR TSN 9003.

%C E421 CPU TIME USED: 000200.9825 SECONDS.

%C E222 PLEASE LOGON.

APPENDIX H

EXAMPLE OF 9 VARIABLES, 120 ON-SET MINTERMS, AND 40 DON'T
CARE MINTERMS.

TOATAL NUMBER OF ON-SET AND DONT CARES = 160
 ON_SET = 120 DONT CARES = 40
 NUMBERS OF VARIABLES = 10

A96

MINT.# 1 = 1	MINT.# 2 = 2	MINT.# 3 = 3
MINT.# 4 = 4	MINT.# 5 = 5	MINT.# 6 = 6
MINT.# 7 = 7	MINT.# 8 = 9	MINT.# 9 = 11
MINT.# 10 = 12	MINT.# 11 = 13	MINT.# 12 = 14
MINT.# 13 = 16	MINT.# 14 = 17	MINT.# 15 = 18
MINT.# 16 = 19	MINT.# 17 = 21	MINT.# 18 = 22
MINT.# 19 = 23	MINT.# 20 = 24	MINT.# 21 = 26
MINT.# 22 = 27	MINT.# 23 = 28	MINT.# 24 = 29
MINT.# 25 = 30	MINT.# 26 = 31	MINT.# 27 = 34
MINT.# 28 = 35	MINT.# 29 = 36	MINT.# 30 = 37
MINT.# 31 = 38	MINT.# 32 = 39	MINT.# 33 = 41
MINT.# 34 = 42	MINT.# 35 = 43	MINT.# 36 = 45
MINT.# 37 = 48	MINT.# 38 = 49	MINT.# 39 = 50
MINT.# 40 = 51	MINT.# 41 = 52	MINT.# 42 = 53
MINT.# 43 = 55	MINT.# 44 = 56	MINT.# 45 = 57
MINT.# 46 = 58	MINT.# 47 = 59	MINT.# 48 = 60
MINT.# 49 = 61	MINT.# 50 = 62	MINT.# 51 = 63
MINT.# 52 = 64	MINT.# 53 = 68	MINT.# 54 = 69
MINT.# 55 = 70	MINT.# 56 = 71	MINT.# 57 = 72
MINT.# 58 = 75	MINT.# 59 = 76	MINT.# 60 = 77
MINT.# 61 = 78	MINT.# 62 = 79	MINT.# 63 = 80
MINT.# 64 = 81	MINT.# 65 = 90	MINT.# 66 = 91
MINT.# 67 = 92	MINT.# 68 = 93	MINT.# 69 = 94
MINT.# 70 = 95	MINT.# 71 = 96	MINT.# 72 = 97
MINT.# 73 = 98	MINT.# 74 = 103	MINT.# 75 = 104
MINT.# 76 = 105	MINT.# 77 = 106	MINT.# 78 = 107
MINT.# 79 = 108	MINT.# 80 = 109	MINT.# 81 = 111
MINT.# 82 = 112	MINT.# 83 = 113	MINT.# 84 = 114
MINT.# 85 = 115	MINT.# 86 = 116	MINT.# 87 = 117
MINT.# 88 = 121	MINT.# 89 = 122	MINT.# 90 = 123
MINT.# 91 = 124	MINT.# 92 = 145	MINT.# 93 = 126
MINT.# 94 = 127	MINT.# 95 = 128	MINT.# 96 = 129
MINT.# 97 = 150	MINT.# 98 = 151	MINT.# 99 = 152
MINT.# 100 = 153	MINT.# 101 = 156	MINT.# 102 = 157
MINT.# 103 = 158	MINT.# 104 = 159	MINT.# 105 = 160
MINT.# 106 = 161	MINT.# 107 = 162	MINT.# 108 = 163
MINT.# 109 = 164	MINT.# 110 = 165	MINT.# 111 = 167
MINT.# 112 = 168	MINT.# 113 = 170	MINT.# 114 = 172
MINT.# 115 = 173	MINT.# 116 = 174	MINT.# 117 = 175
MINT.# 118 = 176	MINT.# 119 = 178	MINT.# 120 = 179
MINT.# 121 = 185	MINT.# 122 = 186	MINT.# 123 = 187
MINT.# 124 = 188	MINT.# 125 = 189	MINT.# 126 = 190
MINT.# 127 = 191	MINT.# 128 = 192	MINT.# 129 = 201
MINT.# 130 = 202	MINT.# 131 = 203	MINT.# 132 = 204
MINT.# 133 = 205	MINT.# 134 = 206	MINT.# 135 = 207
MINT.# 136 = 213	MINT.# 137 = 214	MINT.# 138 = 215
MINT.# 139 = 216	MINT.# 140 = 217	MINT.# 141 = 220
MINT.# 142 = 223	MINT.# 143 = 224	MINT.# 144 = 225
MINT.# 145 = 226	MINT.# 146 = 227	MINT.# 147 = 228
MINT.# 148 = 229	MINT.# 149 = 230	MINT.# 150 = 245
MINT.# 151 = 246	MINT.# 152 = 247	MINT.# 153 = 248
MINT.# 154 = 249	MINT.# 155 = 256	MINT.# 156 = 257
MINT.# 157 = 258	MINT.# 158 = 259	MINT.# 159 = 254
MINT.# 160 = 260	MINT.#	

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 3-CUBE = (81 , 48) = 000--1000-

THIS CUBE COVERS THESE MINTERMS :
 81 48 16 17 49 80 112 113

THIS MINTERM MT(64) = 81 IS A CORE POINT

THE 2-TH CUBE COVERING IS :
 4-CUBE = (90 , 63) = 000--11-1-

THIS CUBE COVERS THESE MINTERMS :
 90 63 26 27 30 31 58 59 62 91 94 95 122 123 126 127

THIS MINTERM MT(65) = 90 IS A CORE POINT

THE 3-TH CUBE COVERING IS :
 4-CUBE = (156 , 63) = 00-0-111--

THIS CUBE COVERS THESE MINTERMS :
 156 63 28 29 30 31 60 61 62 157 158 159 188 189 190 191

THE 4-TH CUBE COVERING IS :
 3-CUBE = (38 , 5) = 0000-001--

THIS CUBE COVERS THESE MINTERMS :
 38 5 4 6 7 36 37 39

THE 5-TH CUBE COVERING IS :
 3-CUBE = (162 , 51) = 00-01-001-

THIS CUBE COVERS THESE MINTERMS :
 162 51 34 35 50 163 178 179

THE 6-TH CUBE COVERING IS :
 3-CUBE = (41 , 123) = 000-1-10-1

THIS CUBE COVERS THESE MINTERMS :
 41 123 43 57 59 105 107 121

THE 7-TH CUBE COVERING IS :
 3-CUBE = (69 , 78) = 000100-1--

THIS CUBE COVERS THESE MINTERMS :
 69 78 68 70 71 76 77 79

THE 8-TH CUBE COVERING IS :
 3-CUBE = (164 , 225) = 001-100-0-

THIS CUBE COVERS THESE MINTERMS :
 164 225 160 161 165 224 228 229

THE 9-TH CUBE COVERING IS :
 3-CUBE = (150 , 31) = 00-001-11-

THIS CUBE COVERS THESE MINTERMS :
 150 31 22 23 30 151 158 159

THE 10-TH CUBE COVERING IS :
 3-CUBE = (174 , 189) = 00101-11--

THIS CUBE COVERS THESE MINTERMS :
 174 189 172 173 175 188 190 191

THE 11-TH CUBE COVERING IS :
 3-CUBE = (116 , 49) = 000-110-0-

THIS CUBE COVERS THESE MINTERMS :
 116 49 48 52 53 112 113 117

THE 12-TH CUBE COVERING IS :
 3-CUBE = (2 , 51) = 0000--001-

THIS CUBE COVERS THESE MINTERMS :
 2 51 3 18 19 34 35 50

THE 13-TH CUBE COVERING IS :
 3-CUBE = (170 , 50) = 00-01--010

THIS CUBE COVERS THESE MINTERMS :
 170 50 34 42 58 162 178 186

THE 14-TH CUBE COVERING IS :
 3-CUBE = (12 , 93) = 000-0-110-

THIS CUBE COVERS THESE MINTERMS :
 12 93 13 28 29 76 77 92

THE 15-TH CUBE COVERING IS :
 3-CUBE = (75 , 59) = 000---1011

THIS CUBE COVERS THESE MINTERMS :
 75 59 11 27 43 91 107 123

THE 16-TH CUBE COVERING IS :
 3-CUBE = (98 , 58) = 000-1--010

THIS CUBE COVERS THESE MINTERMS :
 98 58 34 42 50 106 114 122

THE 17-TH CUBE COVERING IS :
 3-CUBE = (45 , 21) = 0000---101

THIS CUBE COVERS THESE MINTERMS :
 45 21 5 13 29 37 53 61

THE 18-TH CUBE COVERING IS ;
 3-CUBE = (14 , 68) = 000-00-1-0

THIS CUBE COVERS THESE MINTERMS ;
 14 68 4 6 12 70 76 78

THE 19-TH CUBE COVERING IS ;
 3-CUBE = (24 , 62) = 0000-11--0

THIS CUBE COVERS THESE MINTERMS ;
 24 62 26 28 30 56 58 60

THE 20-TH CUBE COVERING IS ;
 3-CUBE = (1 , 23) = 00000-0--1

THIS CUBE COVERS THESE MINTERMS ;
 1 23 3 5 7 17 19 21

THE 21-TH CUBE COVERING IS ;
 3-CUBE = (111 , 91) = 0001--1-11

THIS CUBE COVERS THESE MINTERMS ;
 111 91 75 79 95 107 123 127

THE 22-TH CUBE COVERING IS ;
 3-CUBE = (124 , 30) = 000--111-0

THIS CUBE COVERS THESE MINTERMS ;
 124 30 28 60 62 92 94 126

THE 23-TH CUBE COVERING IS ;
 3-CUBE = (115 , 48) = 000-1100--

THIS CUBE COVERS THESE MINTERMS ;
 115 48 49 50 51 112 113 114

THE 24-TH CUBE COVERING IS :

$$3\text{-CUBE} = (55 , 17) = 0000-10--1$$

THIS CUBE COVERS THESE MINTERMS :

55 17 19 21 23 49 51 53

THE 25-TH CUBE COVERING IS :

$$2\text{-CUBE} = (128 , 161) = 0010-0000-$$

THIS CUBE COVERS THESE MINTERMS :

128 161 129 160

THE 26-TH CUBE COVERING IS :

$$2\text{-CUBE} = (103 , 7) = 000--00111$$

THIS CUBE COVERS THESE MINTERMS :

103 7 39 71

THE 27-TH CUBE COVERING IS :

$$2\text{-CUBE} = (72 , 96) = 0001-0-000$$

THIS CUBE COVERS THESE MINTERMS :

72 96 64 104

THE 28-TH CUBE COVERING IS :

$$2\text{-CUBE} = (168 , 162) = 001010-0-0$$

THIS CUBE COVERS THESE MINTERMS :

168 162 160 170

THE 29-TH CUBE COVERING IS :

$$2\text{-CUBE} = (9 , 5) = 000000--01$$

THIS CUBE COVERS THESE MINTERMS :

9 5 1 13

THE 30-TH CUBE COVERING IS :
 2-CUBE = (145 , 1) = 00-00-0001

THIS CUBE COVERS THESE MINTERMS :
 145 1 17 129

THE 31-TH CUBE COVERING IS :
 2-CUBE = (152 , 217) = 001-01100-

THIS CUBE COVERS THESE MINTERMS :
 152 217 153 216

THE 32-TH CUBE COVERING IS :
 2-CUBE = (108 , 105) = 0001101-0-

THIS CUBE COVERS THESE MINTERMS :
 108 105 104 109

THE 33-TH CUBE COVERING IS :
 2-CUBE = (176 , 50) = 00-01100-0

THIS CUBE COVERS THESE MINTERMS :
 176 50 48 178

THE 34-TH CUBE COVERING IS :
 2-CUBE = (167 , 37) = 00-01001-1

THIS CUBE COVERS THESE MINTERMS :
 167 37 39 165

THE 35-TH CUBE COVERING IS :
 2-CUBE = (97 , 224) = 00-110000-

THIS CUBE COVERS THESE MINTERMS :
 97 224 96 225

0	0	2	2	2	2	4	2	2	4	4	4
4	6	6	2	4	4	6	6	4	6	6	6
8	8	0	0	4	4	4	4	6	6	2	2
4	6	4	4	6	6	6	6	8	0	0	2
2	2	2	2	4	4	4	2	6	6	4	4
6	6	4	4	6	6	8	8	4	4	6	6
6	6	8	6	8	8	8	10	10	2	4	4
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (81 , 48) IS SELECTED,

CUBE (90 , 63) HAS BEEN FOUND,

CHECKING POINTS

8	2	4	2	2	2	2	0	4	1	2	2
9999	9999	16	16	5	5	10	8	9999	9999	10	10
9999	9999	4	4	1	2	0	2	2	4	8	2
9999	9999	16	16	5	10	5	8	16	9999	9999	10
10	9999	9999	4	1	0	0	1	0	4	2	2
8	8	9999	9999	9999	9999	10	5	9999	9999	8	4
4	0	4	4	8	8	1	2	8	9999	9999	16
8	5	5	16	9999	9999	10	2	9999	9999	0	1
1	1	2	4	2	2	8	8	2	2	0	0
0	0	0	0	2	0	0	1	1	2	4	2
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999								

ACCUMULATED AWAYNESS :

8	8	8	10	10	10	10	8	8	10	10	10
0	0	4	4	6	6	6	4	4	4	6	6
6	6	8	8	10	10	10	10	8	8	8	10
0	0	4	4	6	6	6	4	4	4	4	6
6	6	6	8	10	10	10	10	8	8	10	10
10	10	0	0	4	4	6	6	6	6	8	8
8	10	8	8	8	8	10	10	10	0	0	4
4	6	6	4	4	4	6	8	6	6	12	12
10	10	8	8	10	10	10	10	12	12	12	12
14	14	14	12	12	14	14	14	14	8	8	8
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (90 , 63) IS SELECTED,

ALL CORE CUBES HAVE BEEN SEARCHED,
AND THE MAXIMUM CUBE DIMENSION IS : 5

CUBE (156 , 63) HAS BEEN FOUND.

CHECKING POINTS

8	2	4	2	2	2	2	0	5	3	4	4
9999	9999	17	17	9	9	14	12	9999	9999	9999	9999
9999	9999	4	4	1	2	0	2	2	5	9	4
9999	9999	17	17	9	14	9	12	20	9999	9999	9999
9999	9999	9999	4	1	0	0	1	0	4	2	2
9	9	9999	9999	9999	9999	14	7	9999	9999	8	4
4	0	4	4	8	8	1	2	9	9999	9999	16
8	5	5	17	9999	9999	14	2	9999	9999	0	1
3	3	4	8	9999	9999	9999	9999	2	2	0	0
0	0	0	0	3	2	2	3	3	2	5	2
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999								

ACCUMULATED AWAYNESS :

14	14	14	14	14	14	14	12	12	12	12	12
0	0	8	8	8	8	8	6	4	4	6	6
6	6	14	14	14	14	14	14	12	12	12	12
0	0	8	8	8	8	8	6	6	4	4	6
6	6	6	16	16	16	16	16	14	14	14	14
14	14	0	0	4	4	8	8	6	6	16	16
16	16	14	14	14	14	14	14	14	0	0	10
10	10	10	8	4	4	8	12	6	6	18	18
12	12	10	10	10	10	10	10	18	18	18	18
18	18	18	16	16	16	16	16	16	12	12	12
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0								

CUBE (156 , 63) IS SELECTED.

CUBE (38 , 3) HAS BEEN FOUND.

CUBE (38 , 5) HAS BEEN FOUND.

WATCH FOR 3 AND 5

CHECKING POINTS

10	4	6	9999	9999	9999	9999	0	5	5	6	6
9999	9999	17	17	11	11	16	12	9999	9999	9999	9999
9999	9999	6	6	9999	9999	9999	9999	2	5	9	6
9999	9999	17	17	11	16	11	12	20	9999	9999	9999
9999	9999	9999	4	3	2	2	3	0	4	2	2
9	9	9999	9999	9999	9999	14	7	9999	9999	8	4
4	1	4	4	8	8	1	2	9	9999	9999	16

8	5	5	17	9999	9999	14	2	9999	9999	0	1
3	3	4	8	9999	9999	9999	9999	2	2	0	0
2	2	2	0	3	2	2	3	3	2	5	2
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999								

ACCUMULATED AWAYNESS :

16	16	16	14	14	14	14	16	16	14	14	14
0	0	12	12	10	10	10	12	4	4	6	6
6	6	16	16	14	14	14	14	16	16	16	14
0	0	12	12	10	10	10	12	12	4	4	6
6	6	6	20	18	18	18	18	20	20	18	18
18	18	0	0	4	4	14	14	6	6	20	20
20	18	20	20	20	20	18	18	18	0	0	16
16	14	14	16	4	4	14	18	6	6	22	22
16	16	18	18	10	10	10	10	22	22	22	22
20	20	20	22	22	20	20	20	20	18	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0								

CUBE (38 , 5) IS SELECTED.

CUBE (162 , 51) HAS BEEN FOUND.

CUBE (162 , 58) HAS BEEN FOUND.

CUBE (162 , 225) HAS BEEN FOUND.

CHECKING POINTS

10	6	8	9999	9999	9999	9999	0	5	5	6	6
9999	9999	19	19	11	11	16	12	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	2	7	11	6
9999	9999	9999	9999	11	16	13	12	20	9999	9999	9999
9999	9999	9999	4	3	2	2	3	0	4	2	2
9	9	9999	9999	9999	9999	14	7	9999	9999	8	4
6	1	4	4	8	8	1	2	9	9999	9999	18
10	5	5	17	9999	9999	14	2	9999	9999	0	1
3	3	4	8	9999	9999	9999	9999	4	4	9999	9999
2	2	4	0	5	2	2	3	3	3	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999								

ACCUMULATED AWAYNESS :

CUBE (41 , 123) IS SELECTED.

CUBE (69 , 6) HAS BEEN FOUND.

CUBE (69 , 12) HAS BEEN FOUND.

CUBE (69 , 78) HAS BEEN FOUND.

CHECKING POINTS

10	6	8	9999	9999	9999	9999	1	7	7	8	8
9999	9999	19	19	11	11	16	12	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9	9999	8
9999	9999	9999	9999	11	16	13	14	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	1	8	9999	9999
9999	9999	9999	9999	9999	9999	16	9	9999	9999	8	6
6	2	6	9999	10	9999	2	6	13	9999	9999	18
12	5	5	9999	9999	9999	14	2	9999	9999	0	1
3	3	4	8	9999	9999	9999	9999	4	4	9999	9999
2	2	4	0	5	2	2	3	3	3	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

28	28	26	14	14	14	14	28	26	30	28	28
0	0	26	24	26	26	24	28	4	4	6	6
6	6	16	16	14	14	14	14	20	26	18	26
0	0	12	12	26	24	22	26	16	4	4	6
6	6	6	34	34	32	32	30	34	30	34	32
32	30	0	0	4	4	32	30	6	6	32	30
30	28	32	26	30	24	32	30	28	0	0	28
26	30	28	22	4	4	30	36	6	6	40	38
36	34	38	36	10	10	10	10	38	36	22	22
38	36	34	38	36	38	36	36	34	36	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (69 , 78) IS SELECTED.

CUBE (164 , 225) HAS BEEN FOUND.

CHECKING POINTS

10	6	8	9999	9999	9999	9999	1	7	7	8	8
9999	9999	19	19	11	11	16	12	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9	9999	8
9999	9999	9999	9999	11	16	13	14	9999	9999	9999	9999

9999	9999	9999	6	9999	9999	9999	9999	1	8	9999	9999
9999	9999	9999	9999	9999	9999	16	9	9999	9999	10	8
6	2	6	9999	10	9999	2	6	13	9999	9999	18
12	5	5	9999	9999	9999	14	2	9999	9999	1	2
3	3	4	8	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	6	1	5	4	4	3	3	4	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999								

ACCUMULATED AWAYNESS :

32	34	32	14	14	14	14	34	34	36	34	36
0	0	34	32	32	34	32	36	4	4	6	6
6	6	16	16	14	14	14	14	20	32	18	30
0	0	12	12	30	28	28	32	16	4	4	6
6	6	6	38	34	32	32	30	40	38	34	32
32	30	0	0	4	4	40	38	6	6	34	32
34	32	36	26	36	24	36	34	34	0	0	34
32	34	32	22	4	4	36	40	6	6	42	40
42	40	44	42	10	10	10	10	38	36	22	22
38	36	36	40	40	40	38	40	38	38	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (164 , 225) IS SELECTED.

CUBE (150 , 31) HAS BEEN FOUND.

CHECKING POINTS

10	6	8	9999	9999	9999	9999	1	7	7	8	10
9999	9999	21	21	13	9999	9999	12	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9	9999	8
9999	9999	9999	9999	11	16	15	14	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	1	8	9999	9999
9999	9999	9999	9999	9999	9999	16	9	9999	9999	10	8
6	2	6	9999	10	9999	2	6	13	9999	9999	18
12	5	5	9999	9999	9999	14	2	9999	9999	1	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	6	1	5	4	4	3	3	4	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999								

ACCUMULATED AWAYNESS :

38	38	36	14	14	14	14	40	38	40	38	38
----	----	----	----	----	----	----	----	----	----	----	----

0	0	36	34	34	34	32	40	4	4	6	6
6	6	16	16	14	14	14	14	20	38	18	36
0	0	12	12	34	32	30	38	16	4	4	6
6	6	6	46	34	32	32	30	48	44	34	32
32	30	0	0	4	4	44	42	6	6	44	42
42	38	46	26	44	24	44	42	40	0	0	40
38	40	38	22	4	4	42	44	6	6	48	46
42	40	48	46	10	10	10	10	38	36	22	22
38	36	40	48	46	46	44	44	42	44	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0								

CUBE (150 , 31) IS SELECTED.

CUBE (174 , 189) HAS BEEN FOUND.

CHECKING POINTS

10	6	8	9999	9999	9999	9999	1	7	7	8	10
9999	9999	21	21	13	9999	9999	12	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9	9999	10
9999	9999	9999	9999	11	16	15	14	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	1	8	9999	9999
9999	9999	9999	9999	9999	9999	16	9	9999	9999	10	8
6	2	6	9999	10	9999	2	6	13	9999	9999	18
12	5	5	9999	9999	9999	14	2	9999	9999	1	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	2	7	9999	9999	9999	9999	4	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

46	46	44	14	14	14	14	46	44	44	42	42
0	0	44	42	40	34	32	46	4	4	6	6
6	6	16	16	14	14	14	14	20	42	18	38
0	0	12	12	38	36	34	42	16	4	4	6
6	6	6	56	34	32	32	30	56	52	34	32
32	30	0	0	4	4	50	48	6	6	52	50
50	44	52	26	50	24	48	46	44	0	0	48
46	46	44	22	4	4	46	50	6	6	54	52
42	40	52	50	10	10	10	10	38	36	22	22
38	36	42	50	48	46	44	44	42	48	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0								

CUBE (174 , 189) IS SELECTED.

CUBE (116 , 49) HAS BEEN FOUND.

CHECKING POINTS

10	6	8	9999	9999	9999	9999	1	7	7	8	10
9999	9999	21	21	15	9999	9999	12	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9	9999	10
9999	9999	9999	9999	9999	9999	17	16	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	1	8	9999	9999
9999	9999	9999	9999	9999	9999	16	9	9999	9999	12	10
6	2	6	9999	10	9999	2	6	13	9999	9999	20
14	9999	9999	9999	9999	9999	16	2	9999	9999	1	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	2	7	9999	9999	9999	9999	5	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

50	52	50	14	14	14	14	52	52	50	48	50
0	0	48	46	42	34	32	50	4	4	6	6
6	6	16	16	14	14	14	14	20	48	18	42
0	0	12	12	38	36	36	44	16	4	4	6
6	6	6	60	34	32	32	30	62	60	34	32
32	30	0	0	4	4	54	52	6	6	54	52
54	48	56	26	56	24	52	50	50	0	0	50
48	46	44	22	4	4	48	54	6	6	60	58
42	40	58	56	10	10	10	10	38	36	22	22
38	36	48	56	56	46	44	44	42	50	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (116 , 49) IS SELECTED.

CUBE (2 , 23) HAS BEEN FOUND.

CUBE (2 , 39) HAS BEEN FOUND.

CUBE (2 , 51) HAS BEEN FOUND.

WATCH FOR 23 AND 51

CHECKING POINTS

12	9999	9999	9999	9999	9999	9999	1	9	7	8	10
9999	9999	9999	9999	15	9999	9999	12	9999	9999	9999	9999

9999	9999	9999	9999	9999	9999	9999	9999	9999	11	9999	10
9999	9999	9999	9999	9999	9999	19	16	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	1	8	9999	9999
9999	9999	9999	9999	9999	9999	16	9	9999	9999	12	10
8	2	6	9999	10	9999	2	6	13	9999	9999	22
16	9999	9999	9999	9999	9999	16	2	9999	9999	1	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	2	7	9999	9999	9999	9999	5	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

52	52	50	14	14	14	14	56	54	56	54	54
0	0	48	46	46	34	32	54	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	48
0	0	12	12	38	36	38	48	16	4	4	6
6	6	6	64	34	32	32	30	68	64	34	32
32	30	0	0	4	4	62	60	6	6	58	56
56	52	62	26	60	24	60	58	56	0	0	52
50	46	44	22	4	4	56	58	6	6	64	62
42	40	64	62	10	10	10	10	38	36	22	22
38	36	52	62	60	46	44	44	42	54	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (2 , 51) IS SELECTED,

CUBE (170 , 50) HAS BEEN FOUND.

CHECKING POINTS

12	9999	9999	9999	9999	9999	9999	1	9	7	8	10
9999	9999	9999	9999	15	9999	9999	12	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	10
9999	9999	9999	9999	9999	9999	19	18	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	1	8	9999	9999
9999	9999	9999	9999	9999	9999	16	9	9999	9999	12	10
10	2	6	9999	12	9999	2	6	13	9999	9999	24
16	9999	9999	9999	9999	9999	16	2	9999	9999	1	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	3	9999	9999	9999	9999	9999	6	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

58	52	50	14	14	14	14	62	58	62	62	58
0	0	48	46	54	34	32	58	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	54
0	0	12	12	38	36	42	50	16	4	4	6
6	6	6	70	34	32	32	30	74	70	34	32
32	30	0	0	4	4	70	70	6	6	62	62
58	58	66	26	62	24	66	66	62	0	0	54
54	46	44	22	4	4	62	64	6	6	68	68
42	40	68	68	10	10	10	10	38	36	22	22
38	36	56	64	60	46	44	44	42	56	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (170 , 50) IS SELECTED.

CUBE (12 , 69) HAS BEEN FOUND.

CUBE (12 , 70) HAS BEEN FOUND.

WATCH FOR 69 AND 70

CUBE (12 , 93) HAS BEEN FOUND.

CUBE (12 , 94) HAS BEEN FOUND.

CHECKING POINTS

12	9999	9999	9999	9999	9999	9999	2	9	9999	9999	12
9999	9999	9999	9999	17	9999	9999	14	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	12
9999	9999	9999	9999	9999	9999	19	18	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	2	8	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	12	10
10	2	6	9999	12	9999	3	7	13	9999	9999	24
16	9999	9999	9999	9999	9999	18	2	9999	9999	1	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	3	9999	9999	9999	9999	9999	6	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

62	52	50	14	14	14	14	64	62	62	62	60
0	0	48	46	56	34	32	60	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	56
0	0	12	12	38	36	48	54	16	4	4	6
6	6	6	74	34	32	32	30	76	74	34	32
32	30	0	0	4	4	70	70	6	6	68	68

66	64	70	26	68	24	68	68	66	0	0	62
62	46	44	22	4	4	64	70	6	6	74	74
42	40	72	72	10	10	10	10	38	36	22	22
38	36	64	70	60	46	44	44	42	64	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (12 , 93) IS SELECTED,

CUBE (75 , 59) HAS BEEN FOUND,

CUBE (75 , 127) HAS BEEN FOUND,

WATCH FOR 59 AND 127

CHECKING POINTS

12	9999	9999	9999	9999	9999	9999	9999	3	9999	9999	9999	12
9999	9999	9999	9999	17	9999	9999	14	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	12
9999	9999	9999	9999	9999	9999	9999	19	18	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	2	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	12	10
10	2	6	9999	14	9999	3	7	15	9999	9999	24	24
18	9999	9999	9999	9999	9999	18	2	9999	9999	1	2	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	3	9999	9999	9999	9999	9999	6	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

66	52	50	14	14	14	14	66	62	62	62	64
0	0	48	46	62	34	32	64	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	60
0	0	12	12	38	36	52	58	16	4	4	6
6	6	6	80	34	32	32	30	80	74	34	32
32	30	0	0	4	4	70	70	6	6	74	72
70	68	74	26	70	24	74	72	68	0	0	66
64	46	44	22	4	4	70	76	6	6	82	80
42	40	78	76	10	10	10	10	38	36	22	22
38	36	70	76	60	46	44	44	42	72	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (75 , 59) IS SELECTED,

CUBE (98 , 58) HAS BEEN FOUND.

CHECKING POINTS

12	9999	9999	9999	9999	9999	9999	9999	3	9999	9999	9999	12
9999	9999	9999	9999	17	9999	9999	14	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	12
9999	9999	9999	9999	9999	9999	9999	19	20	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	2	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	14	10
9999	2	7	9999	9999	9999	3	7	15	9999	9999	9999	9999
20	9999	9999	9999	9999	9999	18	2	9999	9999	1	2	
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	3	9999	9999	9999	9999	9999	6	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

72	52	50	14	14	14	14	72	62	62	62	68	
0	0	48	46	70	34	32	68	4	4	6	6	
6	6	16	16	14	14	14	14	20	50	18	66	
0	0	12	12	38	36	56	60	16	4	4	6	
6	6	6	84	34	32	32	30	84	74	34	32	
32	30	0	0	4	4	70	70	6	6	76	76	
70	72	76	26	70	24	78	78	72	0	0	66	
66	46	44	22	4	4	74	84	6	6	88	88	
42	40	84	84	10	10	10	10	38	36	22	22	
38	36	76	80	60	46	44	44	42	76	18	18	
0	0	0	0	0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	0	0	0	0	

CUBE (98 , 58) IS SELECTED,

CUBE (45 , 21) HAS BEEN FOUND.

CHECKING POINTS

14	9999	9999	9999	9999	9999	9999	9999	4	9999	9999	9999	12
9999	9999	9999	9999	9999	9999	9999	9999	14	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	21	20	9999	9999	9999	9999
9999	9999	9999	6	9999	9999	9999	9999	9999	2	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	14	10
9999	2	7	9999	9999	9999	3	8	15	9999	9999	9999	9999
20	9999	9999	9999	9999	9999	18	2	9999	9999	9999	1	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	3	9999	9999	9999	9999	9999	6	9999	9999	9999

9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

74	52	50	14	14	14	14	74	62	62	62	72
0	0	48	46	70	34	32	72	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	66
0	0	12	12	38	36	58	64	16	4	4	6
6	6	6	90	34	32	32	30	90	74	34	32
32	30	0	0	4	4	70	70	6	6	82	80
70	76	62	26	70	24	82	80	76	0	0	66
72	46	44	22	4	4	78	88	6	6	94	92
42	40	90	88	10	10	10	10	38	36	22	22
38	36	80	86	60	46	44	44	42	82	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (45 , 21) IS SELECTED.

CUBE (14 , 68) HAS BEEN FOUND.

CUBE (14 , 92) HAS BEEN FOUND.

WATCH FOR 68 AND 92

CHECKING POINTS

14	9999	9999	9999	9999	9999	9999	9999	4	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	14	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	21	20	9999	9999	9999	9999
9999	9999	9999	7	9999	9999	9999	9999	3	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	14	10
9999	2	7	9999	9999	9999	4	8	15	9999	9999	9999	9999
20	9999	9999	9999	9999	9999	18	2	9999	9999	1	2	2
9999	9999	4	8	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	3	9999	9999	9999	9999	9999	6	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

78	52	50	14	14	14	14	78	62	62	62	72
0	0	48	46	70	34	32	76	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	66

CUBE (24 , 62) IS SELECTED.

CUBE (1 , 23) HAS BEEN FOUND.

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	5	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	23	9999	9999	9999	9999	9999
9999	9999	9999	7	9999	9999	9999	9999	3	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	14	10
9999	2	7	9999	9999	9999	4	8	15	9999	9999	9999	9999
20	9999	9999	9999	9999	9999	20	3	9999	9999	1	3	
9999	9999	5	8	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	3	9999	9999	9999	9999	9999	6	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

84	52	50	14	14	14	14	84	62	62	62	72	
0	0	48	46	70	34	32	76	4	4	6	6	
6	6	16	16	14	14	14	14	20	50	18	66	
0	0	12	12	38	36	70	70	16	4	4	6	
6	6	6	102	34	32	32	30	102	74	34	32	
32	30	0	0	4	4	70	70	6	6	98	98	
70	92	98	26	70	24	96	96	92	0	0	66	
90	46	44	22	4	4	92	104	6	6	108	108	
42	40	104	104	10	10	10	10	38	36	22	22	
38	36	98	104	60	46	44	44	42	100	18	18	
0	0	0	0	0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	0	0	0	0	
0	0	0	0	0	0	0	0	0	0	0	0	

CUBE (1 , 23) IS SELECTED.

CUBE (111 , 91) HAS BEEN FOUND.

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	5	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	23	9999	9999	9999	9999	9999
9999	9999	9999	7	9999	9999	9999	9999	3	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	14	10
9999	3	7	9999	9999	9999	4	9	9999	9999	9999	9999	9999
22	9999	9999	9999	9999	9999	20	3	9999	9999	1	3	

```

9999 9999      5      8 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999      8      3 9999 9999 9999 9999 9999      6 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999

```

ACCUMULATED AWAYNESS :

```

84  52  50  14  14  14  14  88  62  62  62  72
0   0  48  46  70  34  32  76  4   4   6   6
6   6  16  16  14  14  14  14  20  50  18  66
0   0  12  12  38  36  74  70  16  4   4   6
6   6   6 108  34  32  32  30 106  74  34  32
32  30   0   0   4   4  70  70   6   6 104 102
70  94 102  26  70  24 100  98  92   0   0   66
92  46  44  22   4   4  96 112   6   6 118 116
42  40 112 110  10  10  10  10  38  36  22  22
38  36 104 112  60  46  44  44  42 110  18  18
0   0   0   0   0   0   0   0   0   0   0   0
0   0   0   0   0   0   0   0   0   0   0   0
0   0   0   0   0   0   0   0   0   0   0   0
0   0   0   0

```

CUBE (111 , 91) IS SELECTED.

CUBE (124 , 30) HAS BEEN FOUND.

CHECKING POINTS

```

9999 9999 9999 9999 9999 9999 9999 9999 5 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 23 9999 9999 9999 9999 9999
9999 9999 9999 7 9999 9999 9999 9999 3 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 14 10
9999 3 7 9999 9999 9999 5 9 9999 9999 9999 9999
22 9999 9999 9999 9999 9999 9999 9999 3 9999 9999 1 3
9999 9999 5 8 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 8 3 9999 9999 9999 9999 9999 6 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
9999 9999 9999 9999

```

ACCUMULATED AWAYNESS :

```

84  52  50  14  14  14  14  94  62  62  62  72
0   0  48  46  70  34  32  76  4   4   6   6
6   6  16  16  14  14  14  14  20  50  18  66
0   0  12  12  38  36  78  70  16  4   4   6
6   6   6 114  34  32  32  30 110  74  34  32

```

32	30	0	0	4	4	70	70	6	6	110	110
70	100	106	26	70	24	102	102	92	0	0	66
98	46	44	22	4	4	96	120	6	6	126	126
42	40	116	116	10	10	10	10	38	36	22	22
38	36	112	118	60	46	44	44	42	116	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (124 , 30) IS SELECTED.

CUBE (115 , 48) HAS BEEN FOUND.

CUBE (115 , 57) HAS BEEN FOUND.

WATCH FOR 48 AND 57

CUBE (115 , 58) HAS BEEN FOUND.

WATCH FOR 48 AND 58

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	5	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	25	9999	9999	9999	9999	9999
9999	9999	9999	7	9999	9999	9999	9999	3	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	15	11	
9999	3	7	9999	9999	9999	5	9	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	3	9999	9999	1	3	
9999	9999	5	8	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	3	9999	9999	9999	9999	9999	7	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

84	52	50	14	14	14	14	100	62	62	62	72
0	0	48	46	70	34	32	76	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	66
0	0	12	12	38	36	80	70	16	4	4	6
6	6	6	118	34	32	32	30	116	74	34	32
32	30	0	0	4	4	70	70	6	6	112	112
70	104	110	26	70	24	108	108	92	0	0	66
98	46	44	22	4	4	96	124	6	6	132	132
42	40	122	122	10	10	10	10	38	36	22	22
38	36	118	124	60	46	44	44	42	118	18	18

0 0 0 0 0 0 0 0 0 0 0 0
 0 0 0 0 0 0 0 0 0 0 0 0
 0 0 0 0 0 0 0 0 0 0 0 0
 0 0 0 0 0 0 0 0 0 0 0 0

CUBE (115 , 48) IS SELECTED,

CUBE (55 , 3) HAS BEEN FOUND.

CUBE (55 , 5) HAS BEEN FOUND.

WATCH FOR 3 AND 5

CUBE (55 , 17) HAS BEEN FOUND.

WATCH FOR 5 AND 17

CUBE (55 , 27) HAS BEEN FOUND.

WATCH FOR 17 AND 27

CUBE (55 , 29) HAS BEEN FOUND.

WATCH FOR 17 AND 29

CUBE (55 , 57) HAS BEEN FOUND.

WATCH FOR 17 AND 57

CHECKING POINTS

9999 9999 9999 9999 9999 9999 9999 5 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 7 9999 9999 9999 9999 3 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 15 11
 9999 3 7 9999 9999 9999 5 9 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 4 9999 9999 1 3
 9999 9999 5 8 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 8 3 9999 9999 9999 9999 9999 7 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999 9999
 9999 9999 9999 9999

ACCUMULATED AWAYNESS :

84 52 50 14 14 14 14 104 62 62 62 72
 0 0 48 46 70 34 32 76 4 4 6 6
 6 6 16 16 14 14 14 14 20 50 18 66
 0 0 12 12 38 36 80 70 16 4 4 6

6	6	6	124	34	32	32	30	124	74	34	32
32	30	0	0	4	4	70	70	6	6	118	116
70	108	118	26	70	24	116	114	92	0	0	66
98	46	44	22	4	4	96	126	6	6	138	136
42	40	128	126	10	10	10	10	38	36	22	22
38	36	122	132	60	46	44	44	42	122	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0								

CUBE (55 , 17) IS SELECTED.

CUBE (128 , 161) HAS BEEN FOUND.

CUBE (128 , 224) HAS BEEN FOUND.

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	5	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	7	9999	9999	9999	9999	3	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	15	11	
9999	3	7	9999	9999	9999	5	9	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	5	9999	9999	9999	9999	9999
9999	9999	5	8	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	8	4	9999	9999	9999	9999	9999	8	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

84	52	50	14	14	14	14	108	62	62	62	72
0	0	48	46	70	34	32	76	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	66
0	0	12	12	38	36	80	70	16	4	4	6
6	6	6	128	34	32	32	30	130	74	34	32
32	30	0	0	4	4	70	70	6	6	122	120
70	116	124	26	70	24	124	122	92	0	0	66
98	46	44	22	4	4	96	128	6	6	138	136
42	40	132	130	10	10	10	10	38	36	22	22
38	36	126	134	60	46	44	44	42	124	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0								

CUBE (128 , 161) IS SELECTED.

ACCUMULATED AWAYNESS :

84	52	50	14	14	14	14	124	62	62	62	72
0	0	48	46	70	34	32	76	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	66
0	0	12	12	38	36	80	70	16	4	4	6
6	6	6	134	34	32	32	30	138	74	34	32
32	30	0	0	4	4	70	70	6	6	128	136
70	116	132	26	70	24	144	142	92	0	0	66
98	46	44	22	4	4	96	154	6	6	138	136
42	40	160	158	10	10	10	10	38	36	22	22
38	36	148	148	60	46	44	44	42	150	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (9 , 5) IS SELECTED.

CUBE (145 , 1) HAS BEEN FOUND.

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	12
9999	9999	9999	9999	9999	9999	9999	6	9	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	5	9	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

84	52	50	14	14	14	14	124	62	62	62	72
0	0	48	46	70	34	32	76	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	66
0	0	12	12	38	36	80	70	16	4	4	6
6	6	6	134	34	32	32	30	138	74	34	32
32	30	0	0	4	4	70	70	6	6	128	140
70	116	132	26	70	24	154	150	92	0	0	66
98	46	44	22	4	4	96	154	6	6	138	136
42	40	164	160	10	10	10	10	38	36	22	22
38	36	154	148	60	46	44	44	42	154	18	18

CUBE (108 , 72) HAS BEEN FOUND.

CUBE (108 , 77) HAS BEEN FOUND.

CUBE (108 , 92) HAS BEEN FOUND.

CUBE (108 , 105) HAS BEEN FOUND.

WATCH FOR 77 AND 105

CHECKING POINTS

9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	13
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9	9999	9999	9999	9999	9999	9999	9999	9	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999
9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999	9999

ACCUMULATED AWAYNESS :

84	52	50	14	14	14	14	124	62	62	62	72
0	0	48	46	70	34	32	76	4	4	6	6
6	6	16	16	14	14	14	14	20	50	18	66
0	0	12	12	38	36	80	70	16	4	4	6
6	6	6	134	34	32	32	30	138	74	34	32
32	30	0	0	4	4	70	70	6	6	128	150
70	116	132	26	70	24	162	158	92	0	0	66
98	46	44	22	4	4	96	154	6	6	138	136
42	40	164	160	10	10	10	10	38	36	22	22
38	36	172	148	60	46	44	44	42	166	18	18
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	0	0	0	0	0	0	0	0

CUBE (108 , 105) IS SELECTED;

CUBE (176 , 50) HAS BEEN FOUND.

CUBE (176 , 162) HAS BEEN FOUND.

WATCH FOR 50 AND 162

APPENDIX I

EXAMPLE OF 9 VARIABLES, 5 ON-SET MINTERMS, AND 503 DON'T
CARE MINTERMS.

TOATAL NUMBER OF ON-SET AND DONT-CARES = 508
 ON-SET = 5 DONT-CARES = 503
 NUMBERS OF VARIABLES = 9

MINT.# 1 = 1	MINT.# 2 = 2	MINT.# 3 = 4
MINT.# 4 = 8	MINT.# 5 = 256	MINT.# 6 = 0
MINT.# 7 = 3	MINT.# 8 = 5	MINT.# 9 = 6
MINT.# 10 = 7	MINT.# 11 = 9	MINT.# 12 = 10
MINT.# 13 = 11	MINT.# 14 = 12	MINT.# 15 = 13
MINT.# 16 = 14	MINT.# 17 = 15	MINT.# 18 = 17
MINT.# 19 = 18	MINT.# 20 = 19	MINT.# 21 = 20
MINT.# 22 = 21	MINT.# 23 = 22	MINT.# 24 = 23
MINT.# 25 = 24	MINT.# 26 = 25	MINT.# 27 = 26
MINT.# 28 = 27	MINT.# 29 = 28	MINT.# 30 = 29
MINT.# 31 = 30	MINT.# 32 = 31	MINT.# 33 = 33
MINT.# 34 = 34	MINT.# 35 = 35	MINT.# 36 = 36
MINT.# 37 = 37	MINT.# 38 = 38	MINT.# 39 = 39
MINT.# 40 = 40	MINT.# 41 = 41	MINT.# 42 = 42
MINT.# 43 = 43	MINT.# 44 = 44	MINT.# 45 = 45
MINT.# 46 = 46	MINT.# 47 = 47	MINT.# 48 = 48
MINT.# 49 = 49	MINT.# 50 = 50	MINT.# 51 = 51
MINT.# 52 = 52	MINT.# 53 = 53	MINT.# 54 = 54
MINT.# 55 = 55	MINT.# 56 = 56	MINT.# 57 = 57
MINT.# 58 = 58	MINT.# 59 = 59	MINT.# 60 = 60
MINT.# 61 = 61	MINT.# 62 = 62	MINT.# 63 = 63
MINT.# 64 = 65	MINT.# 65 = 66	MINT.# 66 = 67
MINT.# 67 = 68	MINT.# 68 = 69	MINT.# 69 = 70
MINT.# 70 = 71	MINT.# 71 = 72	MINT.# 72 = 73
MINT.# 73 = 74	MINT.# 74 = 75	MINT.# 75 = 76
MINT.# 76 = 77	MINT.# 77 = 78	MINT.# 78 = 79
MINT.# 79 = 80	MINT.# 80 = 81	MINT.# 81 = 82
MINT.# 82 = 83	MINT.# 83 = 84	MINT.# 84 = 85
MINT.# 85 = 86	MINT.# 86 = 87	MINT.# 87 = 88
MINT.# 88 = 89	MINT.# 89 = 90	MINT.# 90 = 91
MINT.# 91 = 92	MINT.# 92 = 93	MINT.# 93 = 94
MINT.# 94 = 95	MINT.# 95 = 96	MINT.# 96 = 97
MINT.# 97 = 98	MINT.# 98 = 99	MINT.# 99 = 100
MINT.# 100 = 101	MINT.# 101 = 102	MINT.# 102 = 103
MINT.# 103 = 104	MINT.# 104 = 105	MINT.# 105 = 106
MINT.# 106 = 107	MINT.# 107 = 108	MINT.# 108 = 109
MINT.# 109 = 110	MINT.# 110 = 111	MINT.# 111 = 112
MINT.# 112 = 113	MINT.# 113 = 114	MINT.# 114 = 115
MINT.# 115 = 116	MINT.# 116 = 117	MINT.# 117 = 118
MINT.# 118 = 119	MINT.# 119 = 120	MINT.# 120 = 121
MINT.# 121 = 122	MINT.# 122 = 123	MINT.# 123 = 124
MINT.# 124 = 125	MINT.# 125 = 126	MINT.# 126 = 127
MINT.# 127 = 129	MINT.# 128 = 130	MINT.# 129 = 131
MINT.# 130 = 132	MINT.# 131 = 133	MINT.# 132 = 134
MINT.# 133 = 135	MINT.# 134 = 136	MINT.# 135 = 137
MINT.# 136 = 138	MINT.# 137 = 139	MINT.# 138 = 140
MINT.# 139 = 141	MINT.# 140 = 142	MINT.# 141 = 143

MINT.# 142 = 144	MINT.# 143 = 145	MINT.# 144 = 146
MINT.# 145 = 147	MINT.# 146 = 148	MINT.# 147 = 149
MINT.# 148 = 150	MINT.# 149 = 151	MINT.# 150 = 152
MINT.# 151 = 153	MINT.# 152 = 154	MINT.# 153 = 155
MINT.# 154 = 156	MINT.# 155 = 157	MINT.# 156 = 158
MINT.# 157 = 159	MINT.# 158 = 160	MINT.# 159 = 161
MINT.# 160 = 162	MINT.# 161 = 163	MINT.# 162 = 164
MINT.# 163 = 165	MINT.# 164 = 166	MINT.# 165 = 167
MINT.# 166 = 168	MINT.# 167 = 169	MINT.# 168 = 170
MINT.# 169 = 171	MINT.# 170 = 172	MINT.# 171 = 173
MINT.# 172 = 174	MINT.# 173 = 175	MINT.# 174 = 176
MINT.# 175 = 177	MINT.# 176 = 178	MINT.# 177 = 179
MINT.# 178 = 180	MINT.# 179 = 181	MINT.# 180 = 182
MINT.# 181 = 183	MINT.# 182 = 184	MINT.# 183 = 185
MINT.# 184 = 186	MINT.# 185 = 187	MINT.# 186 = 188
MINT.# 187 = 189	MINT.# 188 = 190	MINT.# 189 = 191
MINT.# 190 = 192	MINT.# 191 = 193	MINT.# 192 = 194
MINT.# 193 = 195	MINT.# 194 = 196	MINT.# 195 = 197
MINT.# 196 = 198	MINT.# 197 = 199	MINT.# 198 = 200
MINT.# 199 = 201	MINT.# 200 = 202	MINT.# 201 = 203
MINT.# 202 = 204	MINT.# 203 = 205	MINT.# 204 = 206
MINT.# 205 = 207	MINT.# 206 = 208	MINT.# 207 = 209
MINT.# 208 = 210	MINT.# 209 = 211	MINT.# 210 = 212
MINT.# 211 = 213	MINT.# 212 = 214	MINT.# 213 = 215
MINT.# 214 = 216	MINT.# 215 = 217	MINT.# 216 = 218
MINT.# 217 = 219	MINT.# 218 = 220	MINT.# 219 = 221
MINT.# 220 = 222	MINT.# 221 = 223	MINT.# 222 = 224
MINT.# 223 = 225	MINT.# 224 = 226	MINT.# 225 = 227
MINT.# 226 = 228	MINT.# 227 = 229	MINT.# 228 = 230
MINT.# 229 = 231	MINT.# 230 = 232	MINT.# 231 = 233
MINT.# 232 = 234	MINT.# 233 = 235	MINT.# 234 = 236
MINT.# 235 = 237	MINT.# 236 = 238	MINT.# 237 = 239
MINT.# 238 = 240	MINT.# 239 = 241	MINT.# 240 = 242
MINT.# 241 = 243	MINT.# 242 = 244	MINT.# 243 = 245
MINT.# 244 = 246	MINT.# 245 = 247	MINT.# 246 = 248
MINT.# 247 = 249	MINT.# 248 = 250	MINT.# 249 = 251
MINT.# 250 = 252	MINT.# 251 = 253	MINT.# 252 = 254
MINT.# 253 = 255	MINT.# 254 = 257	MINT.# 255 = 258
MINT.# 256 = 259	MINT.# 257 = 260	MINT.# 258 = 261
MINT.# 259 = 262	MINT.# 260 = 263	MINT.# 261 = 264
MINT.# 262 = 265	MINT.# 263 = 266	MINT.# 264 = 267
MINT.# 265 = 268	MINT.# 266 = 269	MINT.# 267 = 270
MINT.# 268 = 271	MINT.# 269 = 272	MINT.# 270 = 273
MINT.# 271 = 274	MINT.# 272 = 275	MINT.# 273 = 276
MINT.# 274 = 277	MINT.# 275 = 278	MINT.# 276 = 279
MINT.# 277 = 280	MINT.# 278 = 281	MINT.# 279 = 282
MINT.# 280 = 283	MINT.# 281 = 284	MINT.# 282 = 285
MINT.# 283 = 286	MINT.# 284 = 287	MINT.# 285 = 288
MINT.# 286 = 289	MINT.# 287 = 290	MINT.# 288 = 291
MINT.# 289 = 292	MINT.# 290 = 293	MINT.# 291 = 294
MINT.# 292 = 295	MINT.# 293 = 296	MINT.# 294 = 297
MINT.# 295 = 298	MINT.# 296 = 299	MINT.# 297 = 300

MINT.# 298	=	301	MINT.# 299	=	302	MINT.# 300	=	303
MINT.# 301	=	304	MINT.# 302	=	305	MINT.# 303	=	306
MINT.# 304	=	307	MINT.# 305	=	308	MINT.# 306	=	309
MINT.# 307	=	310	MINT.# 308	=	311	MINT.# 309	=	312
MINT.# 310	=	313	MINT.# 311	=	314	MINT.# 312	=	315
MINT.# 313	=	316	MINT.# 314	=	317	MINT.# 315	=	318
MINT.# 316	=	319	MINT.# 317	=	320	MINT.# 318	=	321
MINT.# 319	=	322	MINT.# 320	=	323	MINT.# 321	=	324
MINT.# 322	=	325	MINT.# 323	=	326	MINT.# 324	=	327
MINT.# 325	=	328	MINT.# 326	=	329	MINT.# 327	=	330
MINT.# 328	=	331	MINT.# 329	=	332	MINT.# 330	=	333
MINT.# 331	=	334	MINT.# 332	=	335	MINT.# 333	=	336
MINT.# 334	=	337	MINT.# 335	=	338	MINT.# 336	=	339
MINT.# 337	=	340	MINT.# 338	=	341	MINT.# 339	=	342
MINT.# 340	=	343	MINT.# 341	=	344	MINT.# 342	=	345
MINT.# 343	=	346	MINT.# 344	=	347	MINT.# 345	=	348
MINT.# 346	=	349	MINT.# 347	=	350	MINT.# 348	=	351
MINT.# 349	=	352	MINT.# 350	=	353	MINT.# 351	=	354
MINT.# 352	=	355	MINT.# 353	=	356	MINT.# 354	=	357
MINT.# 355	=	358	MINT.# 356	=	359	MINT.# 357	=	360
MINT.# 358	=	361	MINT.# 359	=	362	MINT.# 360	=	363
MINT.# 361	=	364	MINT.# 362	=	365	MINT.# 363	=	366
MINT.# 364	=	367	MINT.# 365	=	368	MINT.# 366	=	369
MINT.# 367	=	370	MINT.# 368	=	371	MINT.# 369	=	372
MINT.# 370	=	373	MINT.# 371	=	374	MINT.# 372	=	375
MINT.# 373	=	376	MINT.# 374	=	377	MINT.# 375	=	378
MINT.# 376	=	379	MINT.# 377	=	380	MINT.# 378	=	381
MINT.# 379	=	382	MINT.# 380	=	383	MINT.# 381	=	384
MINT.# 382	=	385	MINT.# 383	=	386	MINT.# 384	=	387
MINT.# 385	=	388	MINT.# 386	=	389	MINT.# 387	=	390
MINT.# 388	=	391	MINT.# 389	=	392	MINT.# 390	=	393
MINT.# 391	=	394	MINT.# 392	=	395	MINT.# 393	=	396
MINT.# 394	=	397	MINT.# 395	=	398	MINT.# 396	=	399
MINT.# 397	=	400	MINT.# 398	=	401	MINT.# 399	=	402
MINT.# 400	=	403	MINT.# 401	=	404	MINT.# 402	=	405
MINT.# 403	=	406	MINT.# 404	=	407	MINT.# 405	=	408
MINT.# 406	=	409	MINT.# 407	=	410	MINT.# 408	=	411
MINT.# 409	=	412	MINT.# 410	=	413	MINT.# 411	=	414
MINT.# 412	=	415	MINT.# 413	=	416	MINT.# 414	=	417
MINT.# 415	=	418	MINT.# 416	=	419	MINT.# 417	=	420
MINT.# 418	=	421	MINT.# 419	=	422	MINT.# 420	=	423
MINT.# 421	=	424	MINT.# 422	=	425	MINT.# 423	=	426
MINT.# 424	=	427	MINT.# 425	=	428	MINT.# 426	=	429
MINT.# 427	=	430	MINT.# 428	=	431	MINT.# 429	=	432
MINT.# 430	=	433	MINT.# 431	=	434	MINT.# 432	=	435
MINT.# 433	=	436	MINT.# 434	=	437	MINT.# 435	=	438
MINT.# 436	=	439	MINT.# 437	=	440	MINT.# 438	=	441
MINT.# 439	=	442	MINT.# 440	=	443	MINT.# 441	=	444
MINT.# 442	=	445	MINT.# 443	=	446	MINT.# 444	=	447
MINT.# 445	=	448	MINT.# 446	=	449	MINT.# 447	=	450
MINT.# 448	=	451	MINT.# 449	=	452	MINT.# 450	=	453

MINT.# 451 = 454	MINT.# 452 = 455	MINT.# 453 = 456
MINT.# 454 = 457	MINT.# 453 = 458	MINT.# 456 = 459
MINT.# 457 = 460	MINT.# 458 = 461	MINT.# 459 = 462
MINT.# 460 = 463	MINT.# 461 = 464	MINT.# 462 = 465
MINT.# 463 = 466	MINT.# 464 = 467	MINT.# 465 = 468
MINT.# 466 = 469	MINT.# 467 = 470	MINT.# 468 = 471
MINT.# 469 = 472	MINT.# 470 = 473	MINT.# 471 = 474
MINT.# 472 = 475	MINT.# 473 = 476	MINT.# 474 = 477
MINT.# 475 = 478	MINT.# 476 = 479	MINT.# 477 = 480
MINT.# 478 = 481	MINT.# 479 = 482	MINT.# 480 = 483
MINT.# 481 = 484	MINT.# 482 = 485	MINT.# 483 = 486
MINT.# 484 = 487	MINT.# 485 = 488	MINT.# 486 = 489
MINT.# 487 = 490	MINT.# 488 = 491	MINT.# 489 = 492
MINT.# 490 = 493	MINT.# 491 = 494	MINT.# 492 = 495
MINT.# 493 = 496	MINT.# 494 = 497	MINT.# 495 = 498
MINT.# 496 = 499	MINT.# 497 = 500	MINT.# 498 = 501
MINT.# 499 = 502	MINT.# 500 = 503	MINT.# 501 = 504
MINT.# 502 = 505	MINT.# 503 = 506	MINT.# 504 = 507
MINT.# 505 = 508	MINT.# 506 = 509	MINT.# 507 = 510
MINT.# 508 = 511	MINT.#	

THE CUBE COVERING LIST :

THE 1-TH CUBE COVERING IS :
 8-CUBE = (256 , 511) = 1-----

THIS CUBE COVERS THESE MINTERMS :

256	511	257	258	259	260	261	262	263	264	265	266	267	268	269	270
271	272	273	274	275	276	277	278	279	280	281	282	283	284	285	286
287	288	289	290	291	292	293	294	295	296	297	298	299	300	301	302
303	304	305	306	307	308	309	310	311	312	313	314	315	316	317	318
319	320	321	322	323	324	325	326	327	328	329	330	331	332	333	334
335	336	337	338	339	340	341	342	343	344	345	346	347	348	349	350
351	352	353	354	355	356	357	358	359	360	361	362	363	364	365	366
367	368	369	370	371	372	373	374	375	376	377	378	379	380	381	382
383	384	385	386	387	388	389	390	391	392	393	394	395	396	397	398
399	400	401	402	403	404	405	406	407	408	409	410	411	412	413	414
415	416	417	418	419	420	421	422	423	424	425	426	427	428	429	430
431	432	433	434	435	436	437	438	439	440	441	442	443	444	445	446
447	448	449	450	451	452	453	454	455	456	457	458	459	460	461	462
463	464	465	466	467	468	469	470	471	472	473	474	475	476	477	478
479	480	481	482	483	484	485	486	487	488	489	490	491	492	493	494
495	496	497	498	499	500	501	502	503	504	505	506	507	508	509	510

THE 2-TH CUBE COVERING IS :
 8-CUBE = (8 , 511) = -----1---

THIS CUBE COVERS THESE MINTERMS :

8	511	9	10	11	12	13	14	15	24	25	26	27	28	29	30
31	40	41	42	43	44	45	46	47	56	57	58	59	60	61	62
63	72	73	74	75	76	77	78	79	88	89	90	91	92	93	94
95	104	105	106	107	108	109	110	111	120	121	122	123	124	125	126
127	136	137	138	139	140	141	142	143	152	153	154	155	156	157	158
159	168	169	170	171	172	173	174	175	184	185	186	187	188	189	190
191	200	201	202	203	204	205	206	207	216	217	218	219	220	221	222
223	232	233	234	235	236	237	238	239	248	249	250	251	252	253	254
255	264	265	266	267	268	269	270	271	280	281	282	283	284	285	286
287	296	297	298	299	300	301	302	303	312	313	314	315	316	317	318
319	328	329	330	331	332	333	334	335	344	345	346	347	348	349	350
351	360	361	362	363	364	365	366	367	376	377	378	379	380	381	382
383	392	393	394	395	396	397	398	399	408	409	410	411	412	413	414
415	424	425	426	427	428	429	430	431	440	441	442	443	444	445	446
447	456	457	458	459	460	461	462	463	472	473	474	475	476	477	478
479	488	489	490	491	492	493	494	495	504	505	506	507	508	509	510

THE 3-TH CUBE COVERING IS :
 8-CUBE = (2 , 511) = -----1-

THIS CUBE COVERS THESE MINTERMS :

2	511	3	6	7	10	11	14	15	18	19	22	23	26	27	30
31	34	35	38	39	42	43	46	47	50	51	54	55	58	59	62
63	66	67	70	71	74	75	78	79	82	83	86	87	90	91	94
95	98	99	102	103	106	107	110	111	114	115	118	119	122	123	126
127	130	131	134	135	138	139	142	143	146	147	150	151	154	155	158
159	162	163	166	167	170	171	174	175	178	179	182	183	186	187	190
191	194	195	198	199	202	203	206	207	210	211	214	215	218	219	222
223	226	227	230	231	234	235	238	239	242	243	246	247	250	251	254
255	258	259	262	263	266	267	270	271	274	275	278	279	282	283	286
287	290	291	294	295	298	299	302	303	306	307	310	311	314	315	318
319	322	323	326	327	330	331	334	335	338	339	342	343	346	347	350
351	354	355	358	359	362	363	366	367	370	371	374	375	378	379	382
383	386	387	390	391	394	395	398	399	402	403	406	407	410	411	414
415	418	419	422	423	426	427	430	431	434	435	438	439	442	443	446
447	450	451	454	455	458	459	462	463	466	467	470	471	474	475	478
479	482	483	486	487	490	491	494	495	498	499	502	503	506	507	510

THE 4-TH CUBE COVERING IS :
 8-CUBE = (4 , 511) = -----1--

THIS CUBE COVERS THESE MINTERMS :

4	511	5	6	7	12	13	14	15	20	21	22	23	28	29	30
31	36	37	38	39	44	45	46	47	52	53	54	55	60	61	62
63	68	69	70	71	76	77	78	79	84	85	86	87	92	93	94
95	100	101	102	103	108	109	110	111	116	117	118	119	124	125	126
127	132	133	134	135	140	141	142	143	148	149	150	151	156	157	158
159	164	165	166	167	172	173	174	175	180	181	182	183	188	189	190
191	196	197	198	199	204	205	206	207	212	213	214	215	220	221	222
223	228	229	230	231	236	237	238	239	244	245	246	247	252	253	254
255	260	261	262	263	268	269	270	271	276	277	278	279	284	285	286
287	292	293	294	295	300	301	302	303	308	309	310	311	316	317	318
319	324	325	326	327	332	333	334	335	340	341	342	343	348	349	350
351	356	357	358	359	364	365	366	367	372	373	374	375	380	381	382
383	388	389	390	391	396	397	398	399	404	405	406	407	412	413	414
415	420	421	422	423	428	429	430	431	436	437	438	439	444	445	446
447	452	453	454	455	460	461	462	463	468	469	470	471	476	477	478
479	484	485	486	487	492	493	494	495	500	501	502	503	508	509	510

THE 5-TH CUBE COVERING IS :

8-CUBE = (1 , 511) = -----1

THIS CUBE COVERS THESE MINTERMS :

1	511	3	5	7	9	11	13	15	17	19	21	23	25	27	29
31	33	35	37	39	41	43	45	47	49	51	53	55	57	59	61
63	65	67	69	71	73	75	77	79	81	83	85	87	89	91	93
95	97	99	101	103	105	107	109	111	113	115	117	119	121	123	125
127	129	131	133	135	137	139	141	143	145	147	149	151	153	155	157
159	161	163	165	167	169	171	173	175	177	179	181	183	185	187	189
191	193	195	197	199	201	203	205	207	209	211	213	215	217	219	221
223	225	227	229	231	233	235	237	239	241	243	245	247	249	251	253
255	257	259	261	263	265	267	269	271	273	275	277	279	281	283	285
287	289	291	293	295	297	299	301	303	305	307	309	311	313	315	317
319	321	323	325	327	329	331	333	335	337	339	341	343	345	347	349
351	353	355	357	359	361	363	365	367	369	371	373	375	377	379	381
383	385	387	389	391	393	395	397	399	401	403	405	407	409	411	413
415	417	419	421	423	425	427	429	431	433	435	437	439	441	443	445
447	449	451	453	455	457	459	461	463	465	467	469	471	473	475	477
479	481	483	485	487	489	491	493	495	497	499	501	503	505	507	509

THE FINAL CHECKING

105	105	105	105	105	109	309	309	309	409	309	309
409	309	409	409	509	209	209	309	209	309	309	409
209	309	309	409	309	409	409	509	209	209	309	209
309	309	409	209	309	309	409	309	409	409	509	109
209	209	309	209	309	309	409	209	309	309	409	309
409	409	509	209	209	309	209	309	309	409	209	309
309	409	309	409	409	509	109	209	209	309	209	309
309	409	209	309	309	409	309	409	409	509	109	209

